

Writing Datomic in Clojure

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INTERNATIONAL SOFTWARE DEVELOPMENT CONFERENCE

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Overview

- What is Datomic?
- Architecture
- Implementation Clojure Applied
- Summary



What is Datomic?

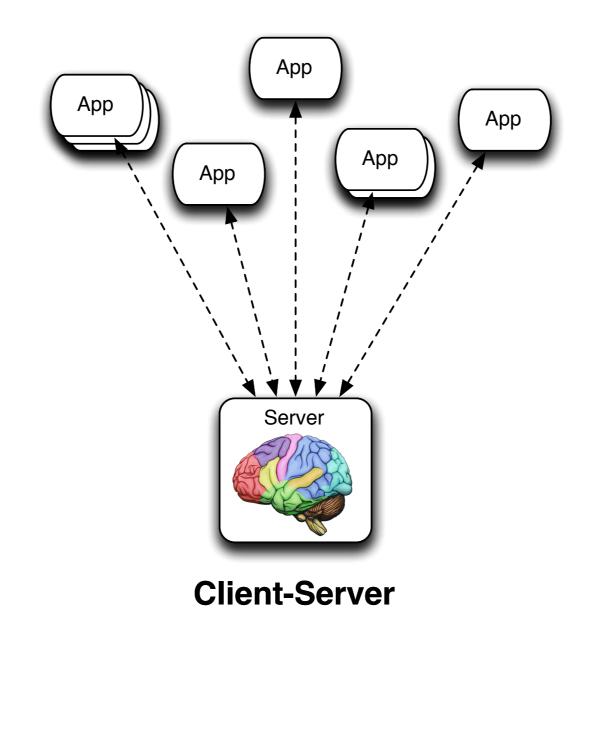
- A new kind of database
- Bringing data power *into* the application
- A sound model of information, with time
- Enabled by architectural advances



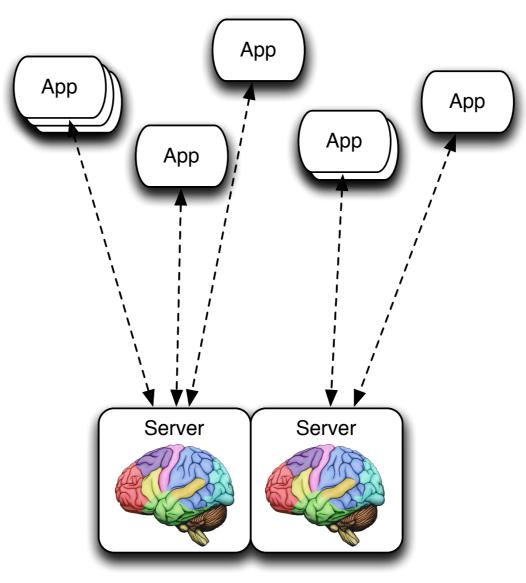
Why Datomic?

- Architecture
- Data Model



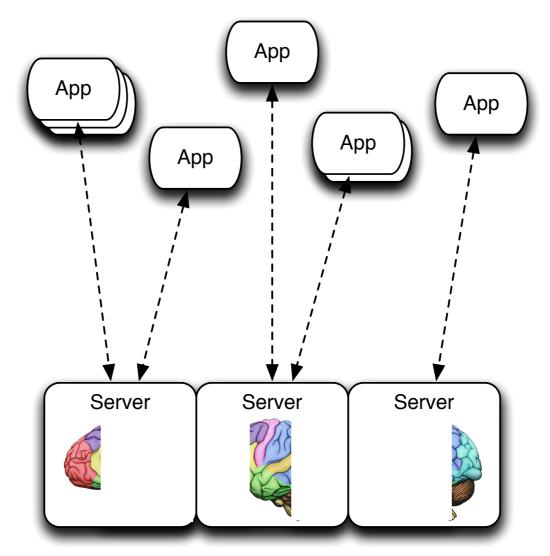






Clustered Client-Server

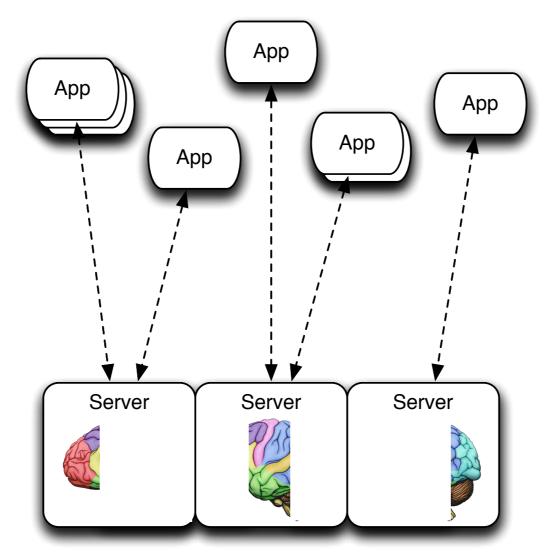




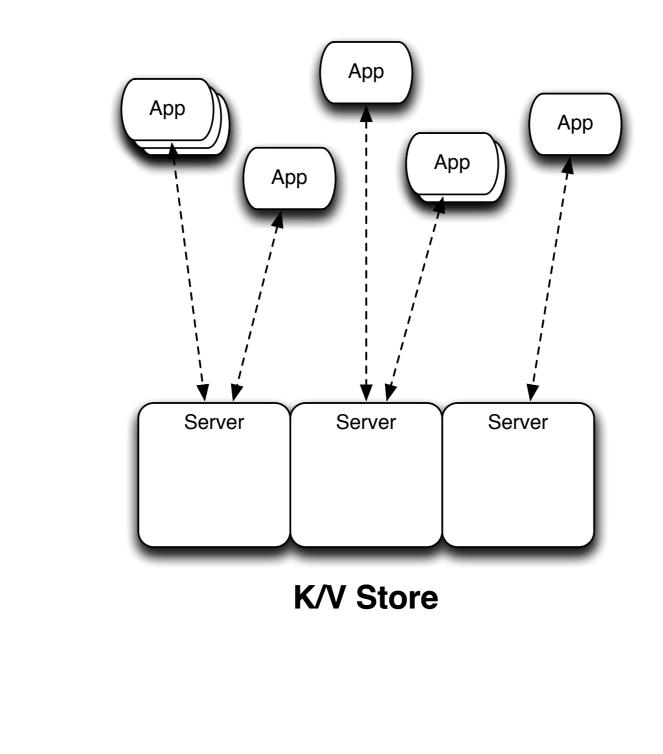
Queries Transactions Consistency Storage

Sharded Client-Server

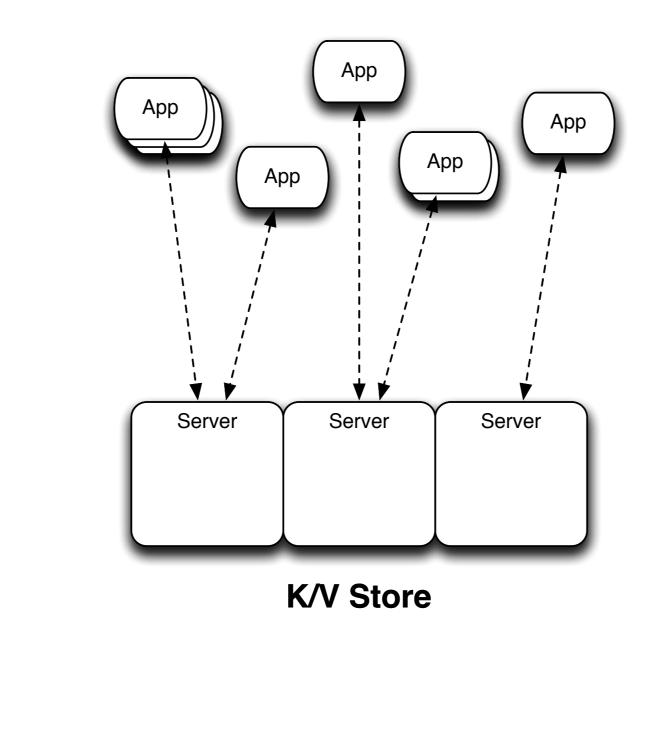




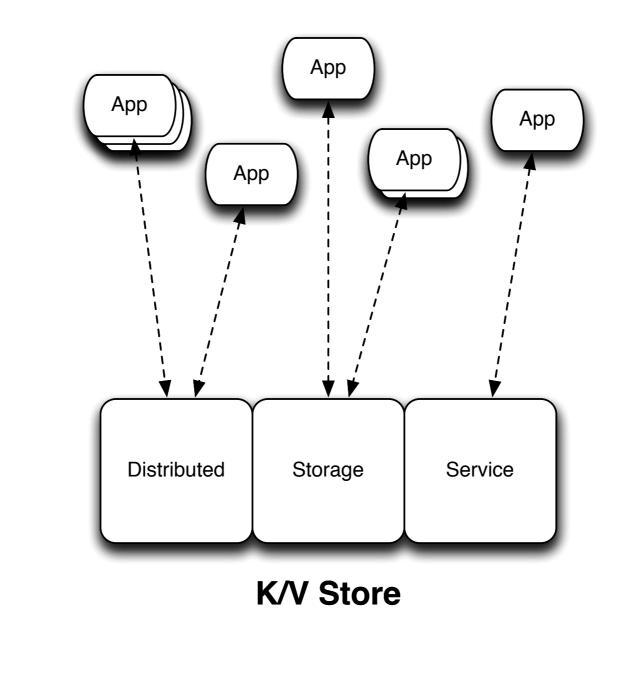
Sharded Client-Server





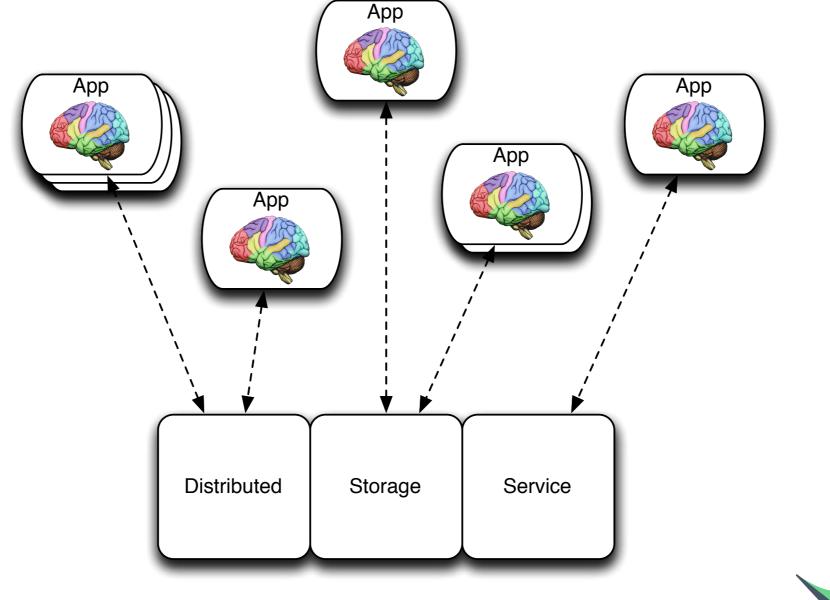




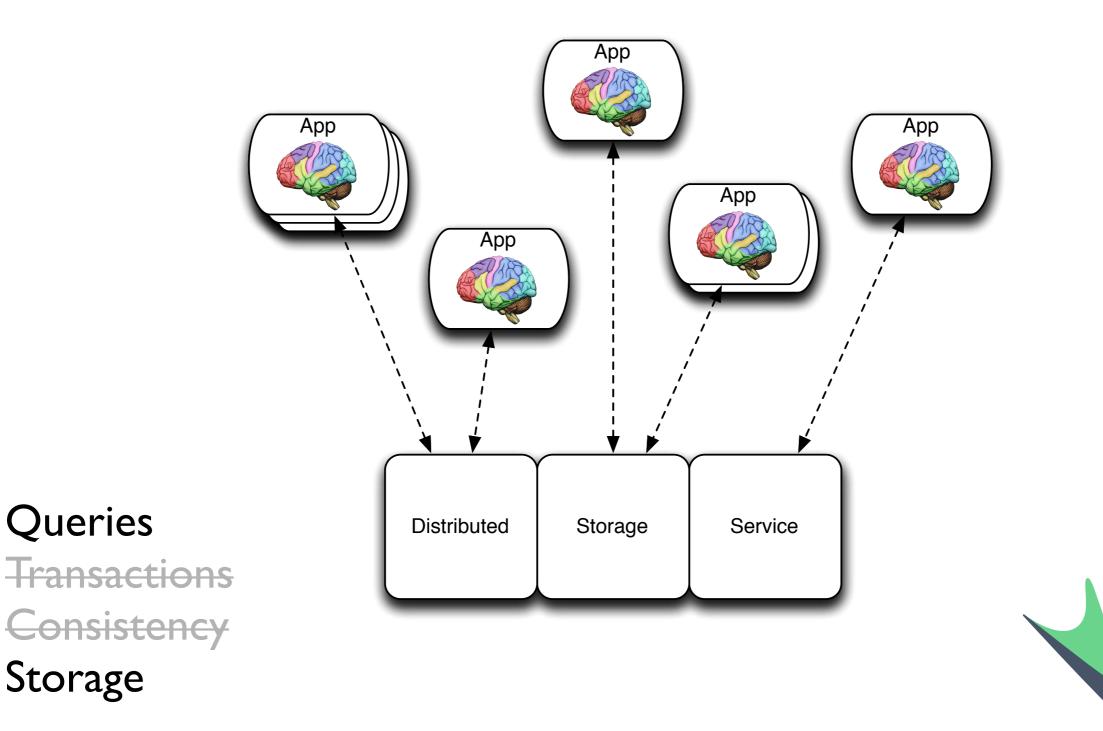


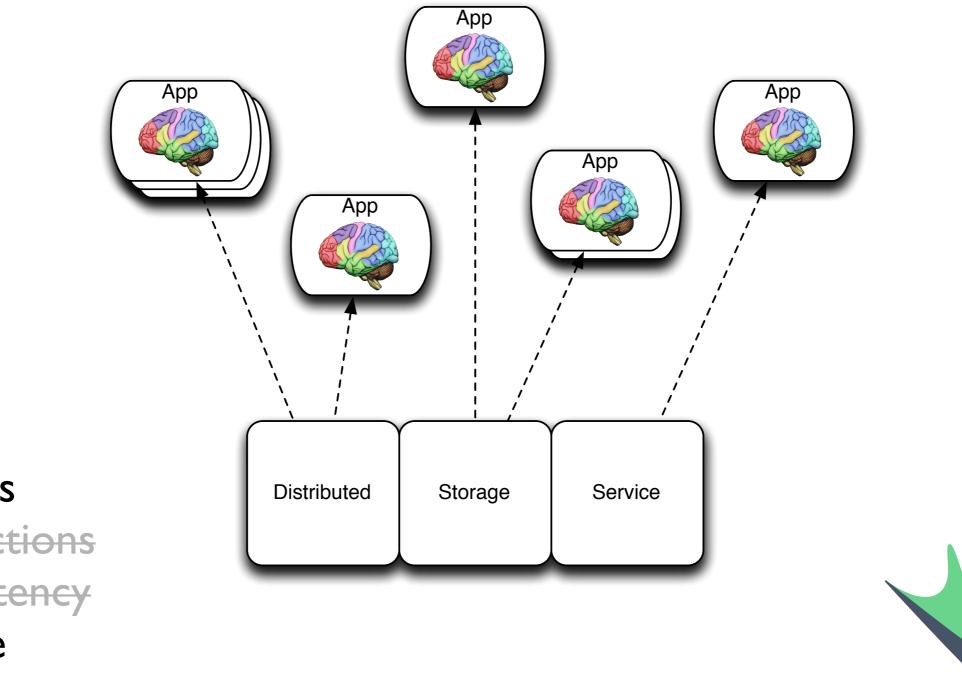


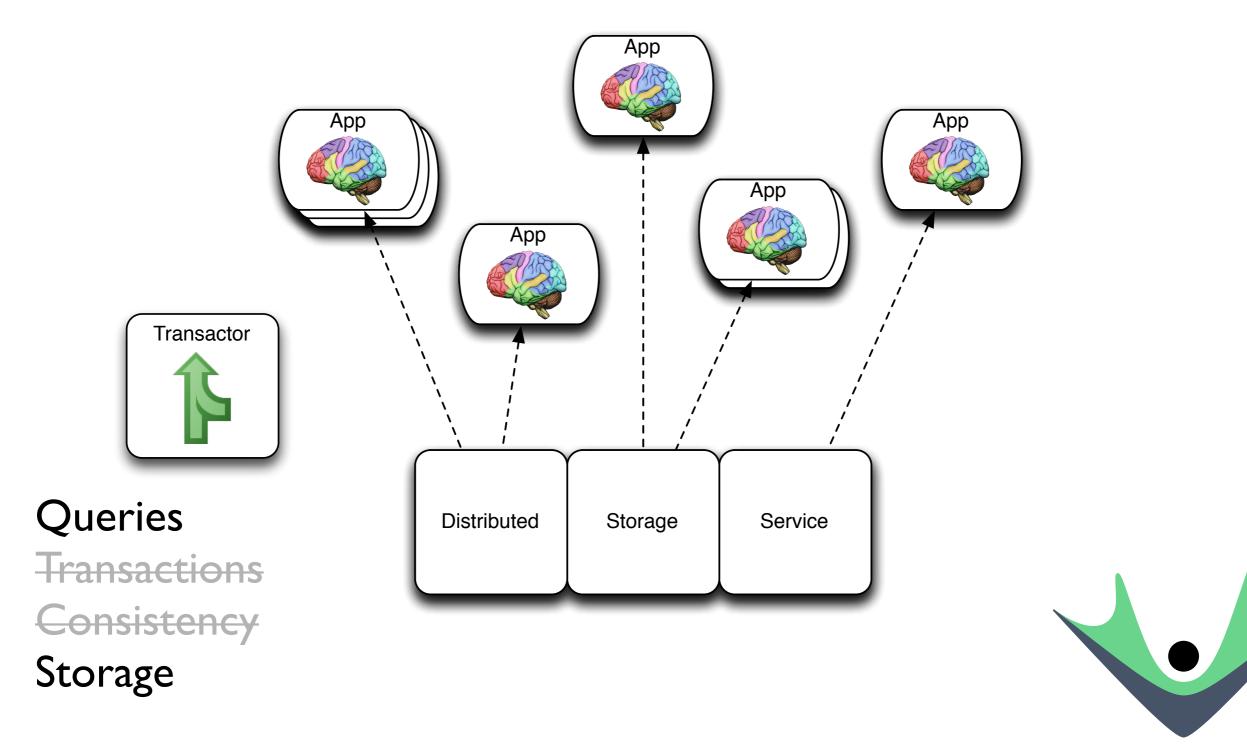


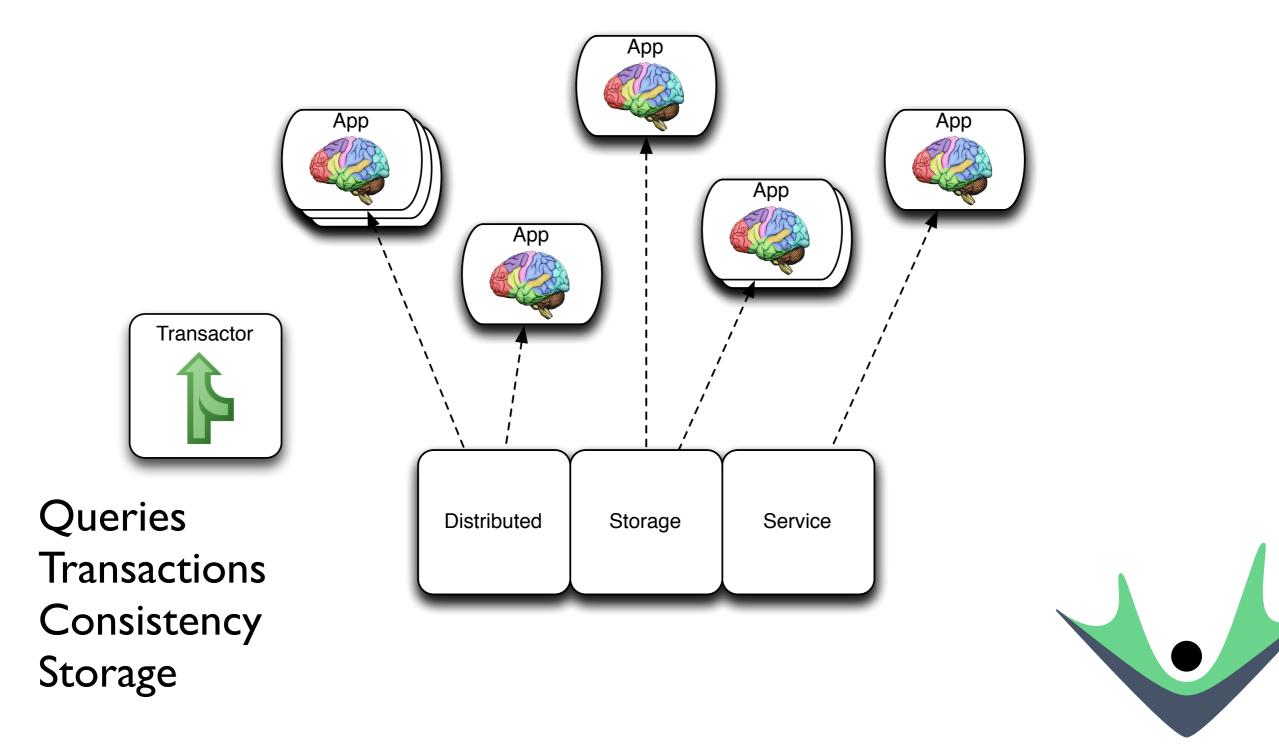


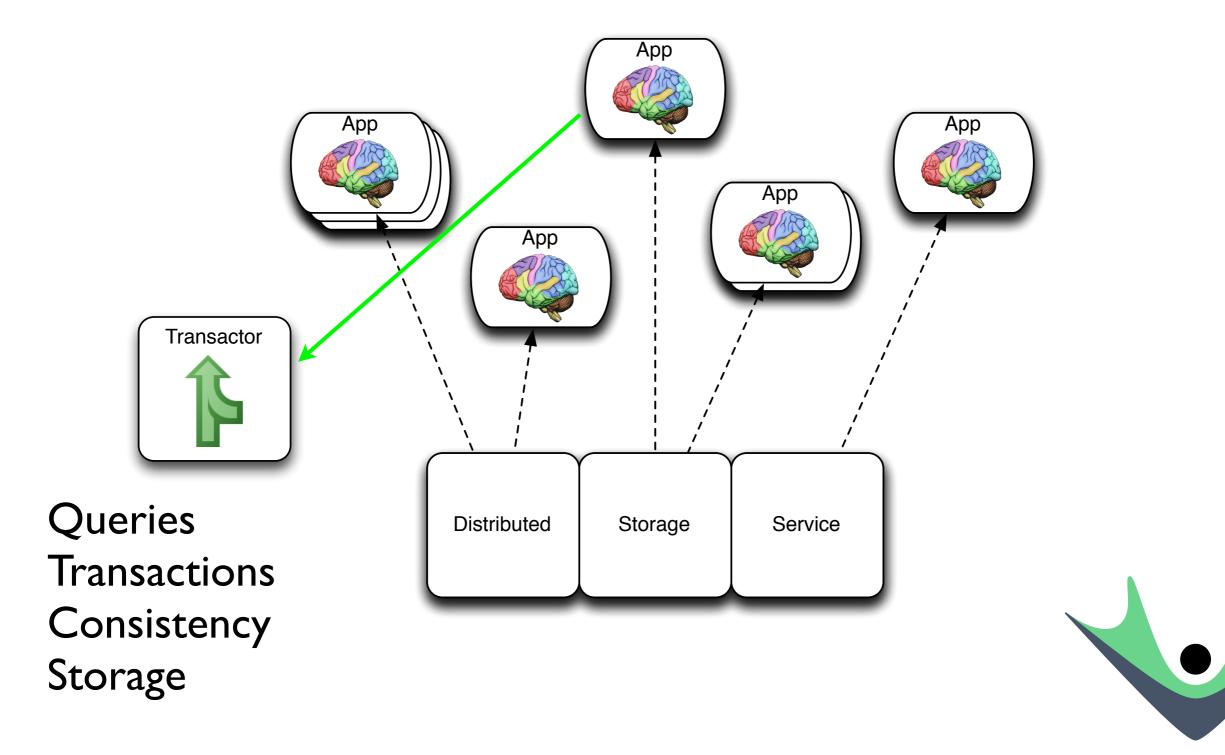


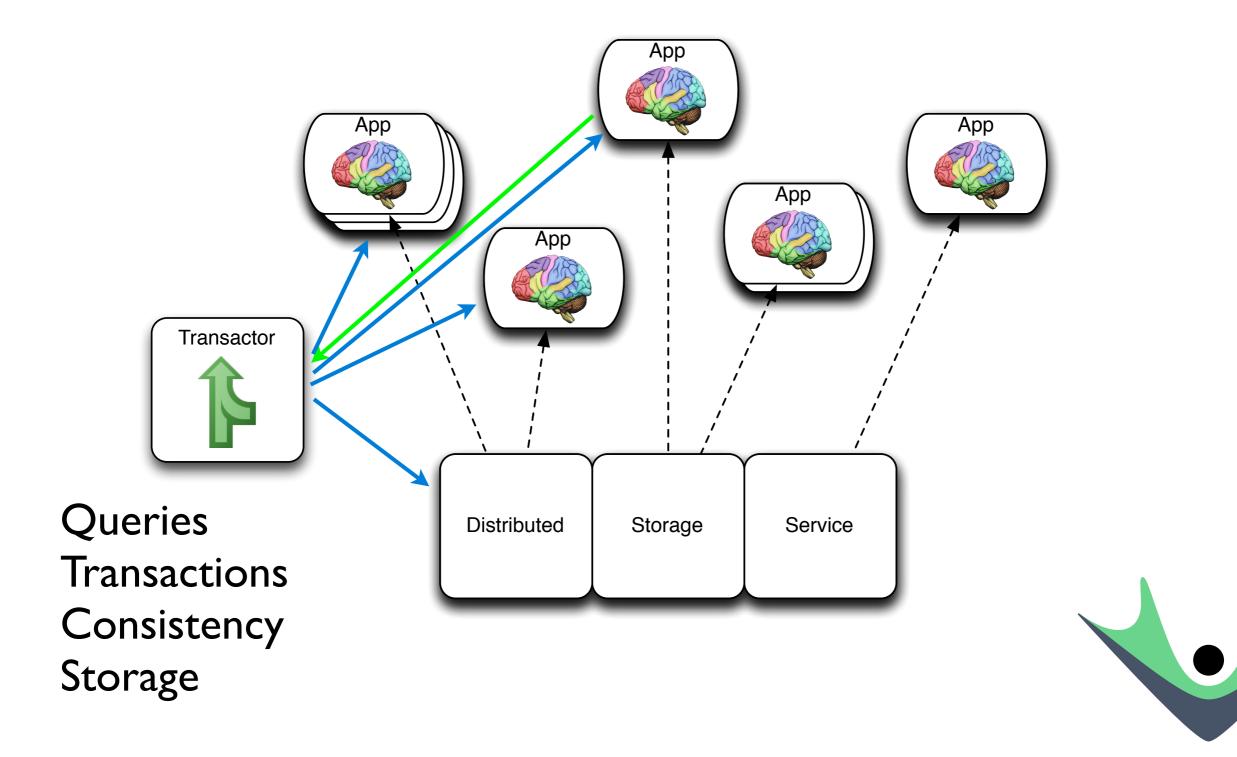


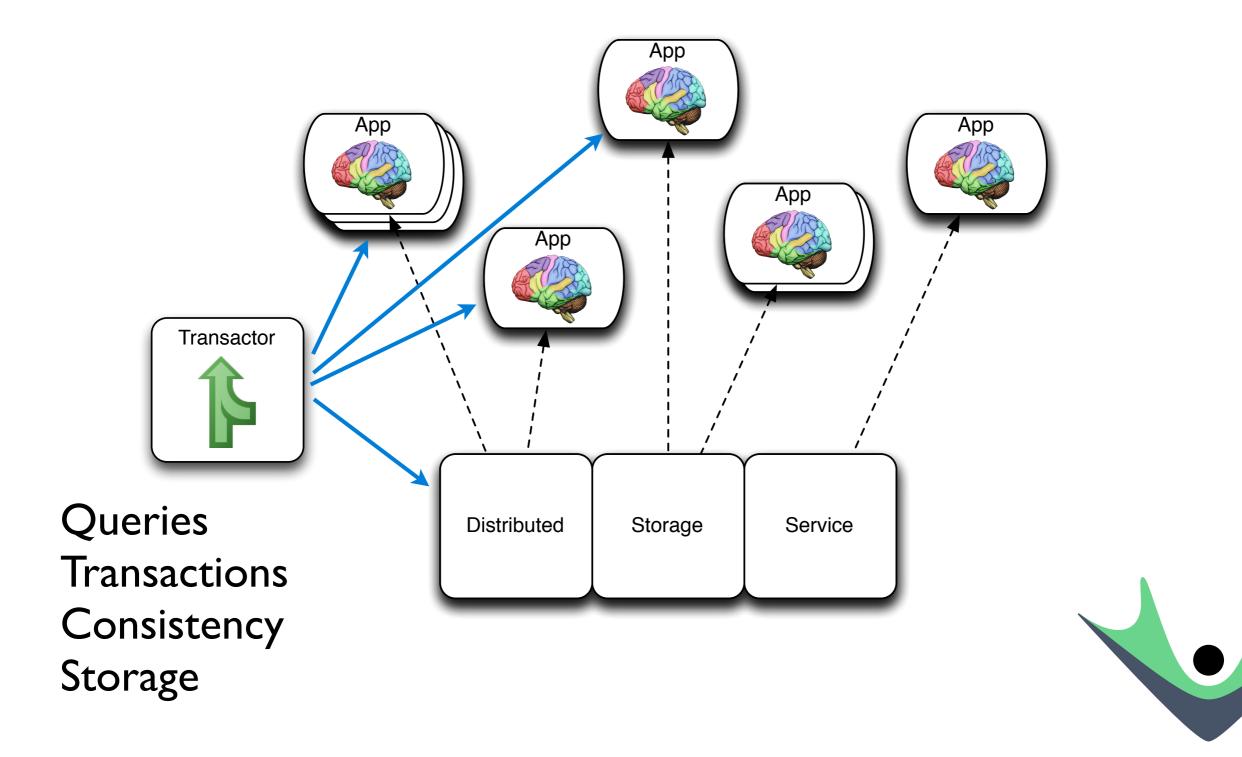


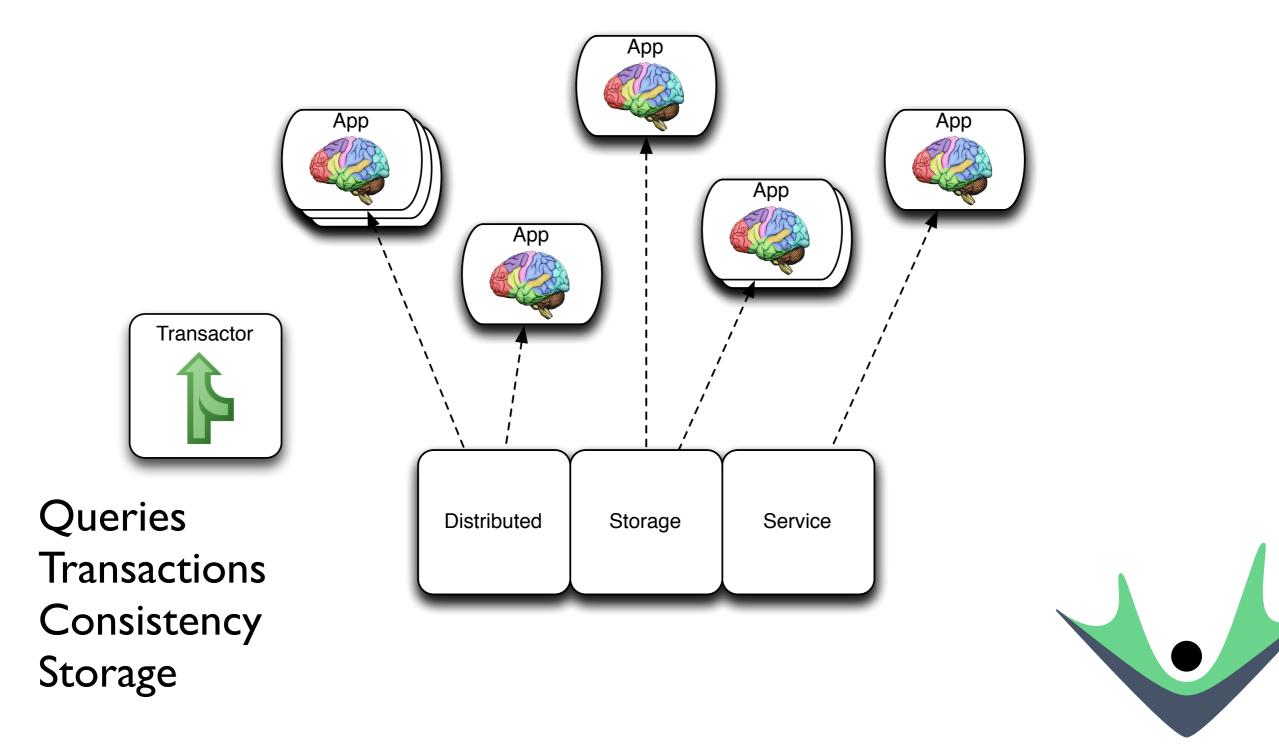




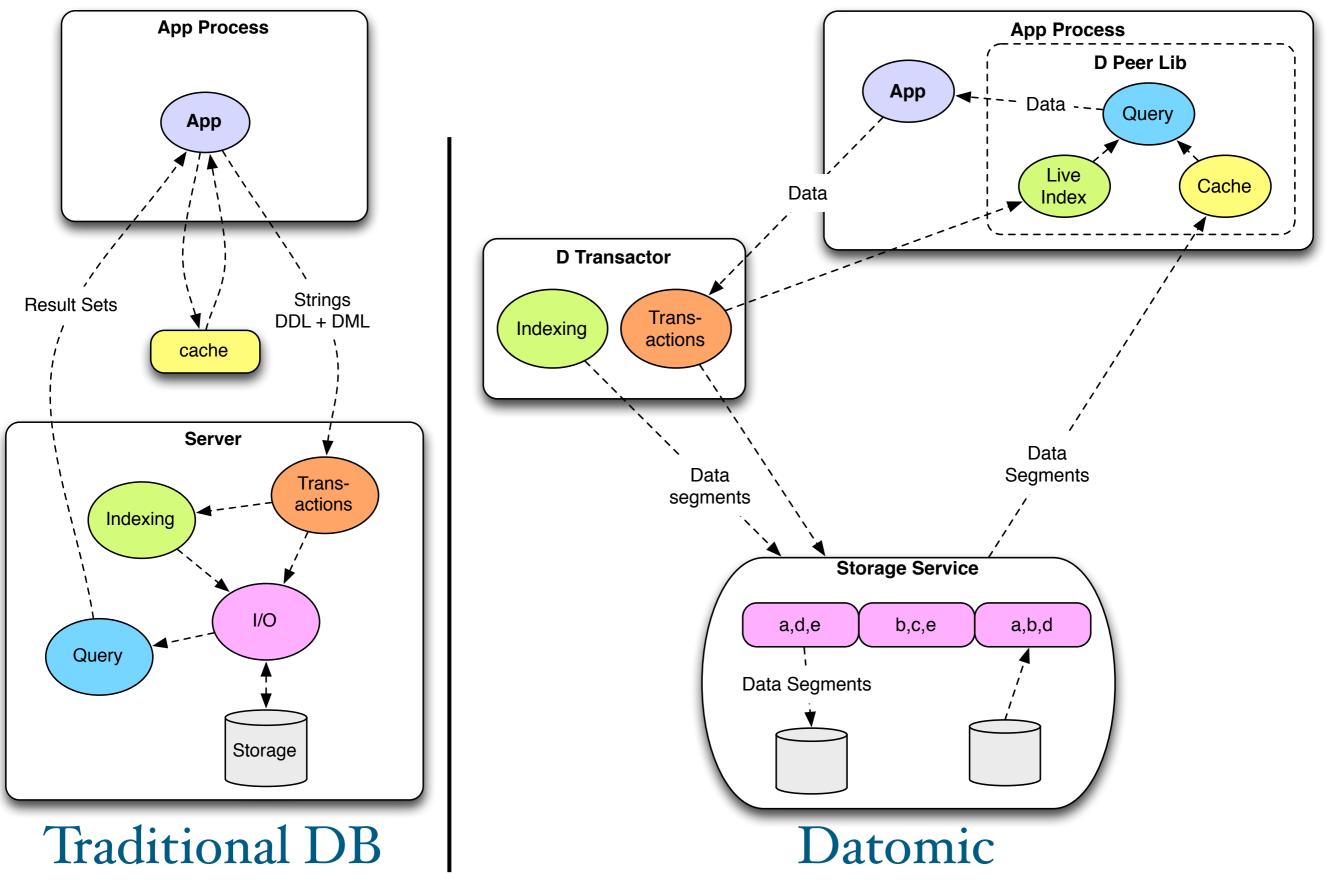








The Database, Deconstructed



Designed for the Cloud

- Ephemeral instances, unreliable disks
- Redundancy in storage service
- Leverages reliable storage services
 - e.g. DynamoDB



Elastic Scaling

- More peers, more power
- Fewer peers, less power, lower cost
- Demand-driven
 - No configuration



Get Your Own Brain

- Query, communication and memory engine
- Goes into your app, making it a peer
- The db is effectively local
- Ad hoc, long running queries ok



Logic

- Declarative search and business logic
- The query language is **Datalog**
 - Simple rules and data patterns
- Joins are implicit, meaning is evident
- db and non-db sources



Perception

- Obtain a queue of transactions
 - not just your own
- Query transactions for filtering/triggering



Consistency

- ACID transactions add new facts
- Database presented to app as a value
- Data in storage service is immutable



Programmability

- Transactions/Rules/Queries/Results are data
- Extensible types, predicates, etc
- Queries can invoke your code



A Database of Facts

- A single storage construct, the datom
 - Entity/Attribute/Value/Transaction
- Attribute definition is the only 'schema'



Adaptability

- Sparse, irregular, hierarchical data
- Single and multi-valued attributes
- No structural rigidity



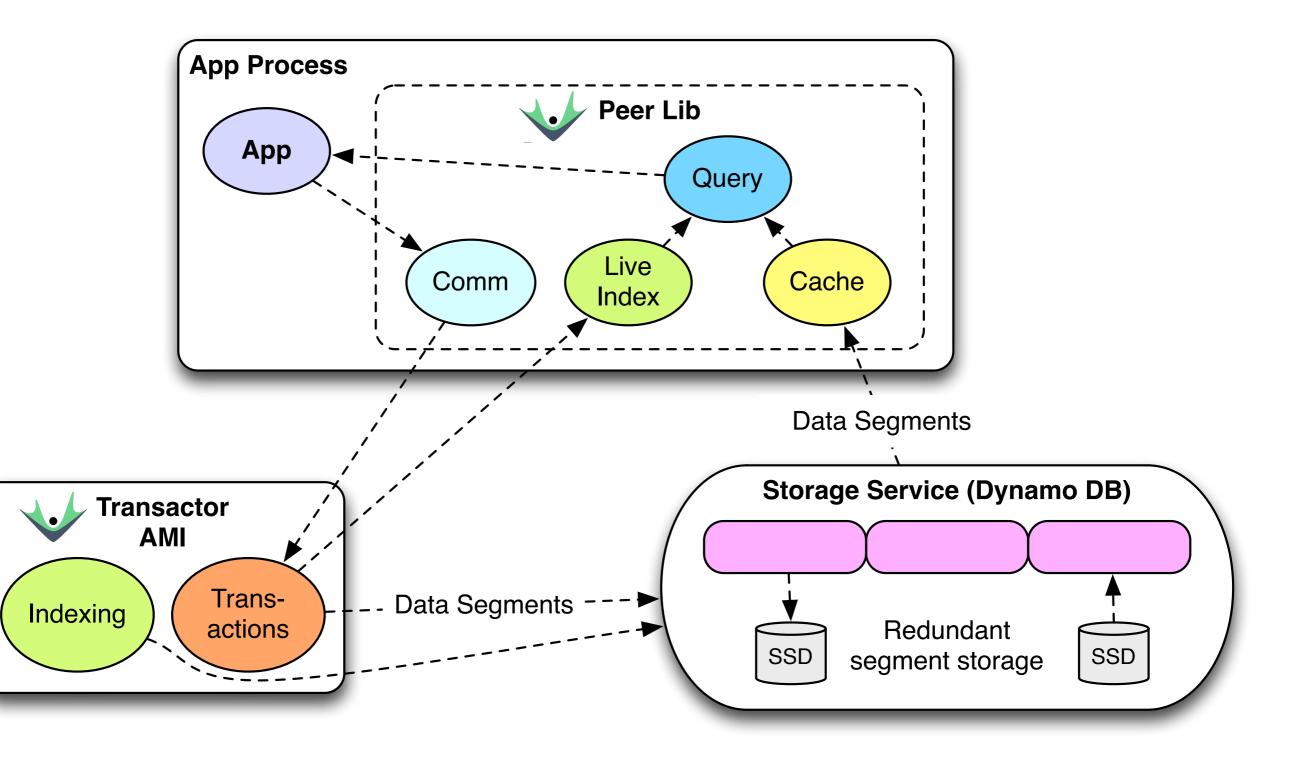
Time Built-in

- Every datom retains its transaction
- Transactions are totally ordered
- Transactions are first-class entities
- Get the db as-of, or since, a point in time



Implementation





State

- Immutable, expanding value
- Must be organized to support query
- Sorted set of facts
- Maintaining sort live in storage bad
 - BigTable mem + storage merge
 - occasional merge into storage
 - persistent trees



Memory Index

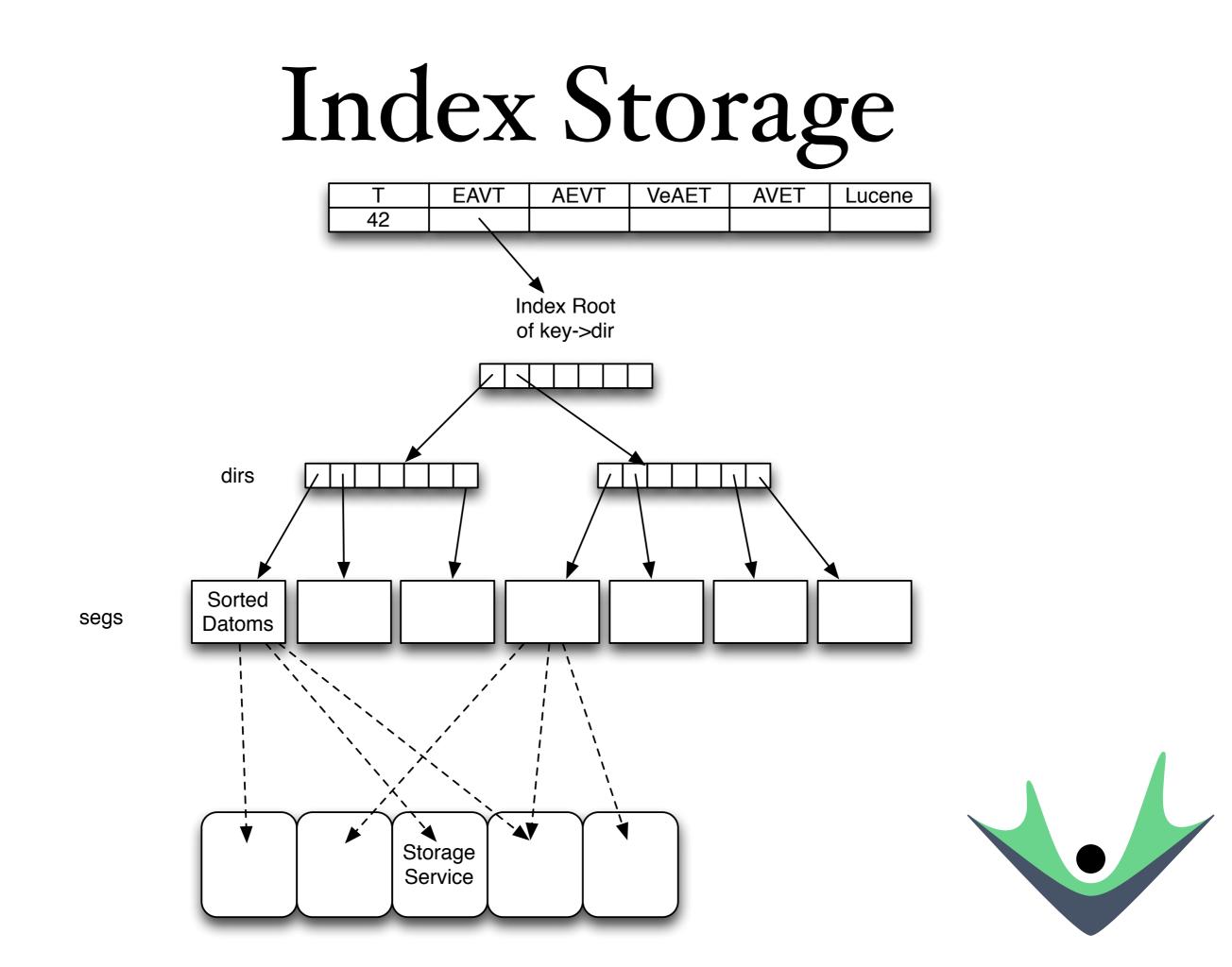
- New persistent sorted set
- Large internal nodes
- Pluggable comparators
- 2 sorts always maintained
 - EAVT, AEVT
- plus AVET, VAET



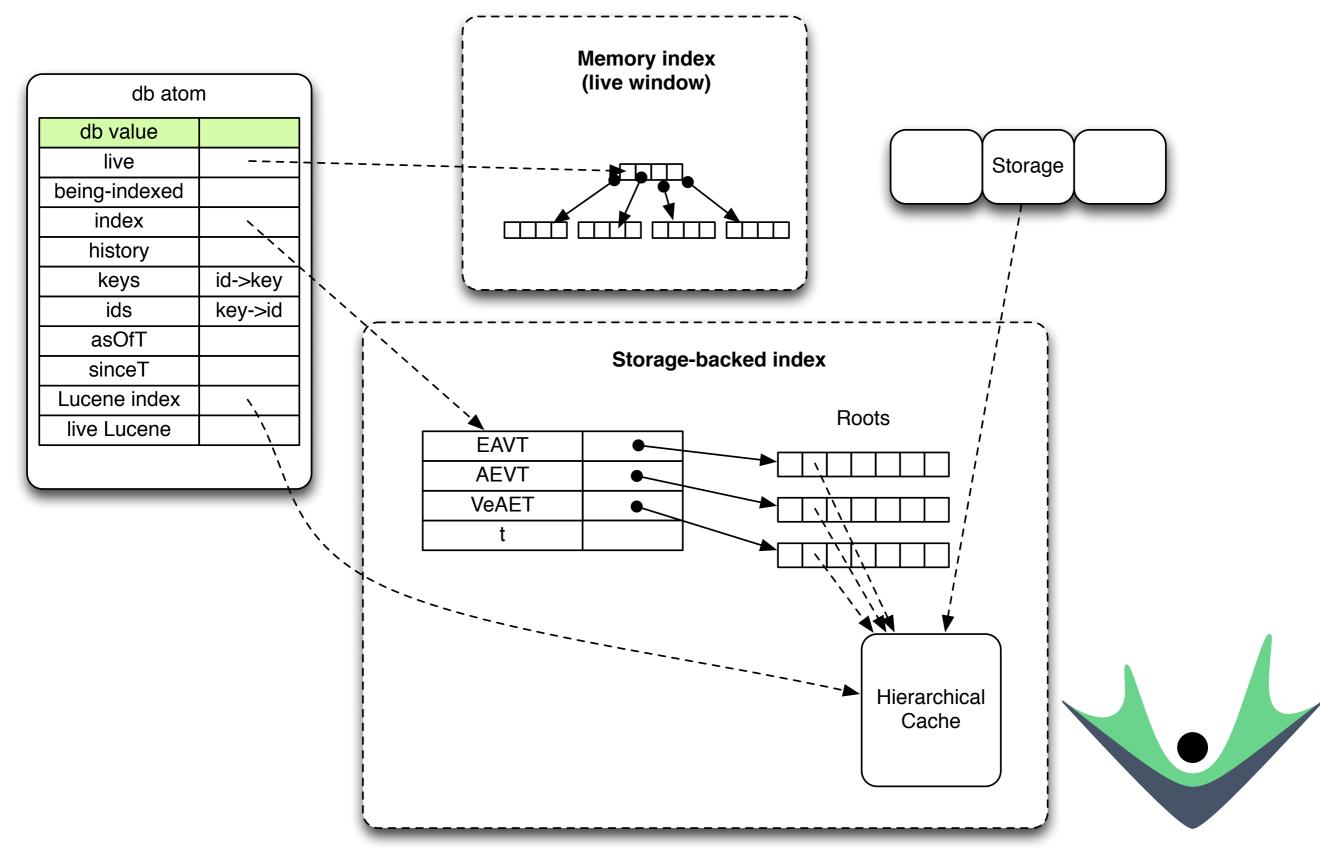
Storage

- Log of tx asserts/retracts (in tree)
- Various covering indexes (trees)
- Storage requirements
 - Data segment values (K->V)
 - atoms (consistent read)
 - pods (conditional put)





What's in a DB Value?

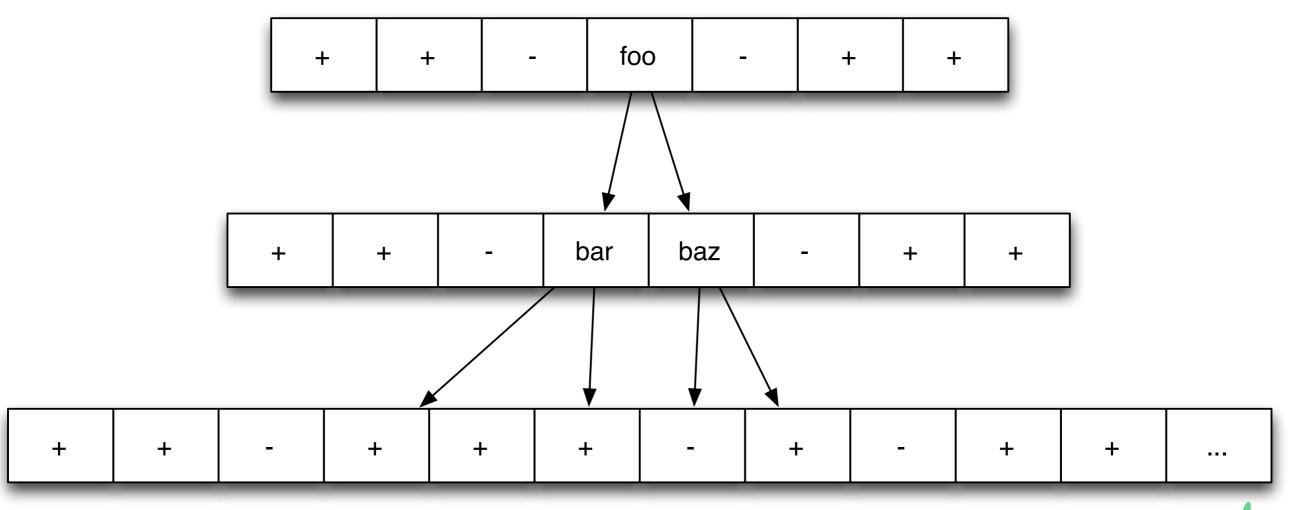


Process

- Assert/retract can't express transformation
- Transaction function:
 - (f db & args) -> tx-data
- tx-data: assert/retract/(tx-fn args...)
- Expand/splice until all assert/retracts



Process Expansion





Transactor

- Accepts transactions
 - Expands, applies, logs, broadcasts
- Periodic indexing, in background
- Indexing creates garbage
 - Storage GC



Transactor Implementation

- HornetQ for transaction communication
- Extensive internal pipelining j.u.c. queues
- Async message decompression
 - transaction expansion/application
 - encoding for, communication with storage
- Java interop to storage APIs

Indexing

- Extensive use of laziness
- Parallel processing
- Parallel I/O
- Async, rejoins via queue



Declarative Programming

- Embedded Datalog
- Takes data sources and rule sets as args
- Extended to work with scalars/collections
- Expression clauses call your code



Datalog Implementation

- Data driven, in and out
- Query/Subquery Recursive (QSQR)
 - Dynamic, set oriented
- DB joins leverage indexes
- Expressions use Clojure compiler
 - caching of transforms at all stages

Over Here

- Peers directly access storage service
- Have own query engine
- Have live mem index and merging
- Two-tier cache
 - Segments (on/off heap)
 - Datoms w/object values (on heap)

Peer Implementation

- HornetQ for transaction communication
- Google Guava caches
- Java APIs for storage
- Entities are like multimaps
 - key -> value(s)
 - reverse attrs



Consistency and Scale

- Process/writes go through transactor
 - traditional server scaling/availability
- Immutability supports consistent reads
 - without transactions
 - scale reads turning knobs on storage
- Query scales with peers
 - dynamic e.g. auto-scaling



Testing

- test.generative was born here
- Functional tests
- Simulation-based testing



Simplicity is Agility

- Key protocols extremely small (< 7 fns)
- Memory, embedded SQL, remote SQL, Infinispan, DynamoDB
- Move from our own dynamo cluster to DynamoDB:
 - 2 weeks
- Support PostgreSQL, Infinispan
 - 1 day each



Leverage

Read/print data Embedded language Runtime compilation Extend standard interfaces/protocols Interop ✓ State model - extended



Summary

- Clojure was made for this kind of app
- Fast enough at all levels
- Most key subsystems < 1000 lines
- A ton of concurrency, no sweat
- Leverage interop Hornetq, Guava etc
- Startup time could be better
- Datomic is Simple

