Netflix's Cloud Data Architecture



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2 Talks @ QCon London



- Whitepaper : Netflix's Transition to High-Availability Storage Systems
- CR Title : Netflix's Cloud Data Architecture
 - About Track : Architectures You've Always Wondered
 - General Overview
 - R Data Replication Deep Dive
- Raghtarrent Title : NoSQL @ Netflix
 - Track : NoSQL : Where and How







Circa 1997

Rent a Movie

Right of first sale



🛯 Buy a Movie

Any retailer (e.g. Walmart) or e-tailer (e.g. Amazon)





A brick-and-mortar store can only hold ~1k-2k titles

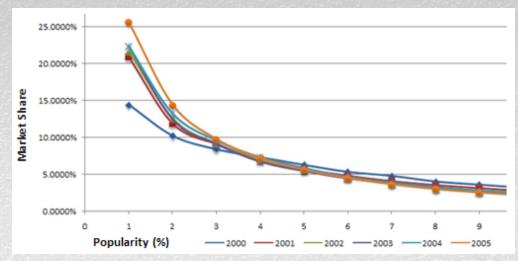
- R Which DVD titles do they pick?
 - R Brand New Releases
 - R Long-standing Hits (e.g. The Godfather)
- R Brand New Releases are expensive
- Stores profit by re-renting the same video within a short time frame (i.e. DVD Availability Window)

Must impose steep rental fees and steeper late fees



Circa 1997

- Retflix saw opportunity in the long-tail business
 - We store >120K titles in our 50+ shipping hubs
 - We recommend movies in the long tail, personalized to the customer, lowering costs







In 1999

- Retflix launches DVD rentals-by-mail
 - ↔ Unlimited rentals for a flat monthly fee
 - R No due dates
 - R No late fees







After a few years...







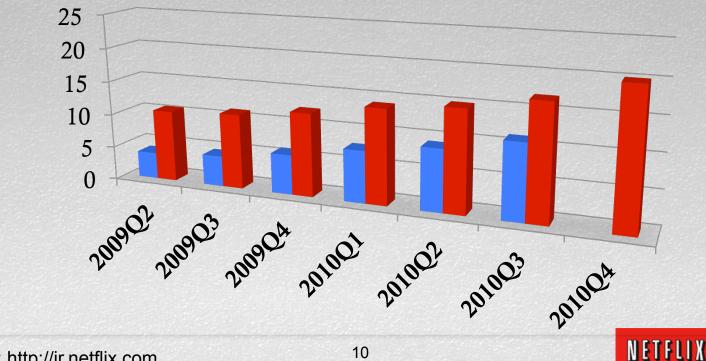
Fast Forward to 2008

- Netflix forecasts the gradual end of the DVD and starts switching to Movie Streaming
- - A We spend \$500MM to \$600MM annually on US Postage for DVD mailing
 - Streaming a movie is 1/100th the cost of shipping a DVD
 - Reasier to craft a business model for international expansion





- We have 20M+ subscribers in the US and Canada
- Uses 20% of US peak downstream bandwidth
- 1st or 2nd largest CDN user in the US





Grew subscribers 2008-2010 by partnering with device makers

- Game Consoles
 - R Xbox 360, Wii, PS3
- Ragina Mobile Devices

Apple IPad, IPhone, IPod Touch, Windows 7





Migration to AWS





Circa late 2008, Netflix had a single data center

- Approaching limits on cooling, power, space, traffic capacity
- Alternatives
 - R Build more data centers
 - Outsource our capacity planning and scale out
 - Allows us to focus on core competencies



- **Winner** : Outsource our capacity planning and scale out
 - CR Leverage a leading Infrastructure-as-a-service provider
 - Amazon Web Services
- **Footnote** : A short 2 years later, we serve >90% of our traffic out of AWS
 - excluding the video streaming bits, which come from CDNs



The Netflix Application Platform





AWS provides various IAAS offerings, but applications need more!

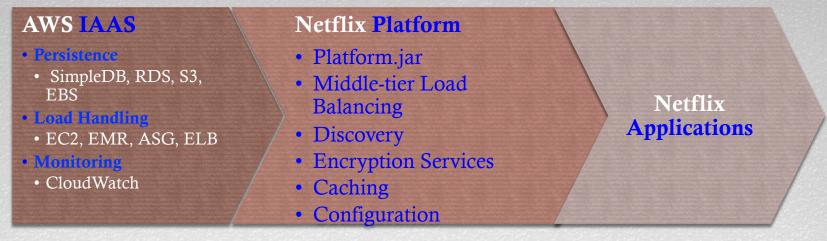




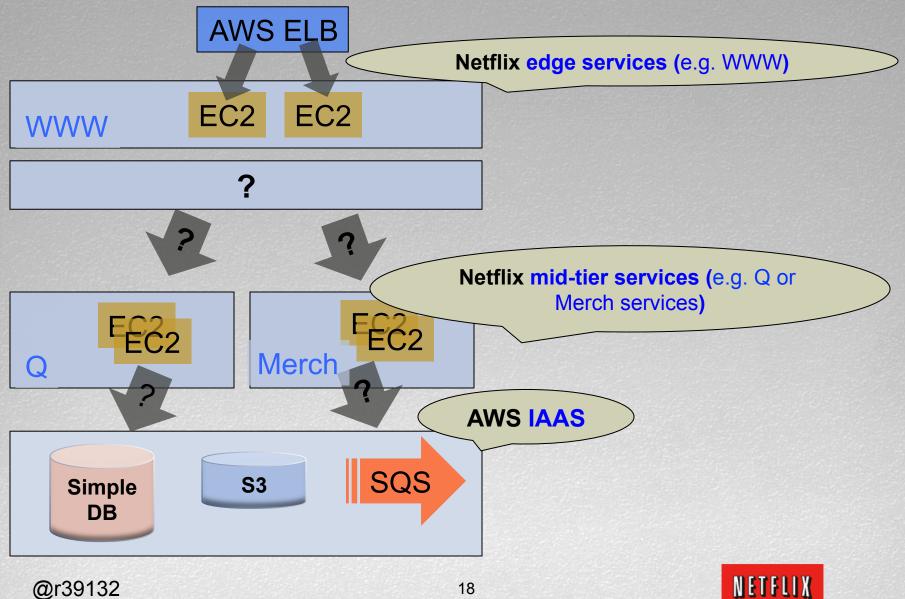


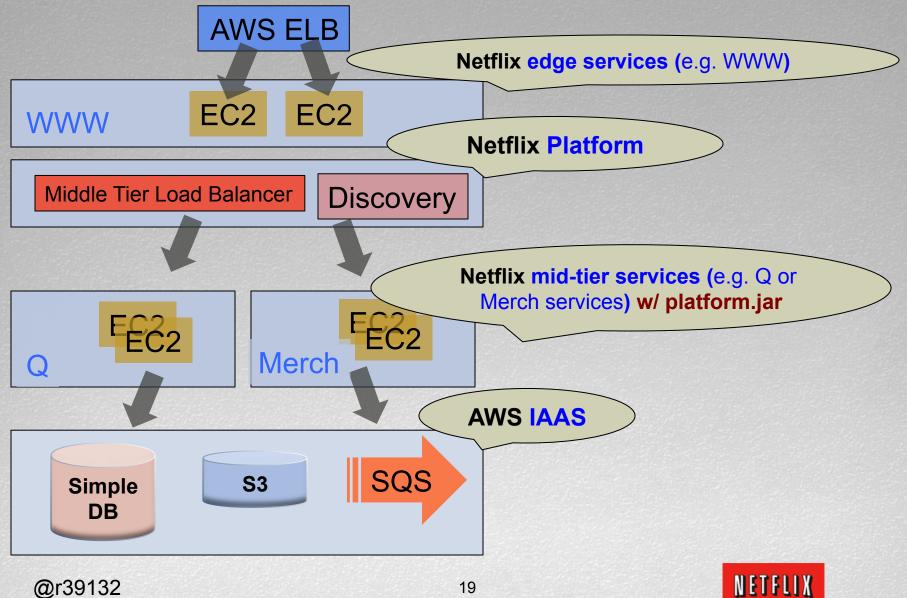
AWS provides various IAAS offerings, but applications need more!

Hence the need for Netflix's infrastructure team to bridge the gap!

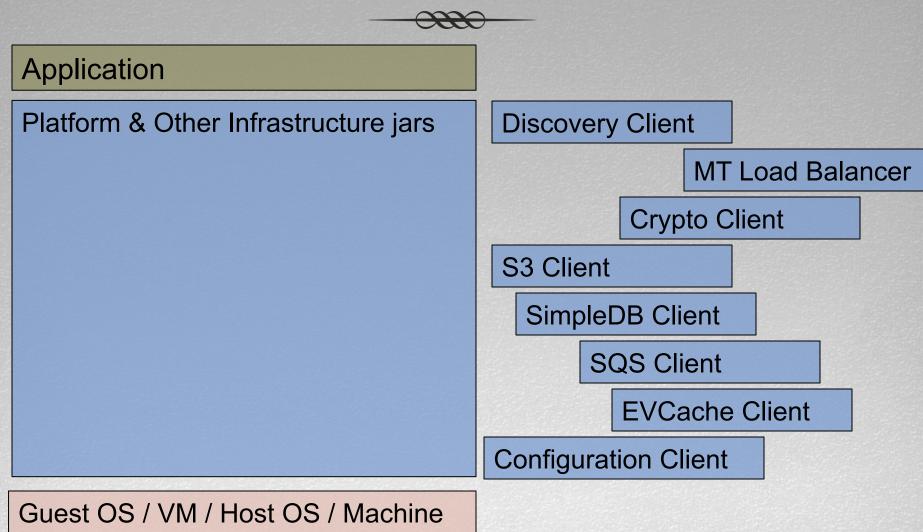








How Netflix Uses AWS (EC2 Stack)





The Netflix Data Platform





- Representation of the Cloud (c.f. NoSQL @ Netflix talk)

 - cr S3
 - R Cassandra
- 🛯 Data Replication
 - R (a.k.a. Item Replicator)
 - \bigcirc HAProxy + Squid + Oracle \rightarrow temporary

Persistence



SimpleDB



Persistence : SimpleDB



Terminology

SimpleDB	Hash Table	Relational Databases
Domain	Hash Table	Table
Item	Entry	Row
Item Name	Key	Mandatory Primary Key
Attribute	Part of the Entry Value	Column



Persistence : SimpleDB



Soccer Players				
Кеу	Value			
ab12ocs12v9 b24h3b3403b	First Name = Harold First Name = Pavel	Last Name = Kewell Last Name = Nedved	Nickname = Wizard of Oz Nickname = Czech Cannon	Teams = Leeds United, Liverpool, Galatasaray Teams = Lazio, Juventus
cc89c9dc892	First Name = Cristiano	Last Name = Ronaldo		Teams = Sporting, Manchester United, Real Madrid

SimpleDB's salient characteristics

- SimpleDB offers a range of consistency options
- SimpleDB domains are sparse and schema-less
- The Key and all Attributes are indexed
- Each item must have a unique Key
- An item contains a set of Attributes
 - Each Attribute has a name
 - Each Attribute has a set of values
 - All data is stored as UTF-8 character strings (i.e. no support for types such as numbers or dates)





Persistence : SimpleDB



- Moved some of our largest transactional data sets to SimpleDB in 2010
 - e.g. **RENTAL HISTORY** : Everything that anyone has watched at Netflix, including streaming and DVD
- We have
 - \sim thousands of domains
 - ~ 1TB of OLTP/transactional data (i.e. no CLOBS or BLOBS)
 - \sim > billions of rows of data (a.k.a. items)
- We execute billions of SimpleDB statements per day



Persistence







Persistence : S3



- Simple KV store organized as objects in buckets
 - Reach AWS account is given a max of 10 buckets
 - Reach bucket can hold an unlimited number of objects
- ↔ We use S3 to store data that does not fit into SimpleDB
 - R Logs from streaming devices
 - R Files used in movie encoding
 - R Truncated-tail of Rental History
 - Good pattern introduced by Greg Kim

Data Replication



Item Replicator

Data Replication between Oracle, SimpleDB, S3, and Cassandra



Data Replication : IR



- Replication Home-grown Data Replication Framework known as IR for Item
- ≪ Keeps data in sync between RDBMS (DC) and NoSQL (Cloud)
 ∞ Mostly unidirectional: DC → Cloud
- - R Trigger-oriented IR
 - All CRUD operations on source table X are copied via a trigger to XLog_i journal tables (i is the shard index)
 - R reads (polls) from a journal table
 - - R reads (polls) from source table X directly



Data Capture Schemes



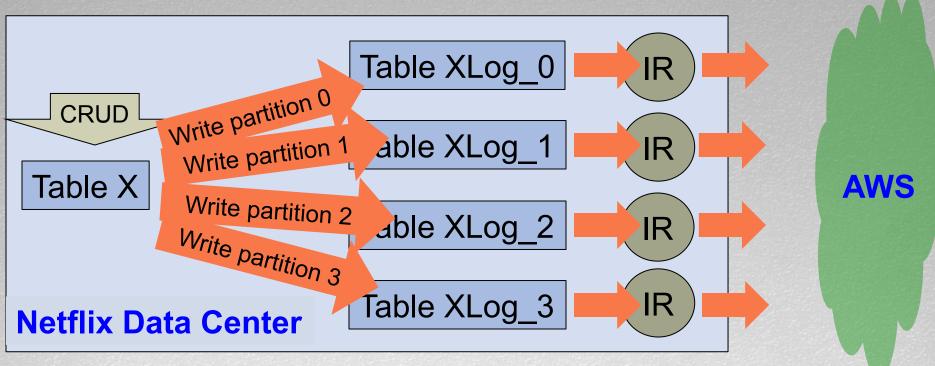




Data Capture Schemes : IR



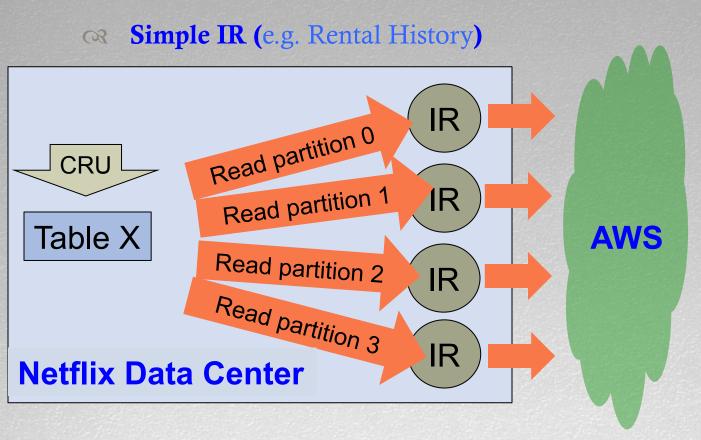
R Trigger-oriented IR (e.g. Movie Q)





Data Capture Schemes : IR







IR's Polling Select





IR's Polling Select



Anatomy of an IR Select Query

- Reserve the SQL below as a recurring poll select * from RENTAL_HISTORY r where r.LAST_MODIFIED_TS > checkpoint value ← (A) and r.LAST_MODIFIED_TS < (now-10 seconds) ← (B) order by r.LAST_MODIFIED_TS asc ← (C)
- (A) preserves progress in case IR crashes and restarts
- (B) handles interesting race condition... will explain in later slide
- (\mathbf{C}) in order read for repeatability as in (A)

IR's Polling Select



Anatomy of an IR Select Query

- Assume previous checkpoint is T0
- \bigcirc (a) time = T4, only the top 2 records are visible and will be replicated
 - (B) hides The Town and Duma
 - (R) {1, The Machinist} and then {2, Blood Diamond} are replicated to the cloud
- \curvearrowright After IR replicates this data, the checkpoint is now T2

Customer_ID	Movie_ID	last_modified_ts	
1	The Machinist	T1	
2	Blood Diamond	T2	
4	The Town	T4 - 3sec	
1	Duma	T4	

NETFLIX



Anatomy of an IR Select Query

- Assume previous checkpoint is T2
- a time = T4+12 sec, The 3 records pointed to by arrows are visible and replicated to the cloud
 - Of these 1 new record, The King's Speech, became visible in time to be replicated Commit-delayed

	Customer_ID	Movie_ID	last_modified_ts	
	1	The Machinist	T1	
	2	Blood Diamond	T2	
	4	The Town	T4 – 3 sec	\leftarrow
new	3	The King's Speech	T4	\leftarrow
	1	Duma	T4	\leftarrow
new	5	Rescue Dawn	T4 + 11 sec	



- Consider 2 update transactions

update RENTAL_HISTORY r
set r.LAST_MODIFIED_TS = systimestamp,
r.MOVIE_ID = 'The Machinest good'
where r.CUSTOMER_ID = 1;
commit;

R Transaction 2

update RENTAL_HISTORY r set r.LAST_MODIFIED_TS = systimestamp, r.MOVIE_ID = 'Blood Diamond good' where r.CUSTOMER_ID = 2; commit;





- - R Transaction 1

update RENTAL_HISTORY r set r.LAST_MODIFIED_TS = systimestamp **@T1** r.MOVIE_ID = 'The Machinest good' where r.CUSTOMER_ID = 1; commit; **@T4**

R Transaction 2

update RENTAL_HISTORY r
set r.LAST_MODIFIED_TS = systimestamp, @T2
r.MOVIE_ID = 'Blood Diamond good'
where r.CUSTOMER_ID = 2;
commit; @T3

Customer_ID	Movie_ID	last_modified_ts				
1	The Machinist bad	Т0				
2	Blood Diamond bad	Т0				
Transaction 2 commits at T3, but records T2						
Customer_ID	Movie_ID	last_modified_ts				
1	The Machinist bad	Т0				
2	Dlagd Diamond agod	T 2				
2	Blood Diamond good	T2				

- -- IR replicates "Blood Diamond good" and sets checkpoint = T2
- -- Transaction 1 commits at T4, but records T1

Customer_ID	Movie_ID	last_modified_ts
1	The Machinist good	T1
2	Blood Diamond good	T2

-- IR will never see "The Machinist good" because checkpoint already advanced past T1 to T2





Solving the "Commit-delayed" Race Condition

R Transaction 1

update RENTAL_HISTORY r

set r.LAST_MODIFIED_TS = systimestamp,

r.MOVIE_ID = 'The Machinest good'

where $r.CUSTOMER_ID = 1$

and r.last_modified_ts < (now - 10 sec);

commit;

```
R Transaction 2
```

update RENTAL_History r
set r.LAST_MODIFIED_TS = systimestamp,
r.MOVIE_ID = 'Blood Diamond good'
where r.CUSTOMER_ID = 2;
and r.last_modified_ts < (now - 10 sec);
commit;</pre>

Forklifting Historical Data



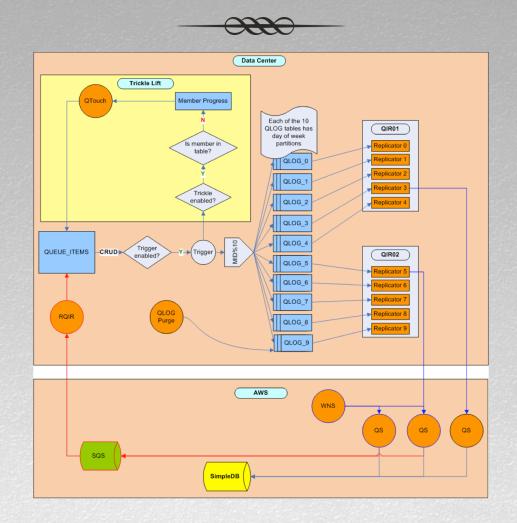


Forklifting Historical Data



- R Trigger-oriented IR
 - R Trickle-lifting
 - R Parallel fork-lifting and incremental replication
- - Typically need to forklift by recovering a snapshot in a second database
 - R Execute forklift of that data
 - \bigcirc Assume database snapshot @ time = T1
 - In the primary database, buffer modifications until forklift complete
 - One forklift is complete, replicate data changed after time=T1

Forklifting Historical Data



Forklifting Historical Data



- - \bigcirc User with id=6 adds a movie to his queue
 - Trigger fires to replicate that data to $QLOG_6$ (e.g. 6 % 10 shards = 6)
 - Along the way, the trigger checks whether user with id=6 is found in the member_progress table
 - If not, the user id is inserted into the member_progress table
 - R If so, no action is taken
 - A separate program called QTouch will notice a new record while polling the member_progress table : user id = 6
 - QTouch will execute a harmless update of all records for the user to fork-lift that user's data into the cloud



Best Practices : IR



- - R Checkpoint Column
 - Choose a Timestamp not Date data type (i.e. for microsecond granularity) or an ordered sequence
 - Need a Before Trigger to set the timestamp or sequence
 Don't trust what is passed in by DB clients

Data Replication



HAProxy + Squid + Oracle

Data lives in Oracle but is cached in the cloud



Data Replication : HAProxy+Squid+Oracle

