

# QCon 2012

Progressive Architectures at the Royal Bank of Scotland

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Ben Stopford

# Overview

History – Front-Office-Centric

‘The City’ has been the source of industry innovation

FPML

AMQP

OpenAdaptor

LMAX Disruptor



# New Drivers

Change: Even more than before!:

- Regulation
- Competition – The ‘Race to Zero’
- Utilisation – data-centres, ‘Just Enough’ hardware

All in the face of even tighter budgets

However – there is hope!: New technologies and architectures

We will show you some of the challenges and our approaches



**Necessity has become an even greater Mother of Innovation**

# Agenda

Ben: Pauseless Java in Low Latency Trading  
Beyond Messaging with ODC

Fuzz: Risk: Increased Accuracy - Big Compute and Big Data

Mark: Data Virtualisation  
Collaboration

Questions



# Development Areas

UX

Data & Resource Consolidation

GPGPUs, FPGA

Big Data (for us)

Data Visualisation

Data Virtualisation

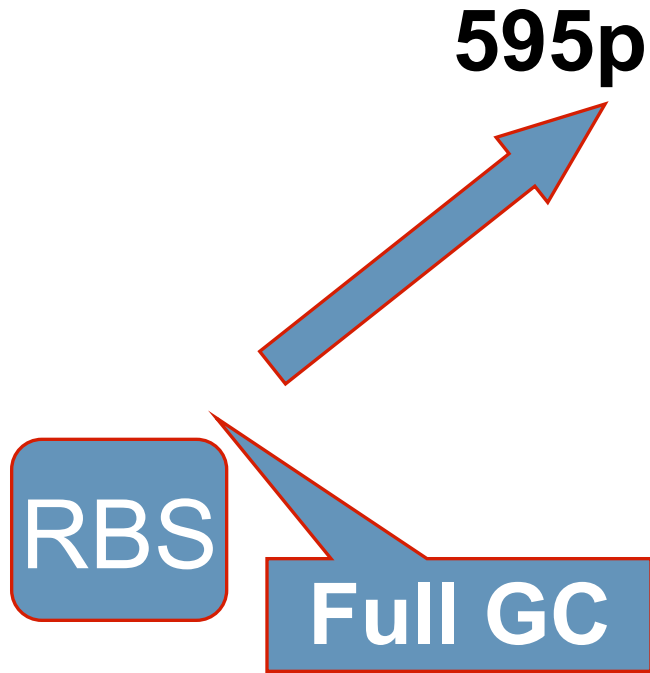


We're driving the shape and features of products – particularly in the last two

# The Manhattan Processor

Avoiding GC Pauses in High Frequency Trading

# The Problem:



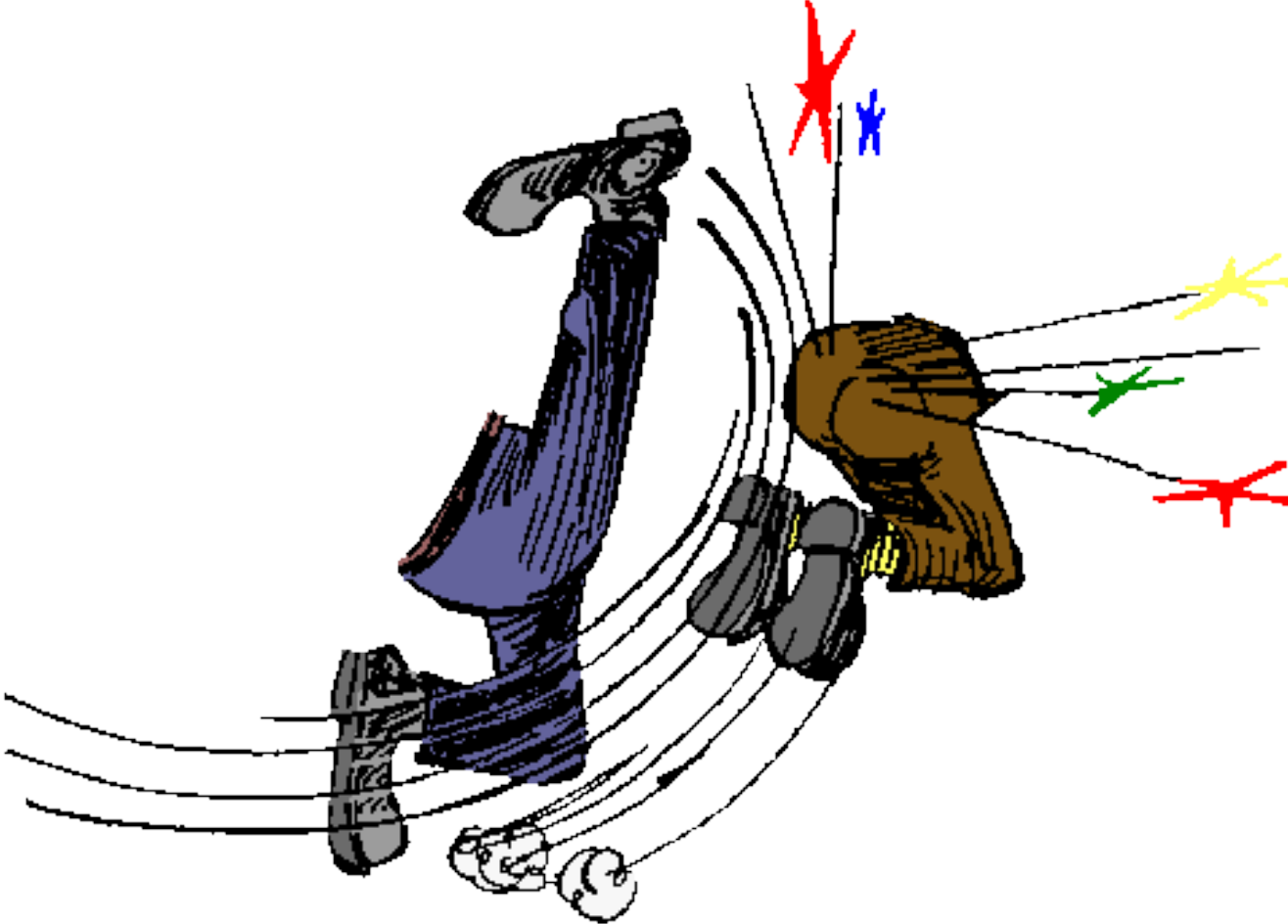
**615p**

**Apple announces  
Financial Results  
above expectations**



**Buy @ 595p  
Short @ 615p**

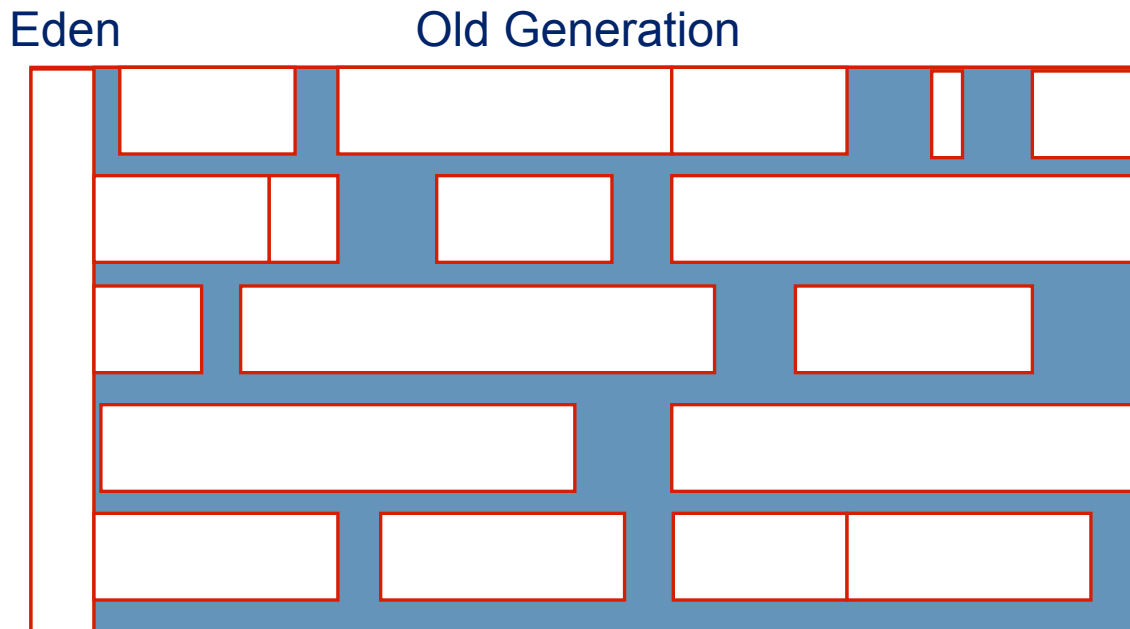
Angry Trader kicks techies collective behinds for losing one months P&L in a one second GC!





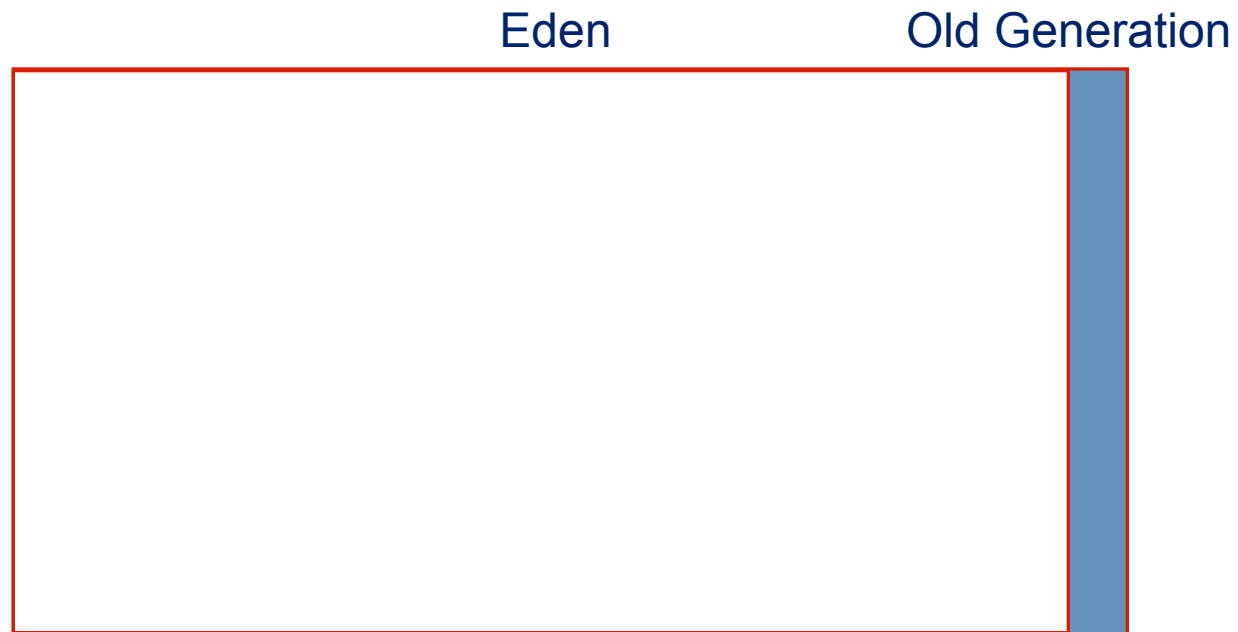
## So can we just tune the collector?

- Cost ~ 1 sec per GB of heap at best
- Stop-the-World happens more than you may think, even with CMS
  - Eden collections
  - Mark/Remark in CMS
  - Full GC (really hurts)
  - CMS falls back to Full GC on ‘Concurrent Mode Failure’



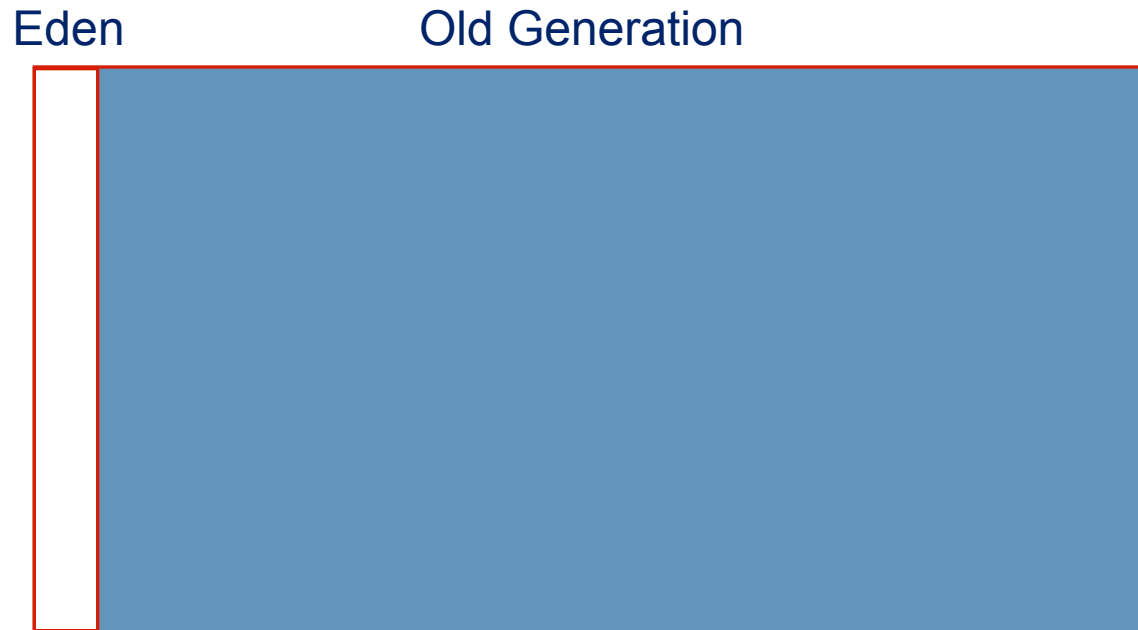
## Avoiding Pauses: Model 1

- Use a huge Eden space (several GB' s)
- Create minimal Garbage



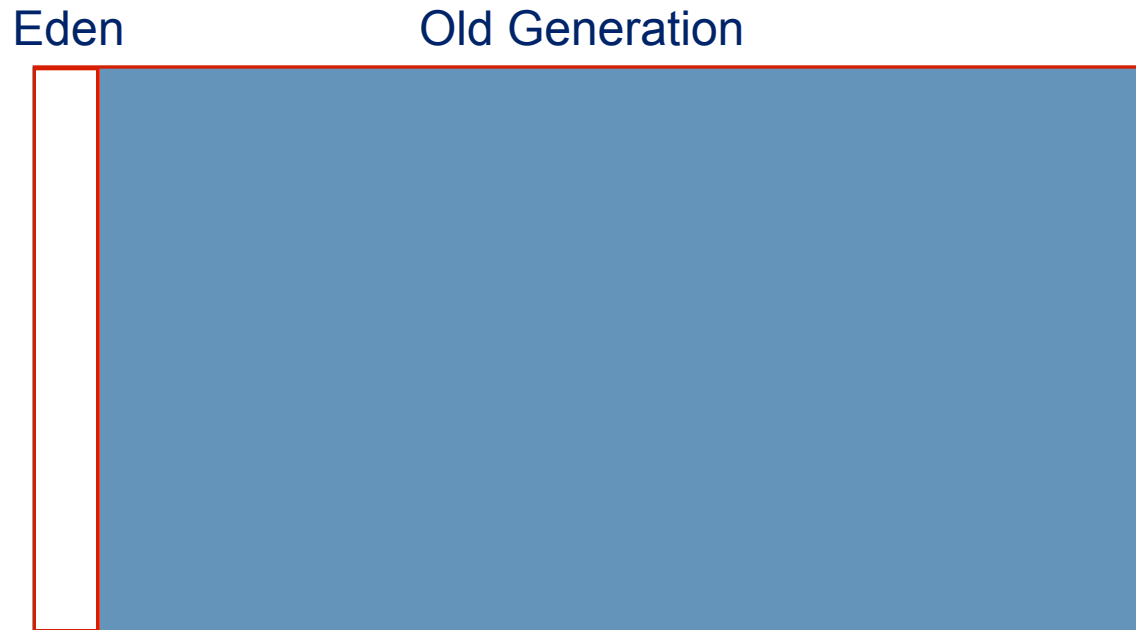
## Avoiding Pauses: Model 2

Use a small Eden that can be collected in a known time (~ 1ms typically) and avoid FullGC



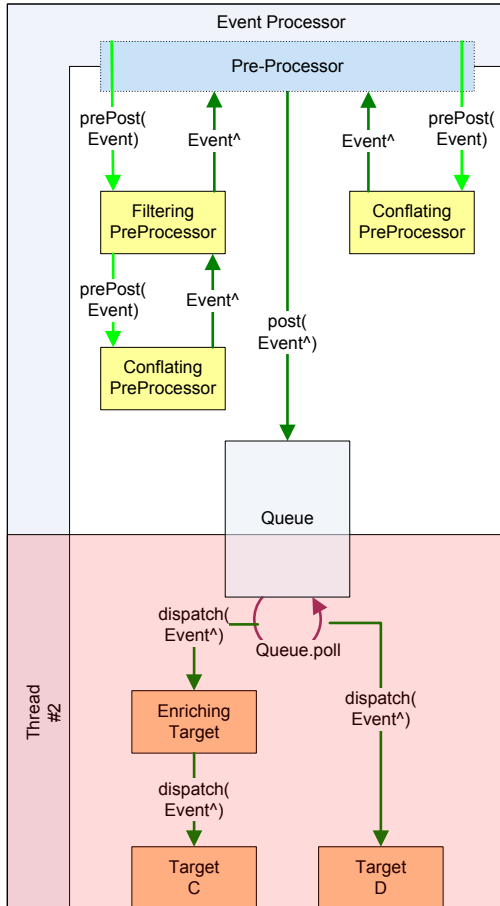
## We use Model 2

- No control over 3<sup>rd</sup> party
- It's very hard to eliminate the production of garbage completely
- We can tolerate very small collections

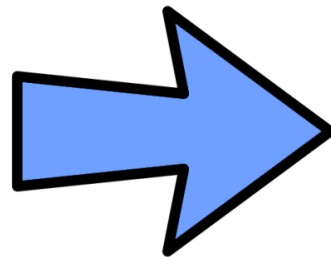


Manhattan Processor:  
Concurrent processing with a  
very low garbage footprint.

# Manhattan Processor



- In-house non-allocating queue specialised for multiple producer, single consumer requirement
- Non-blocking for offer and poll
- Non-blocking object pool
- Pluggable & chainable pre- and post-processors
- Direct in-thread dispatch where there is thread affinity between producer and consumer



Minimal Garbage

## Result

- Fast, pluggable architecture for concurrent processing
- Predictable pauses <1ms
- No full collection pauses in 15hr trading day

Risk: Accuracy → Big Compute and Data

Market and Credit Risk Technology



## What's the Big Deal?

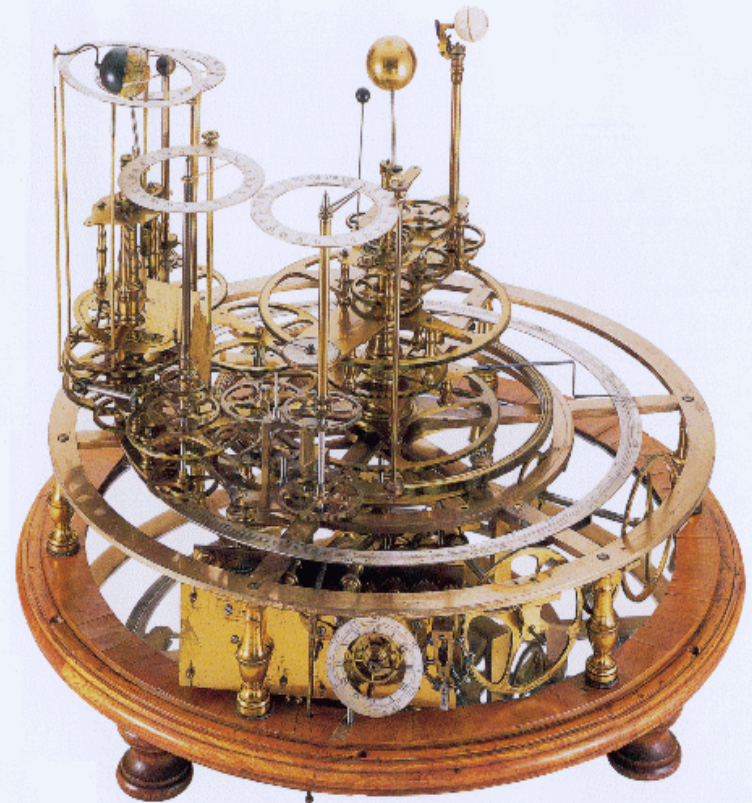
- Risk: Predicting the Future, looking for worst case scenarios
- Measuring Risk is not simple
  - Not as difficult as predicting the weather, but not far off
- Doing it right implies a lot of data and compute
- Impact on architecture and technology choices

Disclaimer:

**The numbers quoted in this presentation are indicative of magnitude only.**

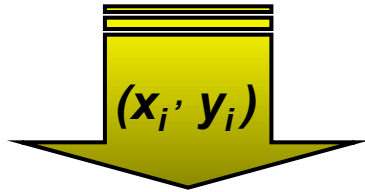
## Historically, Banks have....

- Linear approximation
- Single-step Monte Carlo
- Multi-step Monte Carlo

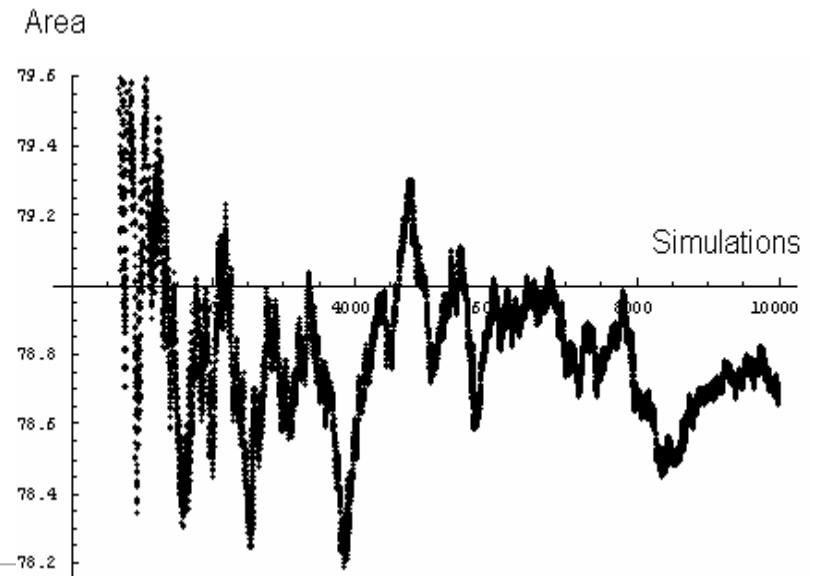


The Market is a complex system and needs to be modelled as such

# Monte-Carlo: Example



$$\text{Hit}_i = \text{InsideCircle}(x_i, y_i)$$



# Single Step vs Multi Step Monte-Carlo

Get directions My places

Car Bus Walk

A Highbury Corner, London

B Queen Elizabeth II Conference Centre

Add Destination - Show options

Depart at 08/03/12 08:00

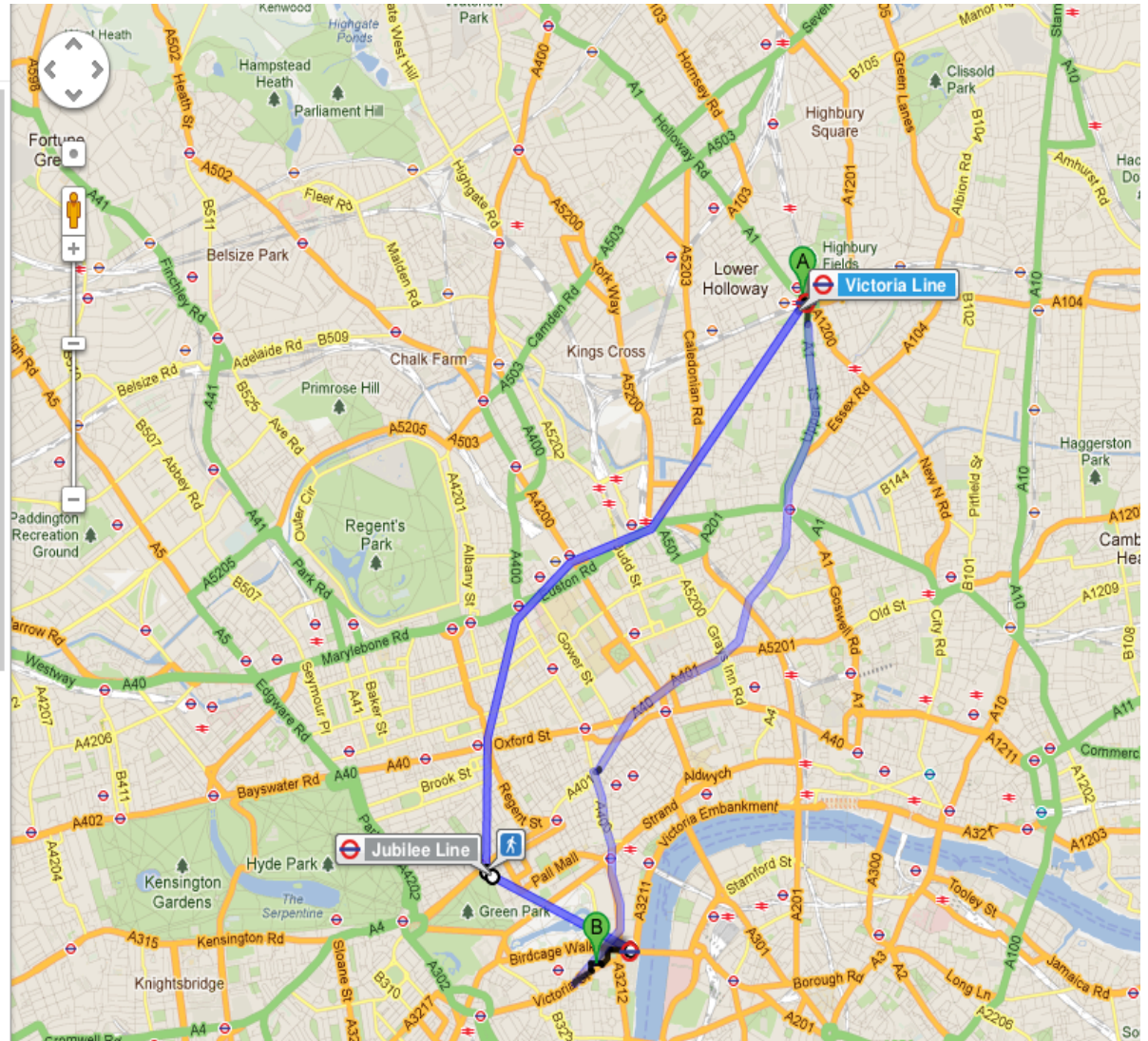
GET DIRECTIONS

## Suggested routes

- 27 mins  
08:01 - 08:28
  - 27 mins  
08:03 - 08:30
  - 27 mins  
08:05 - 08:33
  - 52 mins  
08:02 - 08:54
- Or Drive/Taxi 15 mins

## Public transport directions to Queen Elizabeth II Conference Centre

A Highbury Corner  
London N5 1RD, UK



# Monte-Carlo of my route to QCon

- 10s of relevant routes to Westminster
- 100s of events that affect the journey (weather, traffic jams, dropping off the kids to school ...)
- Each of these is correlated with other events
  - e.g. Traffic and Weather
- Events can happen at many discrete points in time
- Need 1000' s of random numbers to simulate this

# Monte-Carlo: Data Volumes

## Single Step Simulation

- 1' 000 paths
- 1 correlated random number per path
- ~50 observables
- Updated quarterly

## Multi Step Simulation

- 10' 000 paths
- 25' 000 correlated random numbers per path
- 1000' s observables
- Updated monthly



Multi-Step requires significantly more Data

# How much Compute?

## Single Step Simulation

< 4,000 compute hours

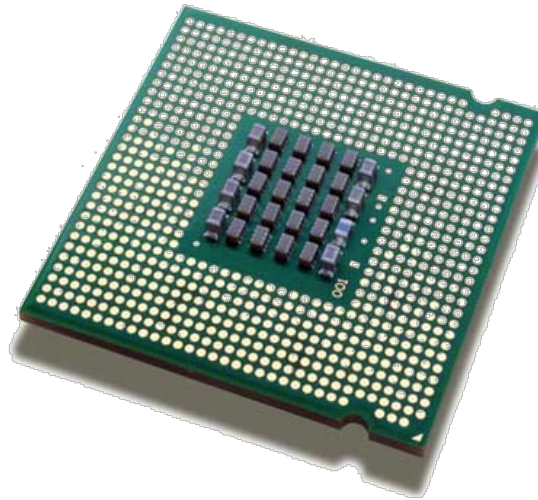
< 1,200 CPU cores



## Multi Step Simulation

~ 20,000 compute hours

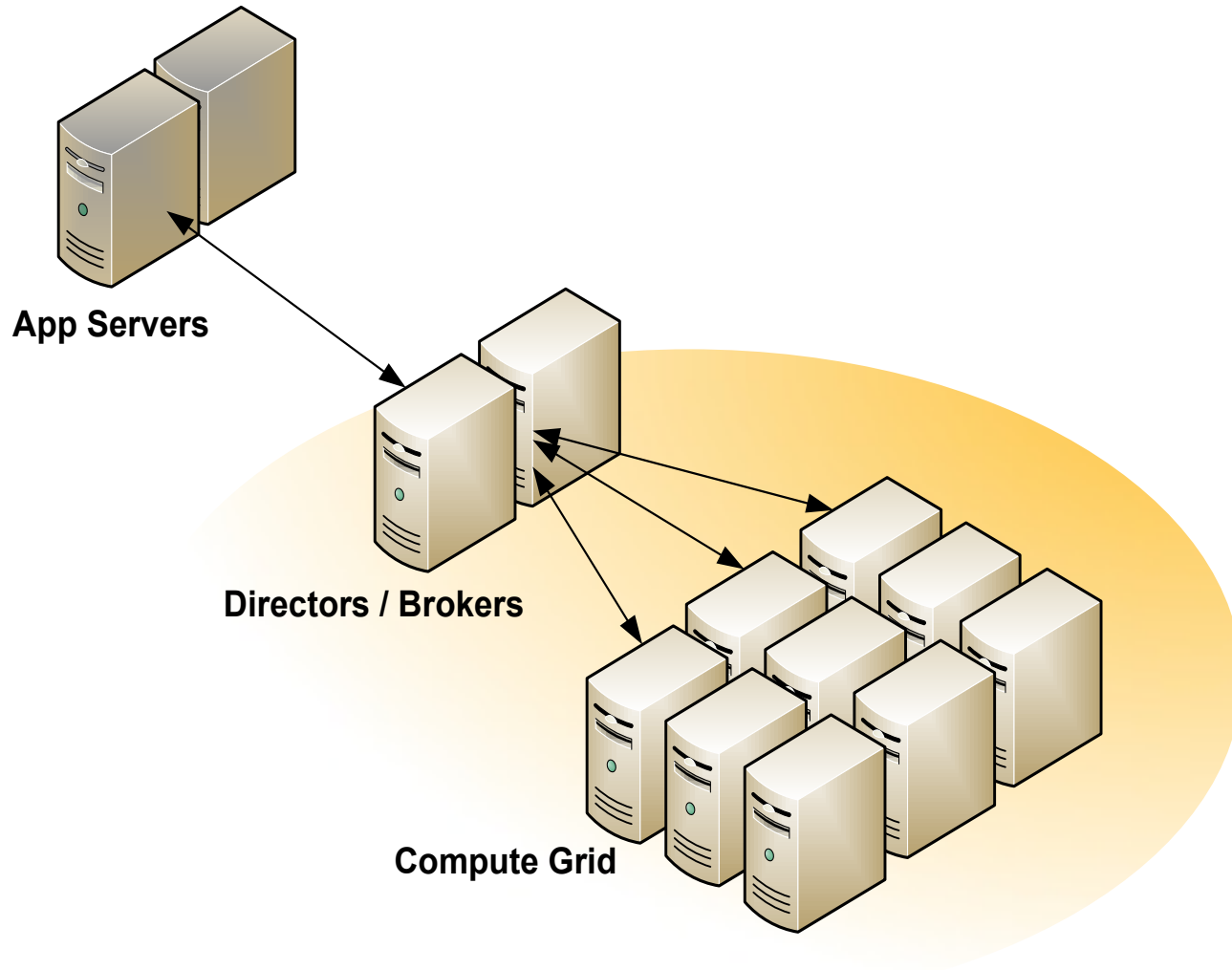
~ 5,000 CPU cores



Massive growth in Computations



# How do we handle that?



**Standard Compute Grid Architecture**

# Huge difference in required Data Volumes

Single Step Simulation

Multi Step Simulation

~ 2 MB

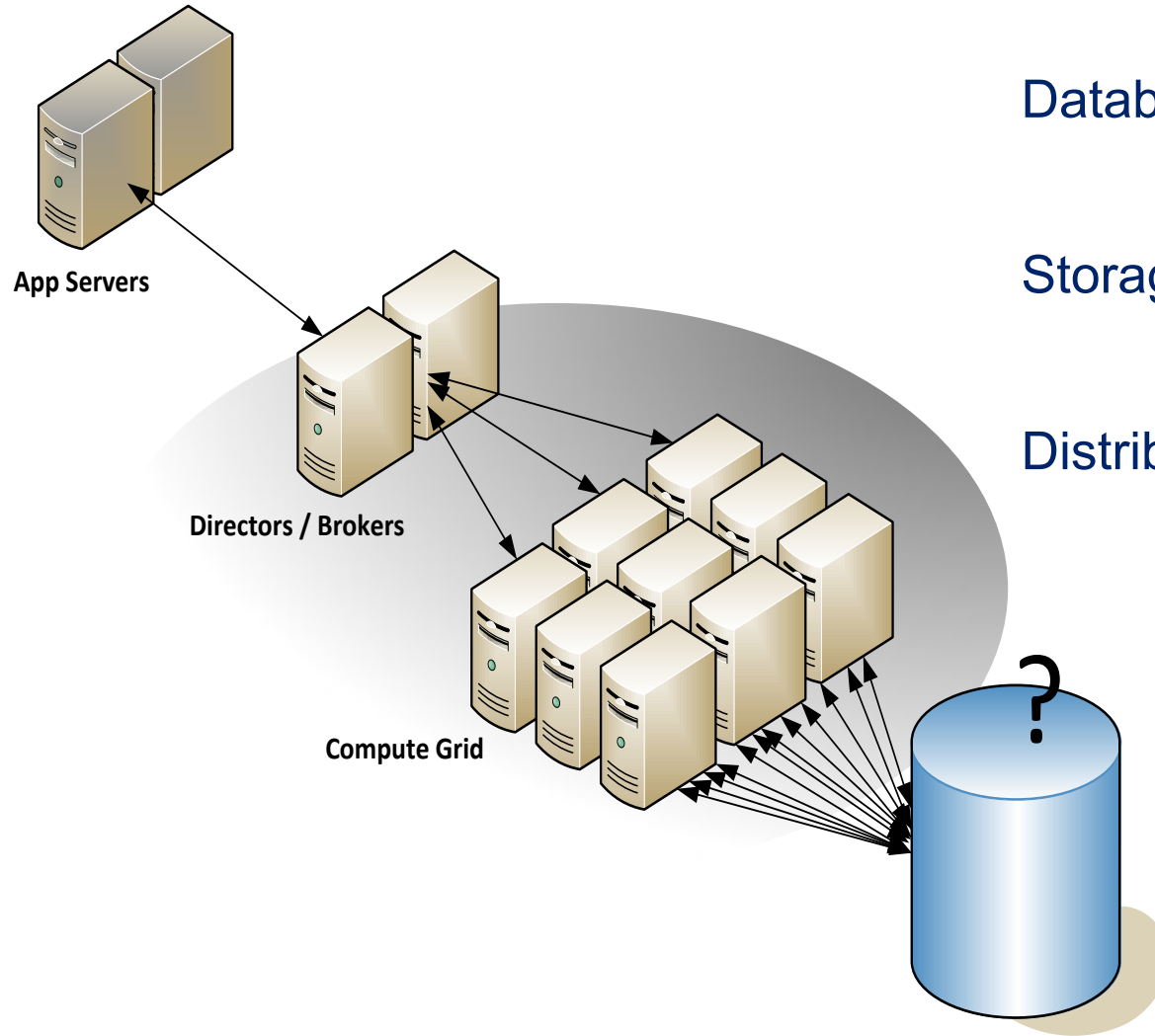


~ 10 TB



How will this data be accessed?

# Storage .....



Database?

Storage Appliance?

Distributed Filesystem?

What's the right solution?

## Requirements

Speed	Concurrent Access Bandwidth Transaction Throughput
Scalability	Concurrent Access Rack awareness
Robustness	Resilience and Recovery Guaranteed Correctness
Interoperability	Access from both Linux and Windows Server Grid Engines.
Support Infrastructure	Logging and Performance Monitoring Support Infrastructure

# Distributed File Systems



## Selection of Top Users



**1-100 nodes, many clusters**



**532 nodes – 5.3 PB storage**

The Facebook logo, which is a dark blue rectangular box with the word 'facebook' written in white, lowercase, sans-serif font. A registered trademark symbol is at the end of the word.

**facebook**

**1100 nodes, 8800 cores, 12 PB raw storage**



**120, 580, 1200 nodes**

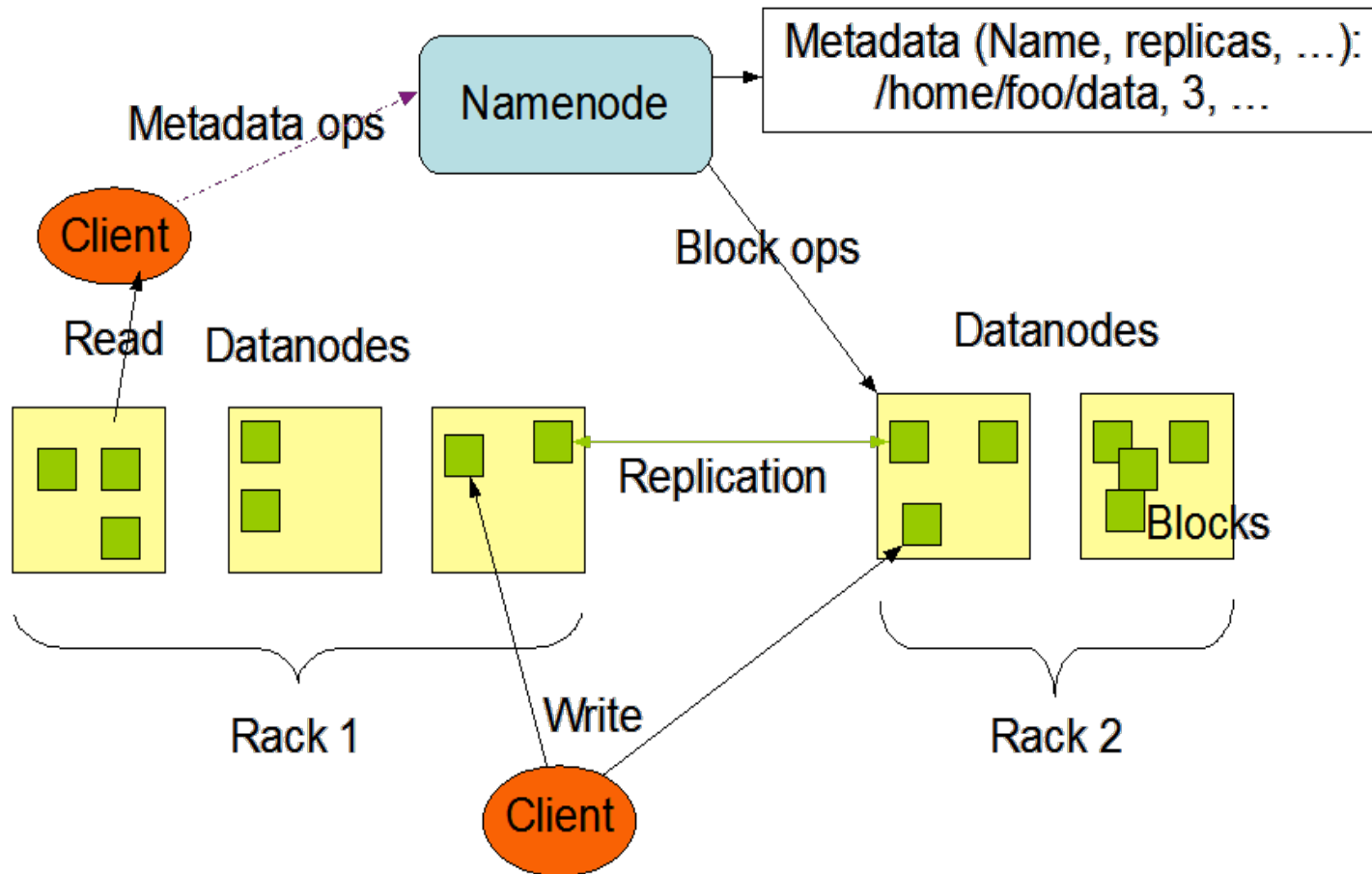


**40,000 nodes**

**Hadoop scales and is mainstream**

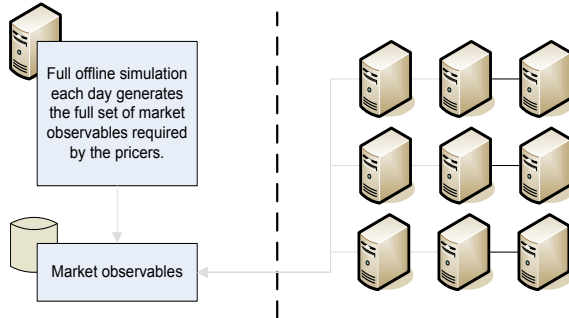
# Quick dip into HDFS

## HDFS Architecture



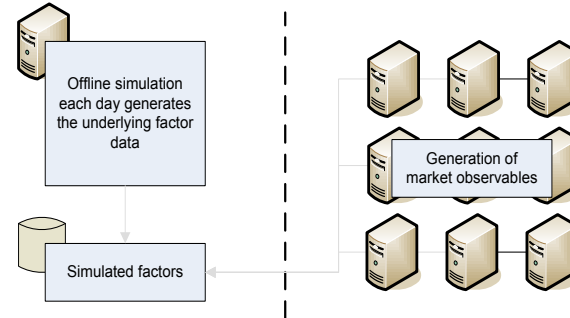
# We tried several architectures

Option 1



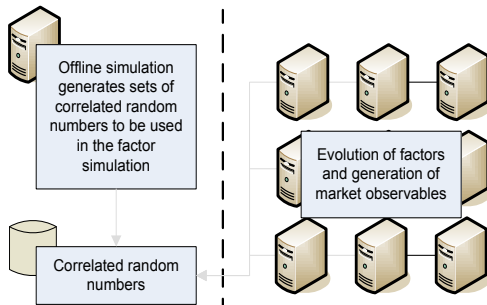
Distributed calculation nodes access the market observables directly on demand. The storage and remote access technology is not defined, but the following could be used:  
Disk based storage with access over HTTP  
iRODS server with storage in HDF5

Option 2



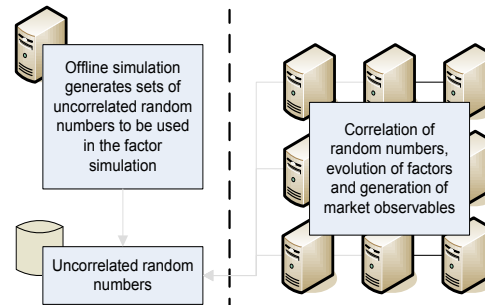
Distributed calculation nodes access the underlying simulated factors and then generate the market observables as required. The storage and remote access technology is not defined, but the nature of the simulated factors may lend itself to a technology that supports some type of sub-element access.

Option 3a



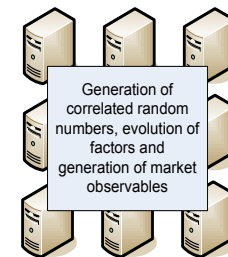
Distributed calculation nodes access the correlated random numbers on demand. The factors are then evolved and the market observables generated. The storage and remote access technology is not defined, but the nature of the correlated random numbers may lend itself to a technology that supports some type of sub-element access.

Option 3b



Distributed calculation nodes access the correlated random numbers on demand. The factors are then evolved and the market observables generated. The storage and remote access technology is not defined, but the nature of the correlated random numbers may lend itself to a technology that supports some type of sub-element access.

Option 4

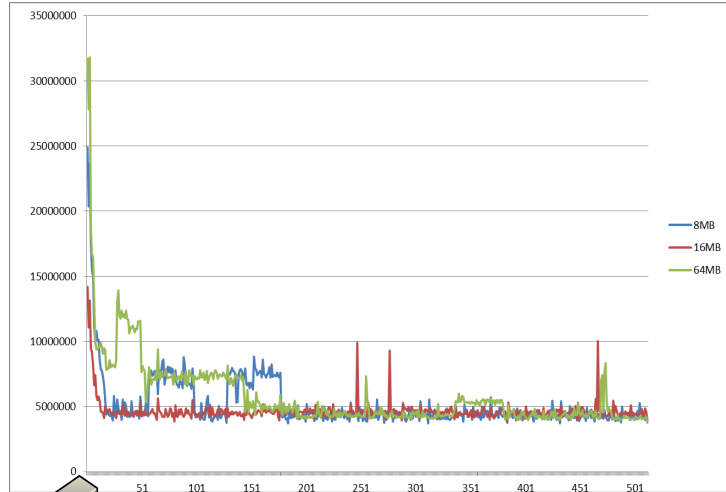


Distributed calculation nodes generate the correlated random numbers, the factor evolution and the market observables as required.



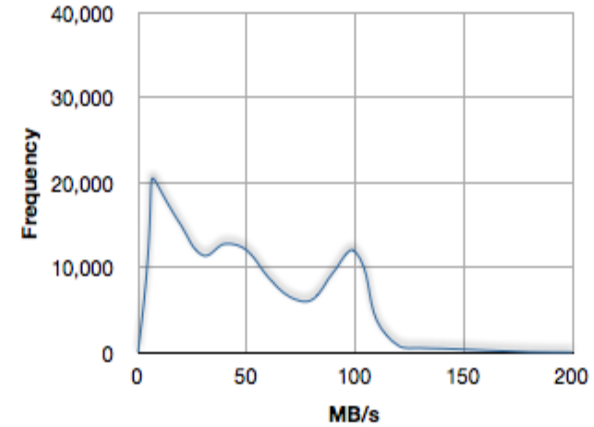
# And carried out a lot of analysis ...

Transfer speeds as a function of HDFS block size

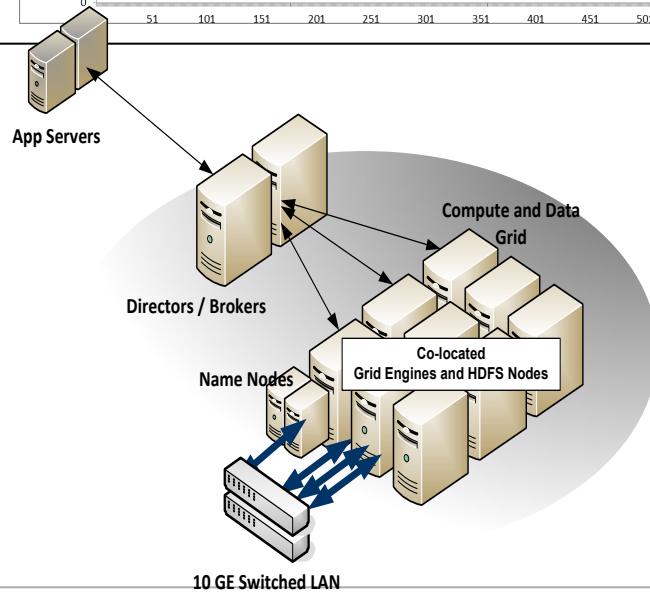
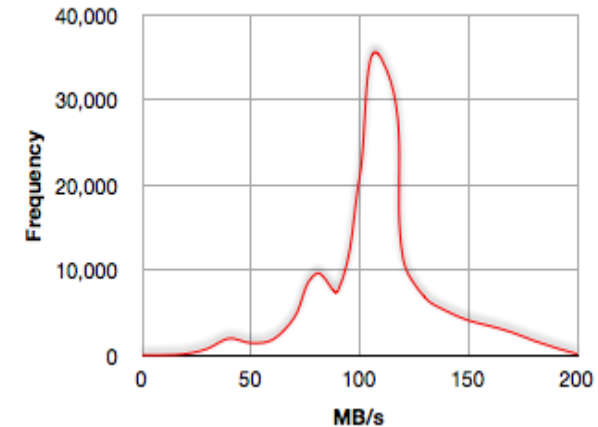


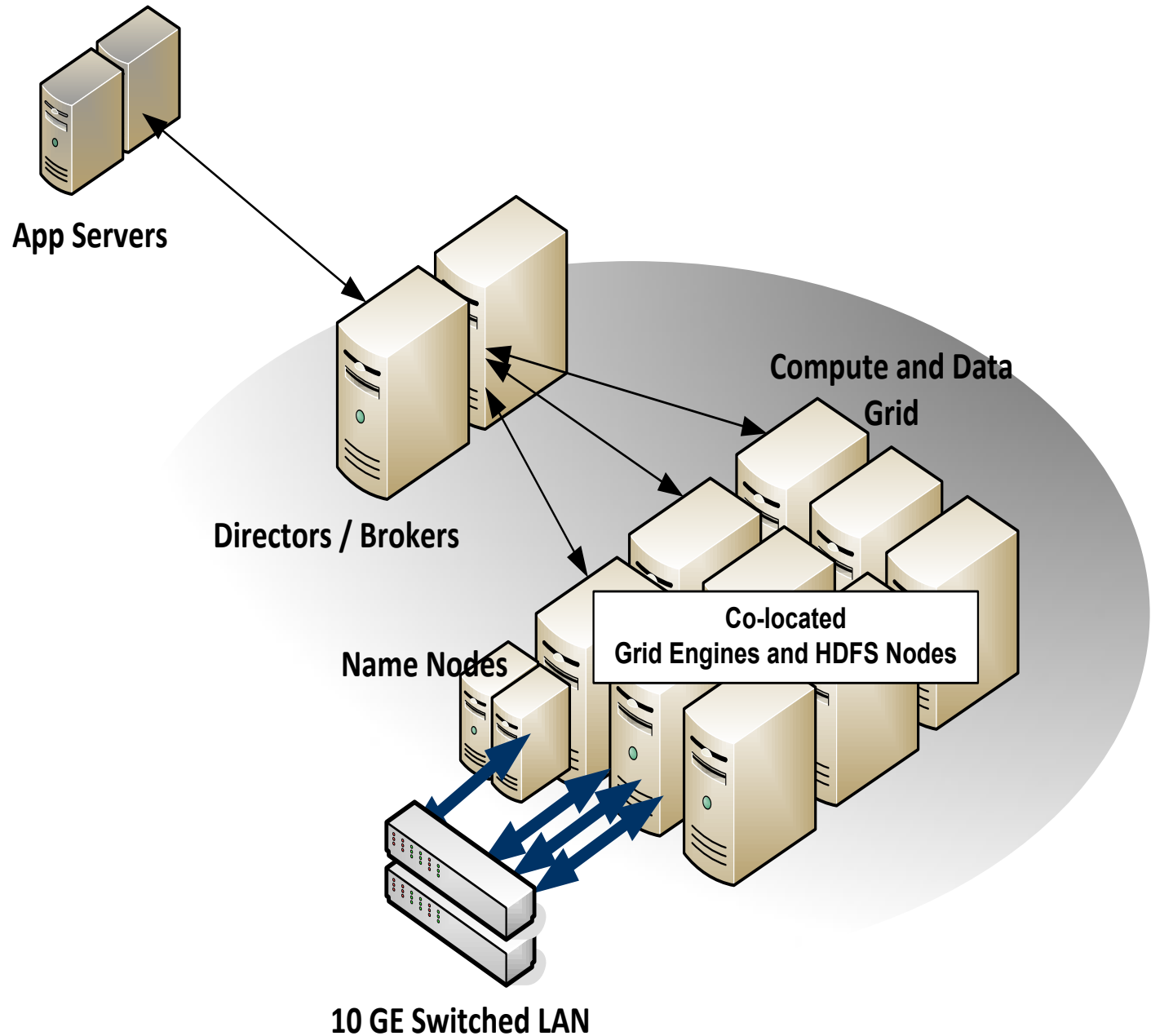
Varying replication count significantly affects average read speed:

3 of 14 = ~ 50MB/s      14 of 14 = ~ 100MB/s  
**Replication factor 3 over 14 nodes**



**Replication factor 14 over 14 nodes**





# Where Next?

## Software Optimisations

Data Structures

Distribution Algorithms

Uncorrelated and Correlated Random number generation

## Data Compression

## Novel Hardware Architectures



££ / Compute

Optimisations are being carried out continuously to reduce costs

# Summary

Risk is challenging and interesting!

Our focus is accuracy, speed and cost

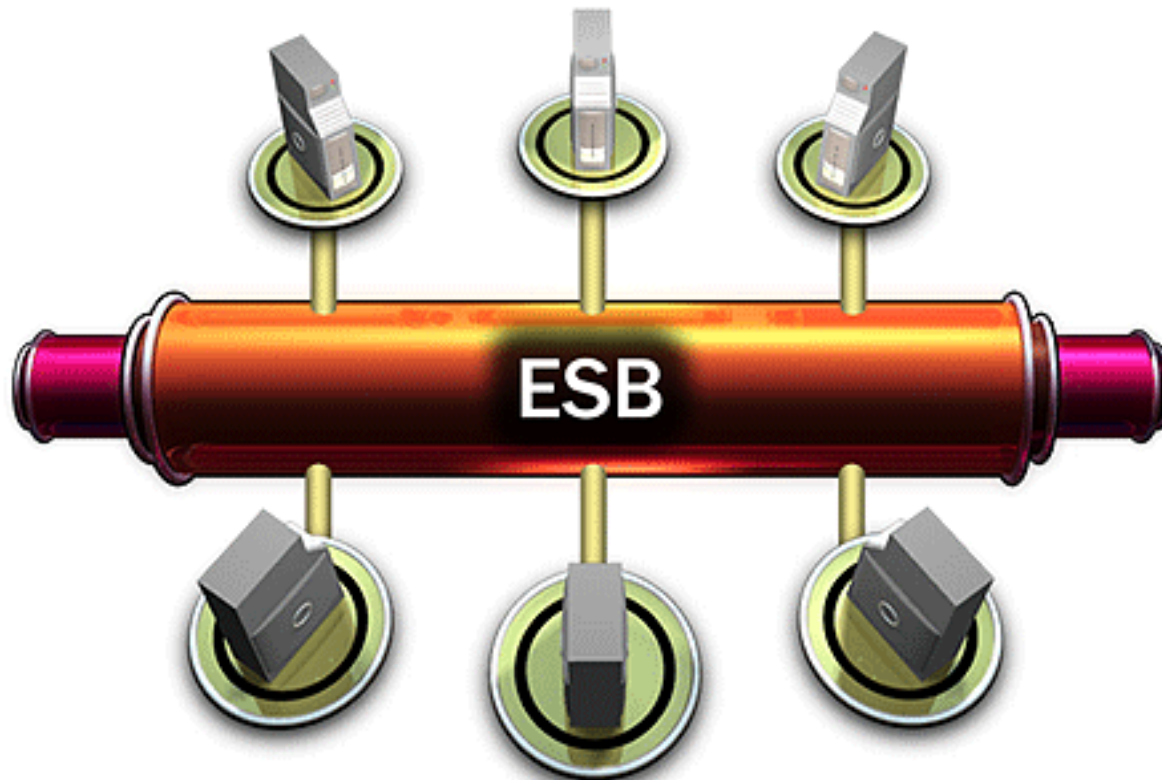
We are innovating methodologies, architecture & engineering

Thank you

ODC

One Version of the Truth For Everyone

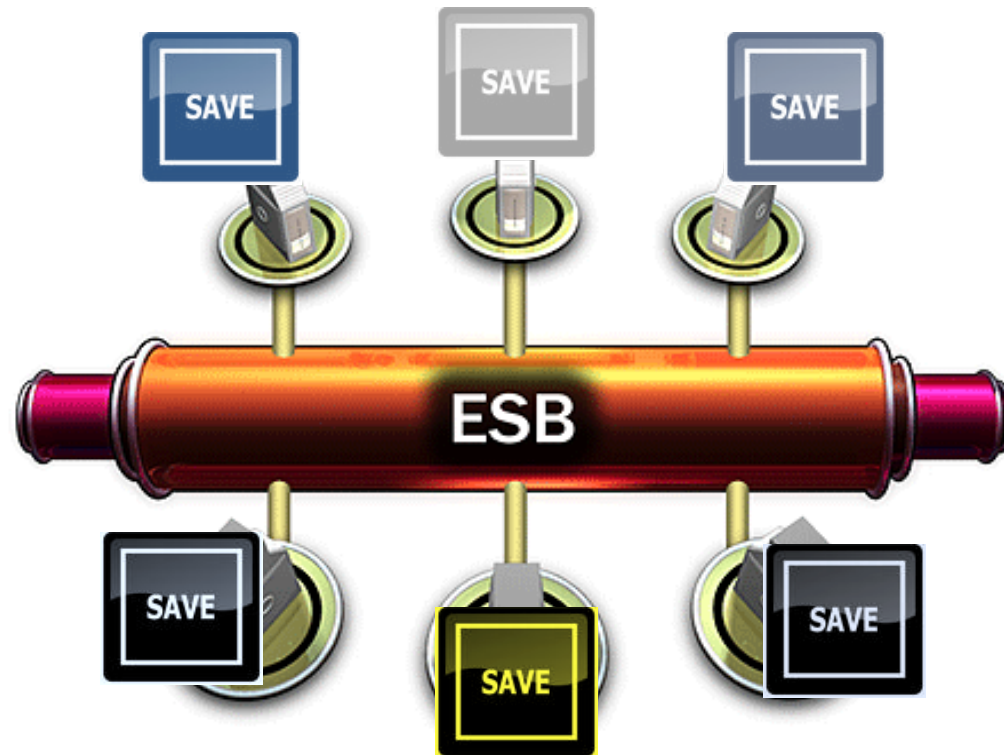
# Messaging allows us to disseminate facts



...but messaging has its limits

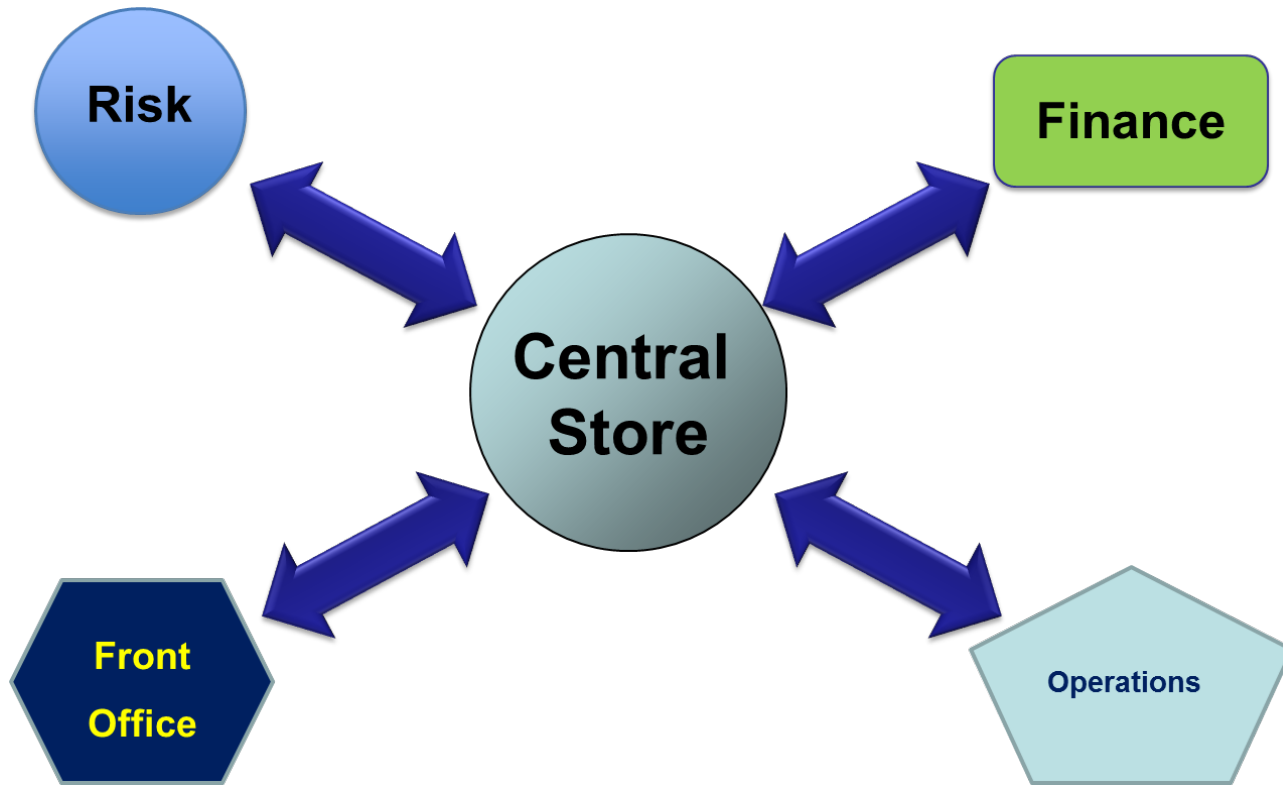


It leaves interpretation of different facts to the various consumers





# Centralising data gets all eyes on a single version of the truth



# Originating from concept of a central ‘brain’ within the bank:

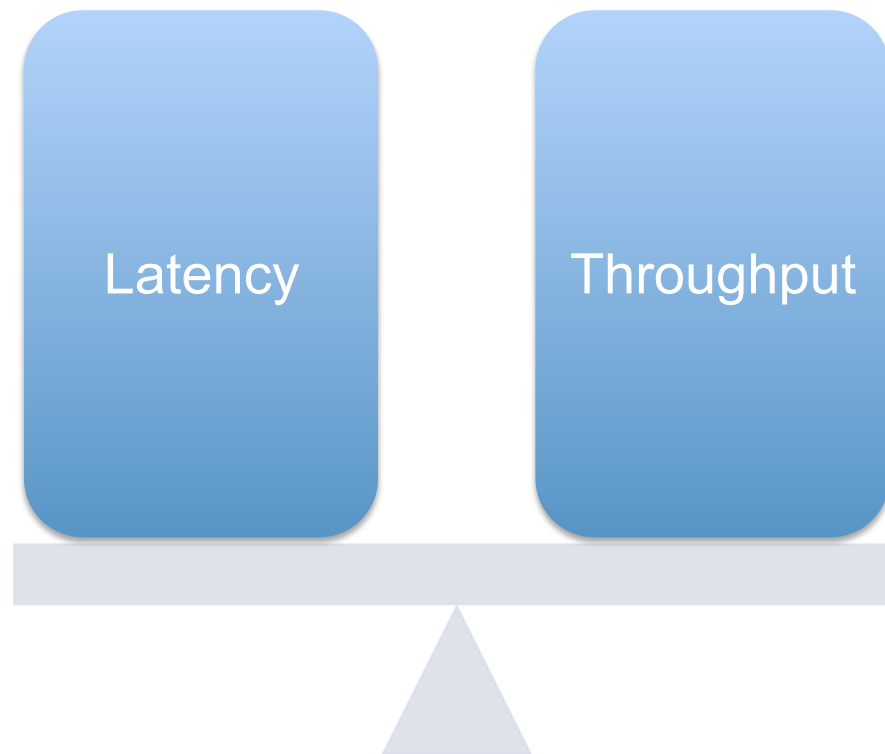
*“The copying of data lies at the route of many of the bank’s problems. By supplying a single real-time view that all systems can interface with we remove the need for reconciliation and promote the concept of truly shared services” - Scott Marcar (Head of Risk and Finance Technology)*



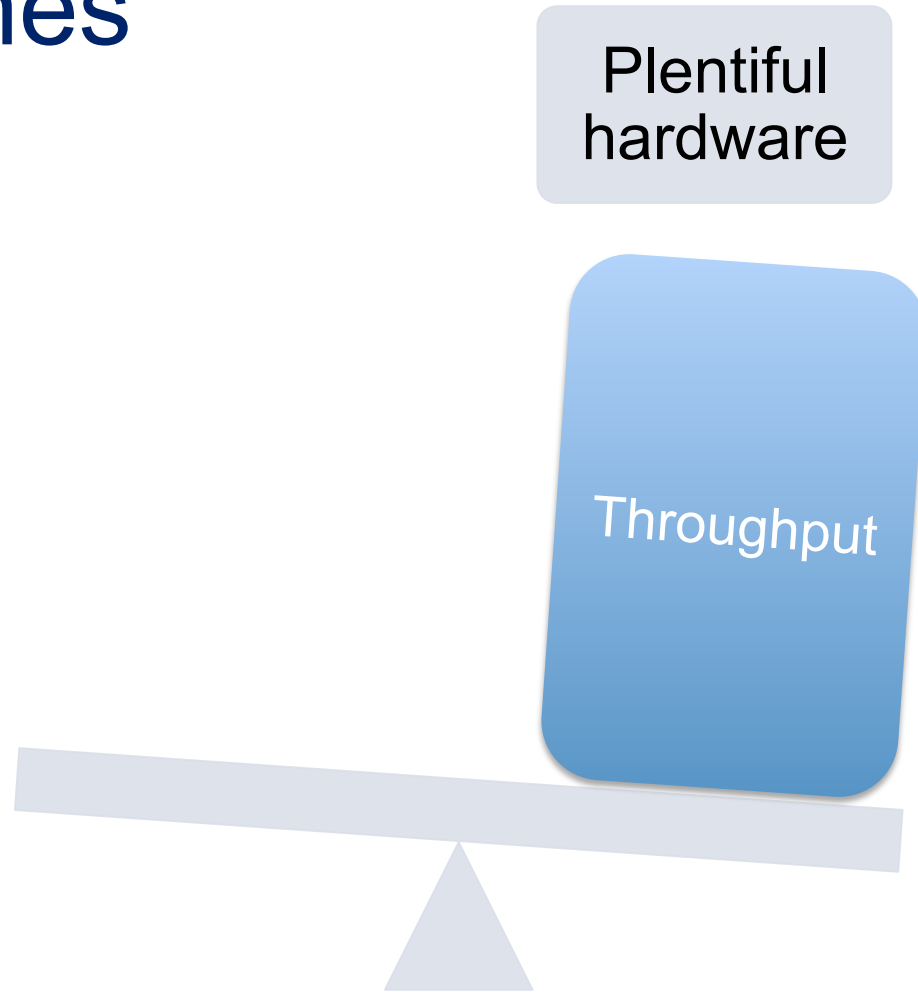
# Nice idea, but how do you do it?



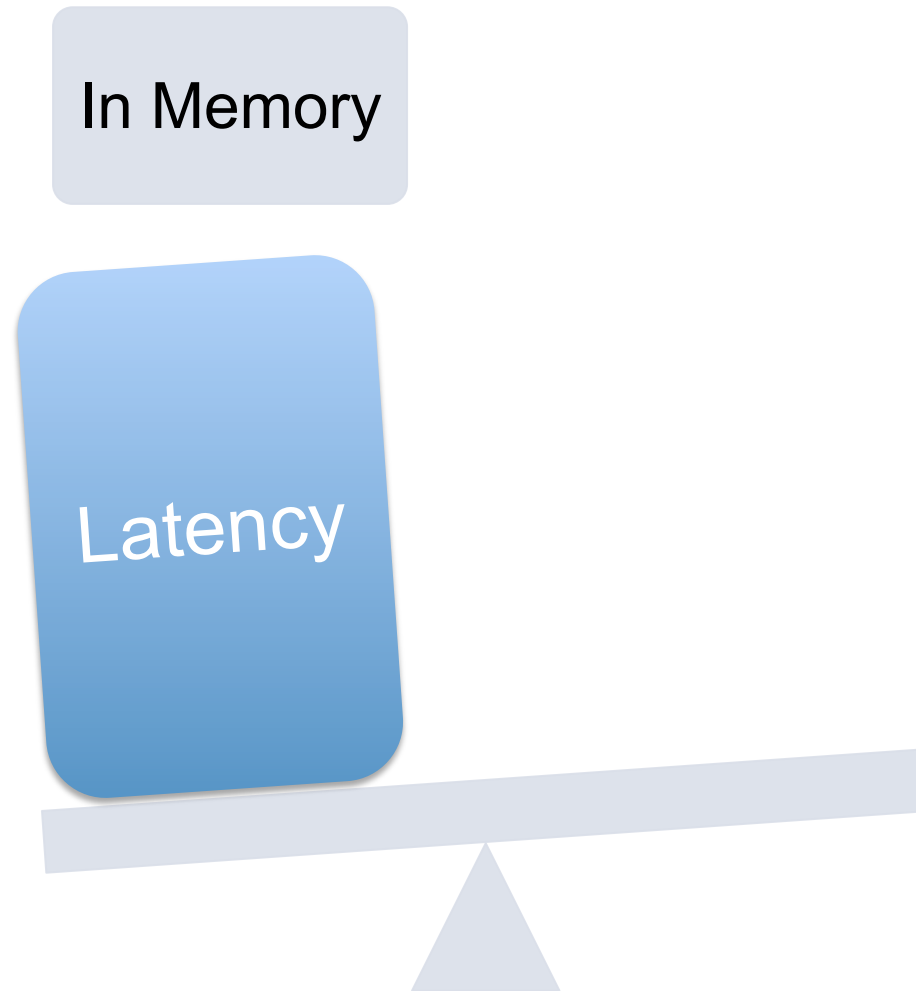
# Needs a balance of low latency and high throughput



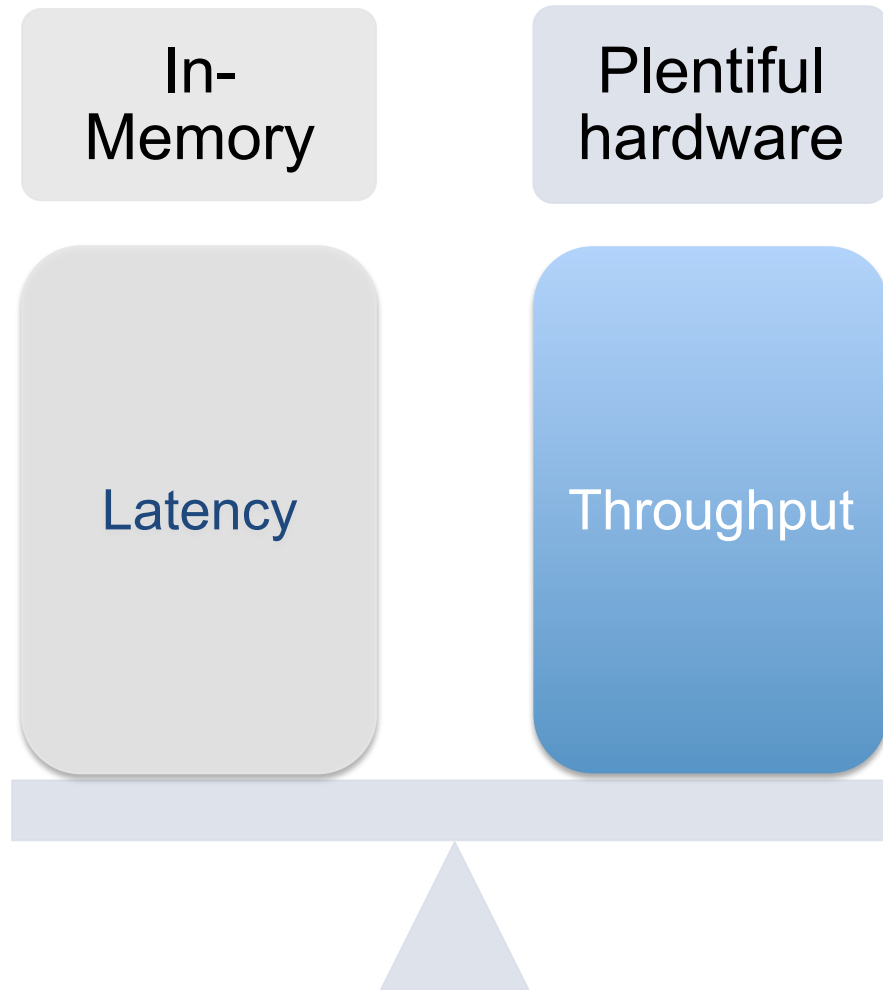
# Throughput generally equates to lots of machines



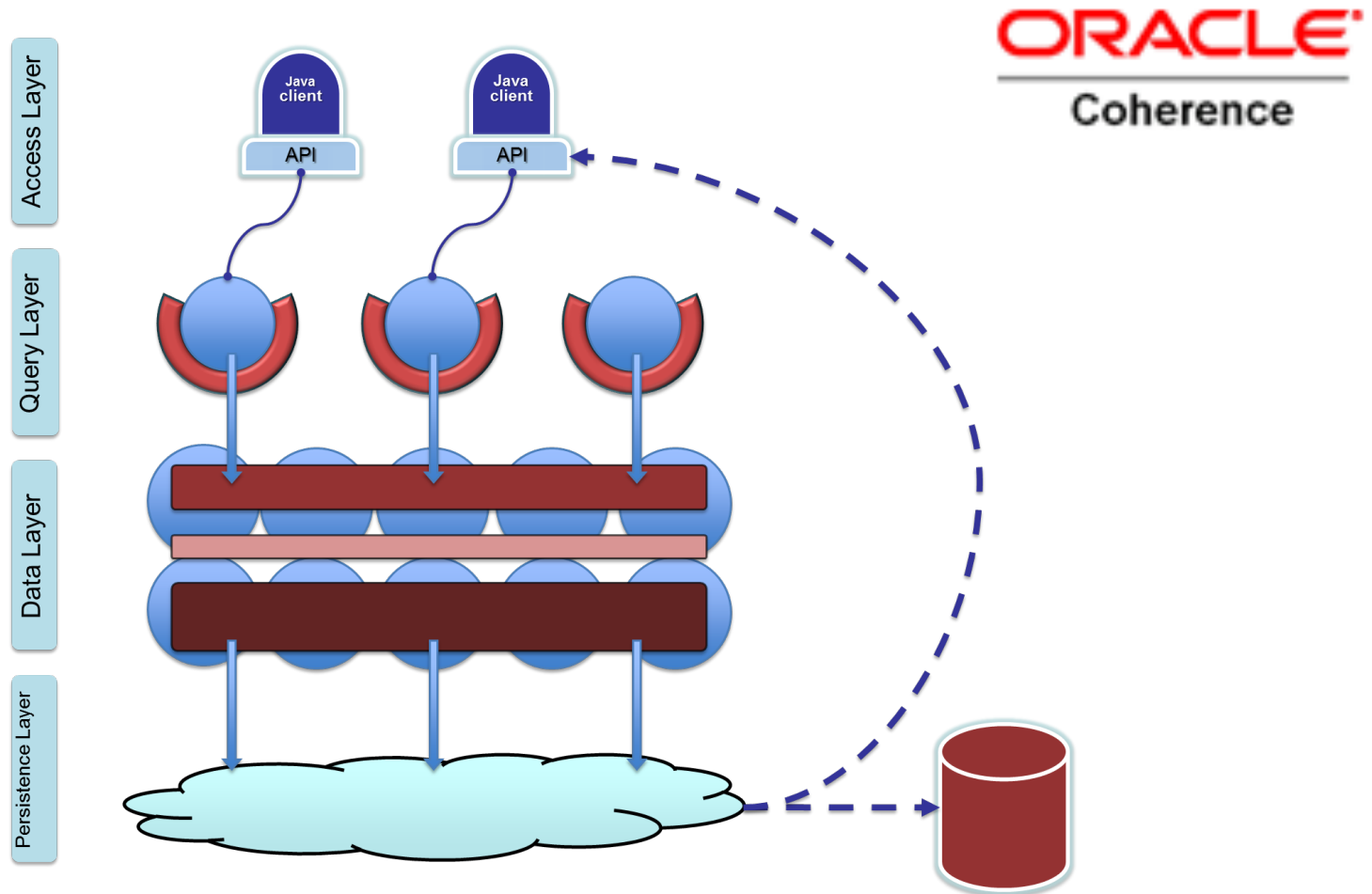
# Low latency = in-memory / keyed access



# ODC is a balance

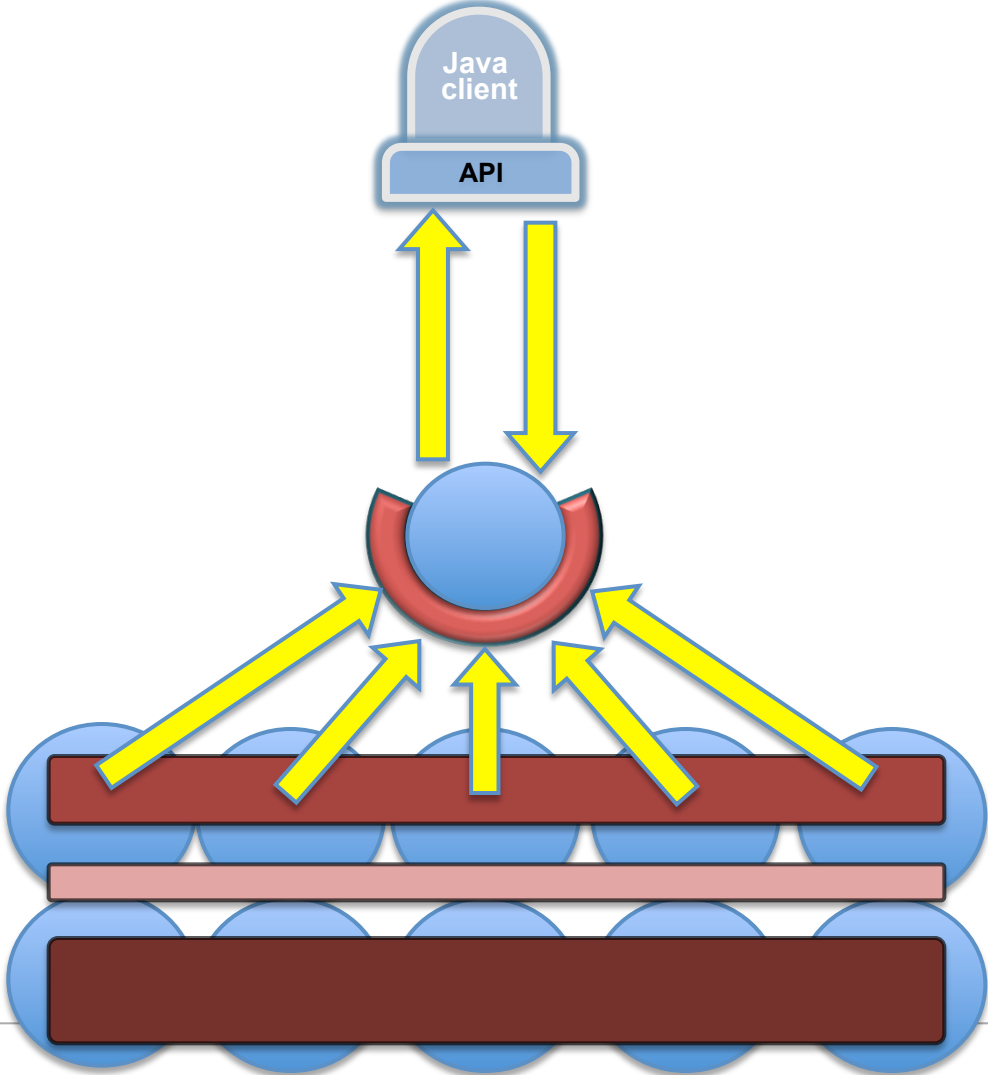


# Oracle Coherence provides the data fabric





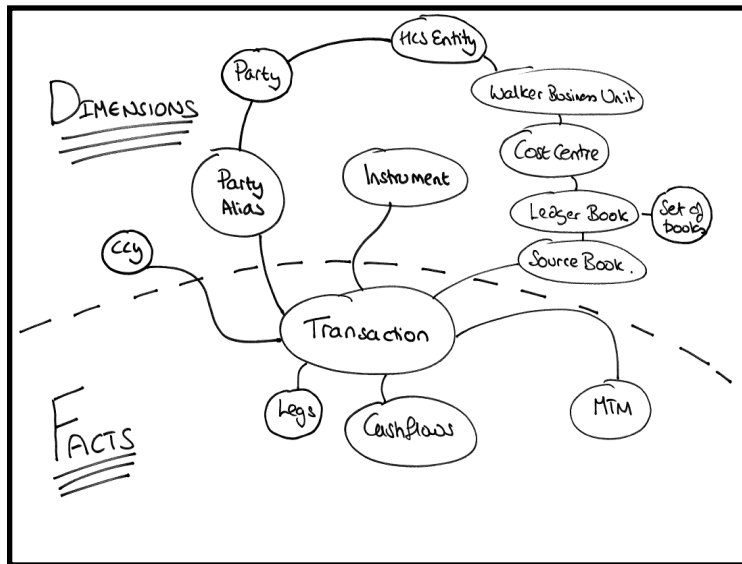
# Coherence is an In-Memory (NoSQL) caching technology with a Shared-Nothing Architecture



# ...but ODC adds normalisation... ...which means it does joins!

```
QueryDefinition query = odc.fromTransaction()  
    .where(valueOf(transaction().getSourceBook().getName()).equalTo("CITI7061"))  
    .joinToValuation().where(valueOf(valuation()).equalTo(today())).createQuery();
```

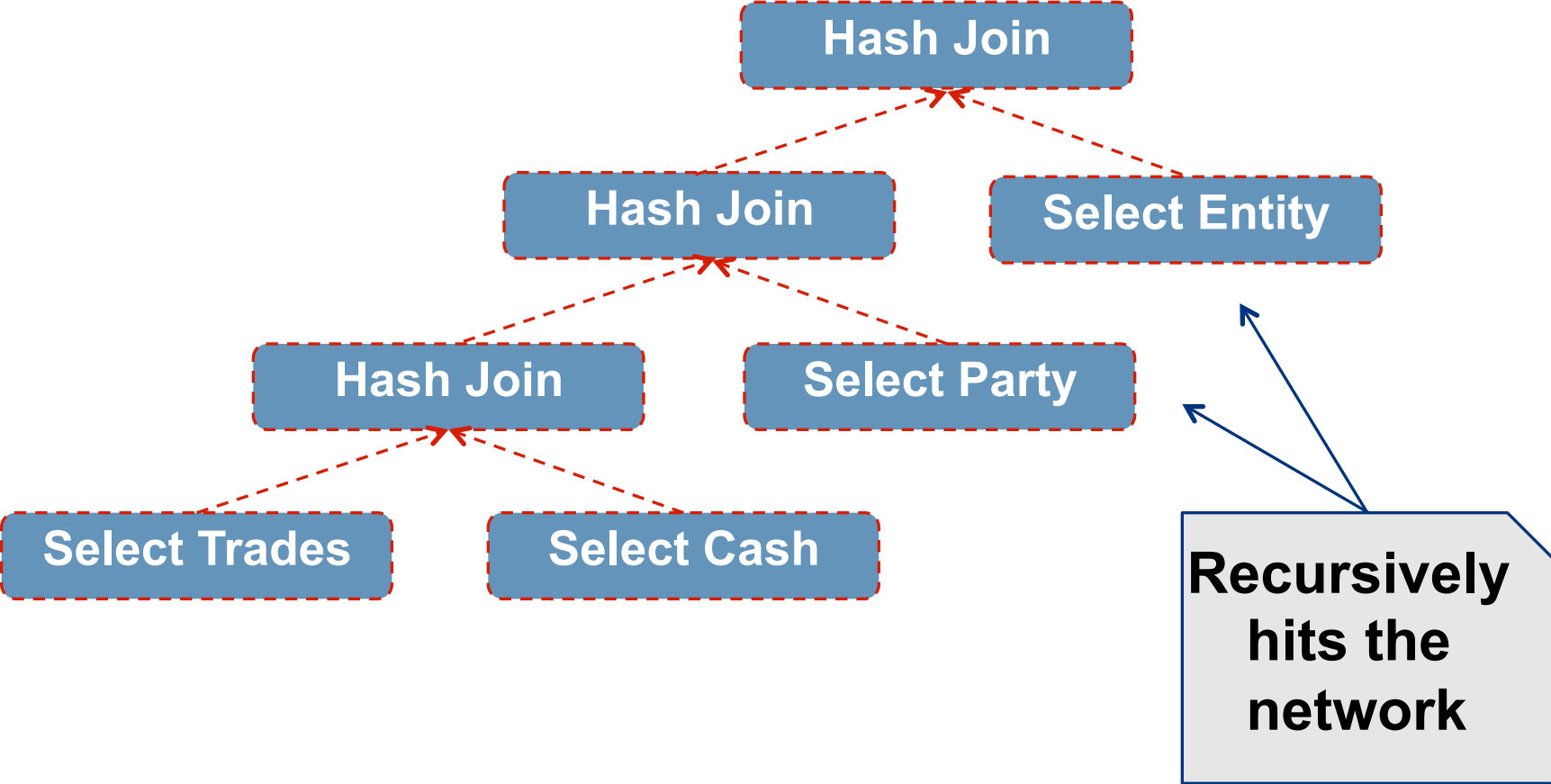
getBusinessCentre()	String
getCurveSystemId()	SystemInstanceId
getEndDateMarkedRate()	BigDecimal
getId()	DailyTransactionValuationId
getNextValuationDueDate()	BusinessDate
getProjectedEndDate()	BusinessDate
getRate()	BigDecimal
getStartDateMarkedRate()	BigDecimal



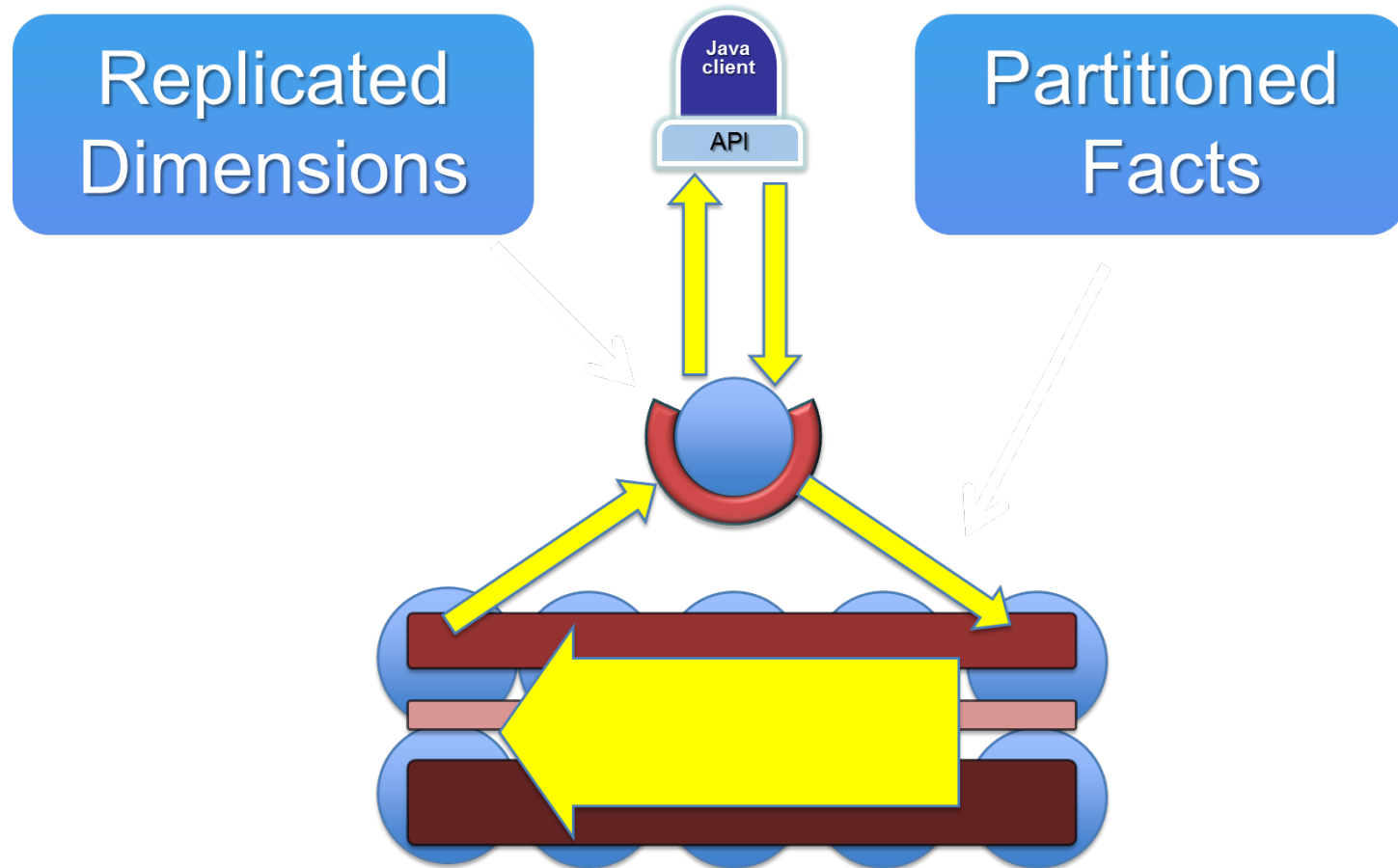
# Why Joins?

- Objects need to be versioned as changes come from many different sources
- Recompose different slices of data (for example bi-temporally)
- Fat objects (Distributed Cache) / Documents (NoSQL) don't allow this

But the performance of arbitrary joins in a distributed architecture tend to suck!

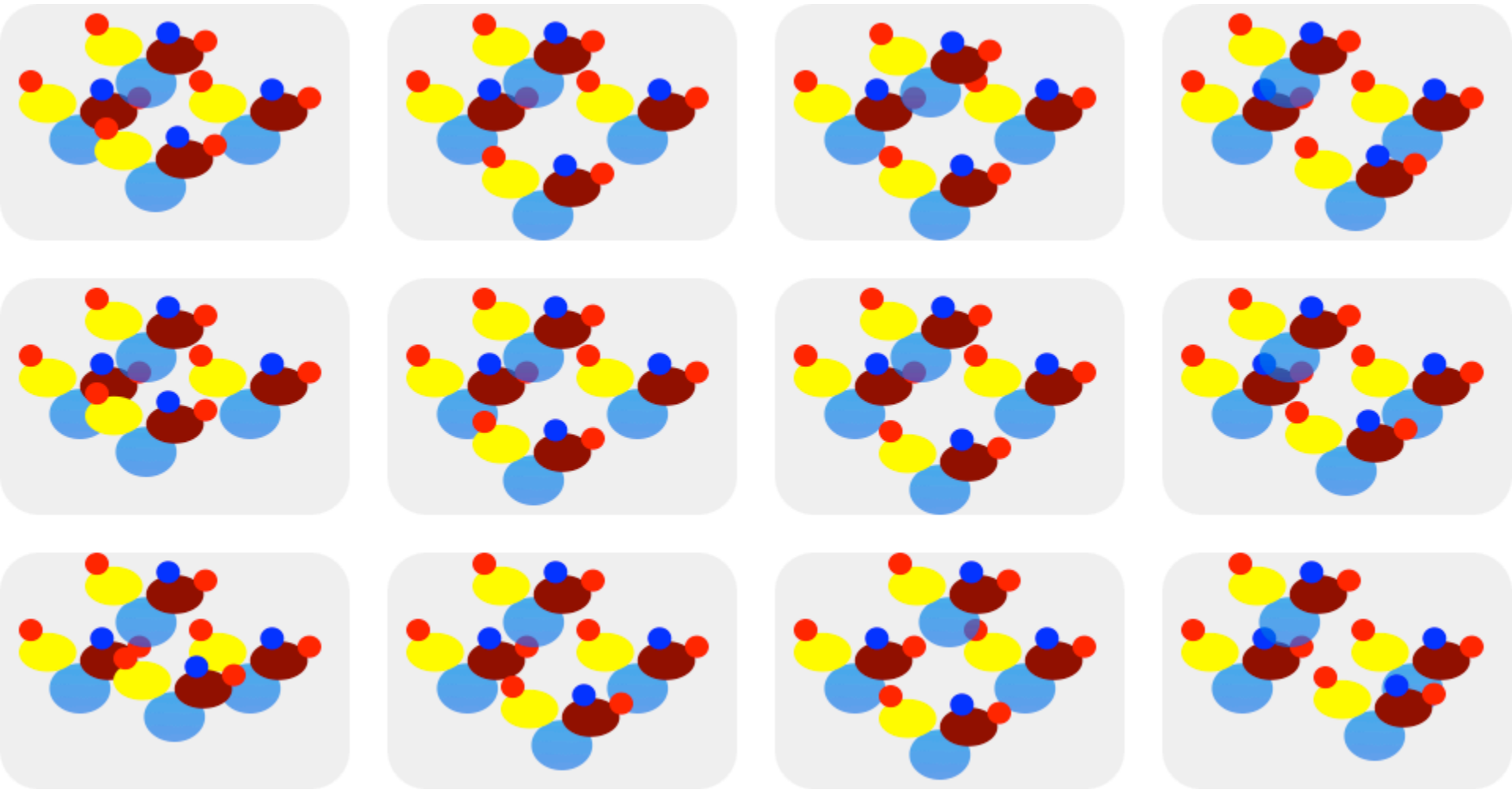


ODC prevents this by replicating data that won't shard with the primary key



We use replication to get  
around the distributed  
join problem

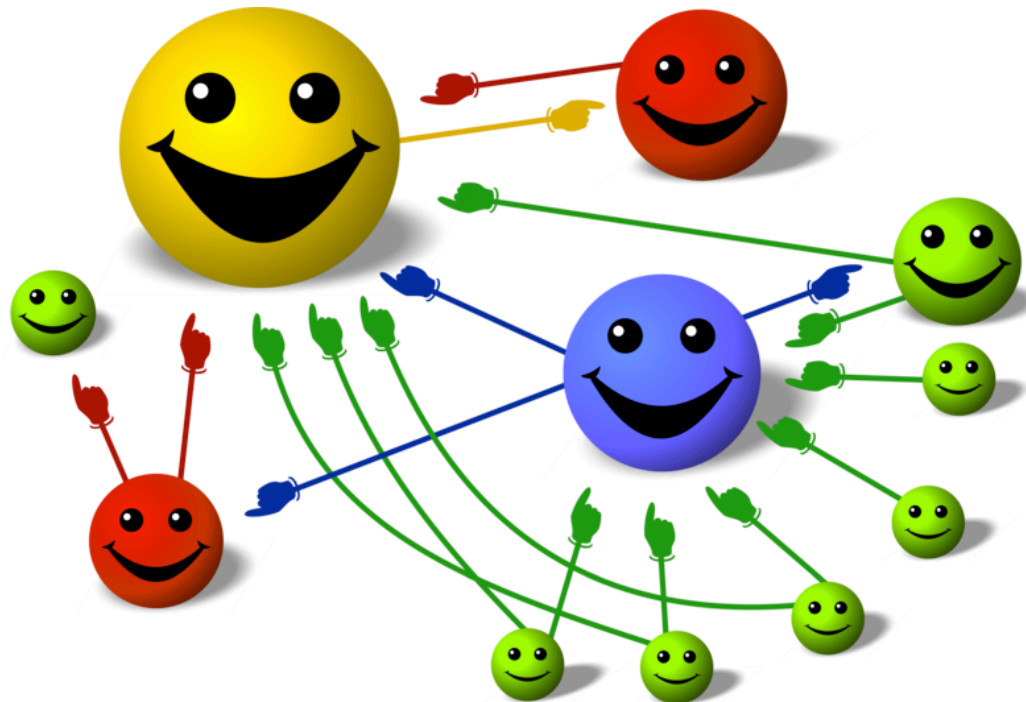
The problem with replication is that it uses up all your RAM



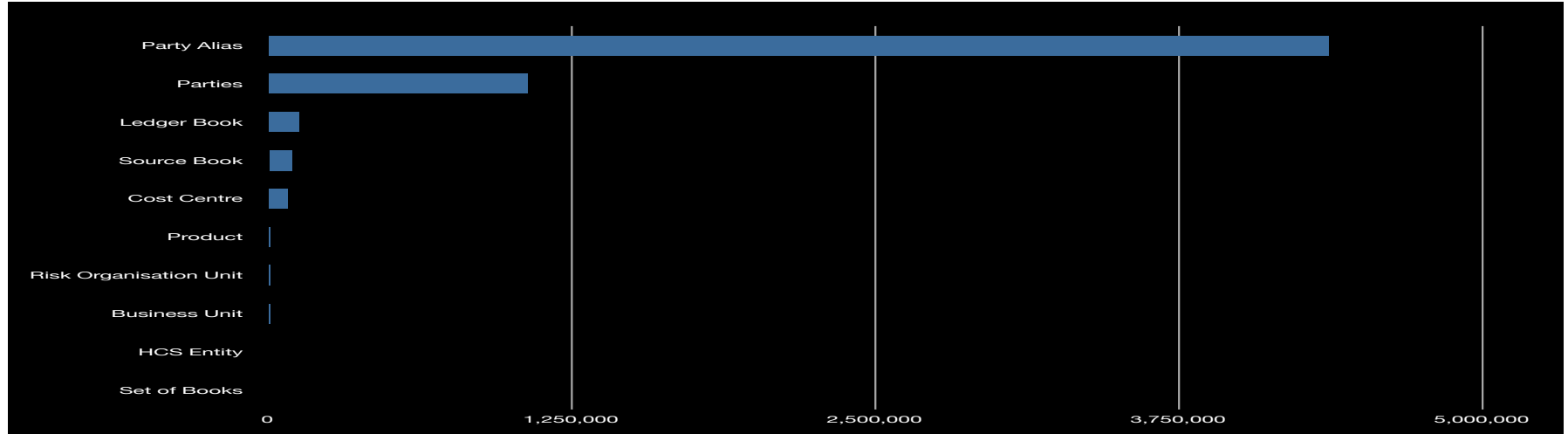
Replication eats into  
your storage!



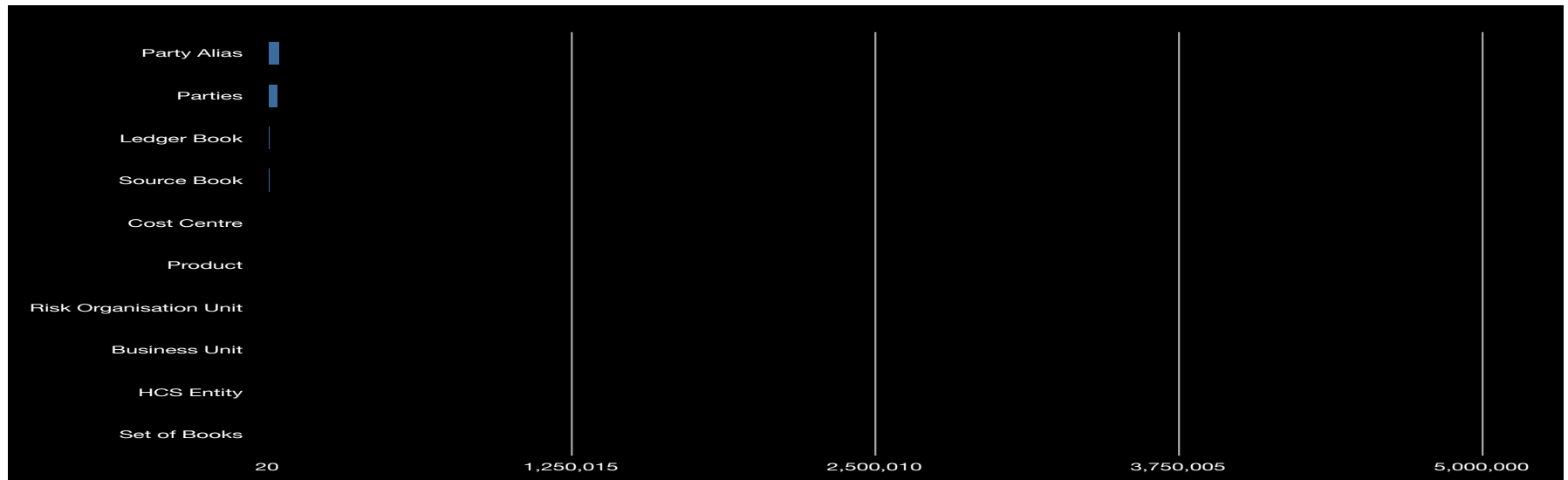
# 'Connected Replication Pattern' only replicate data that is referenced



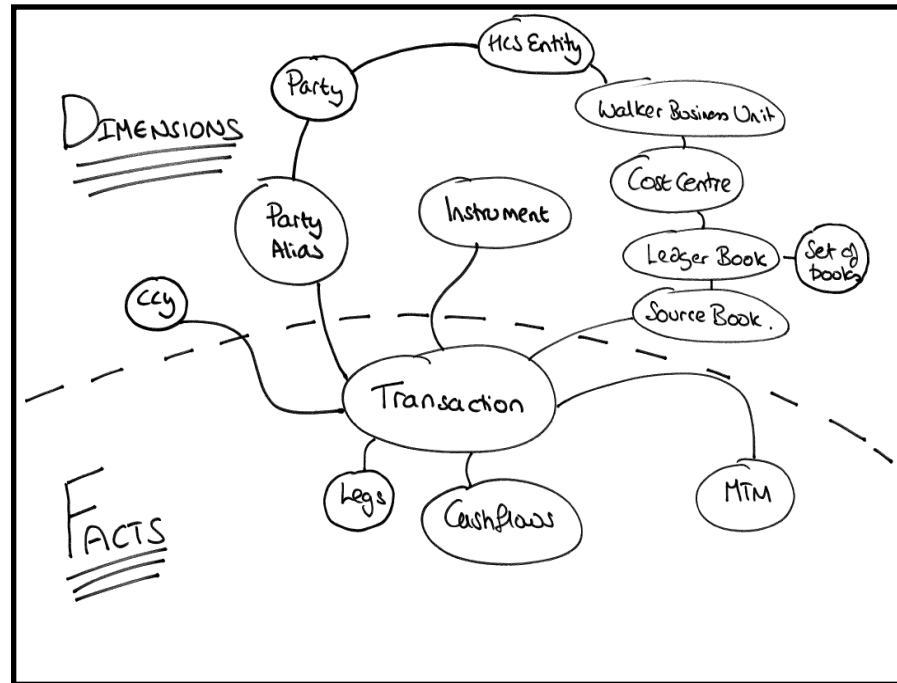
# All Data:



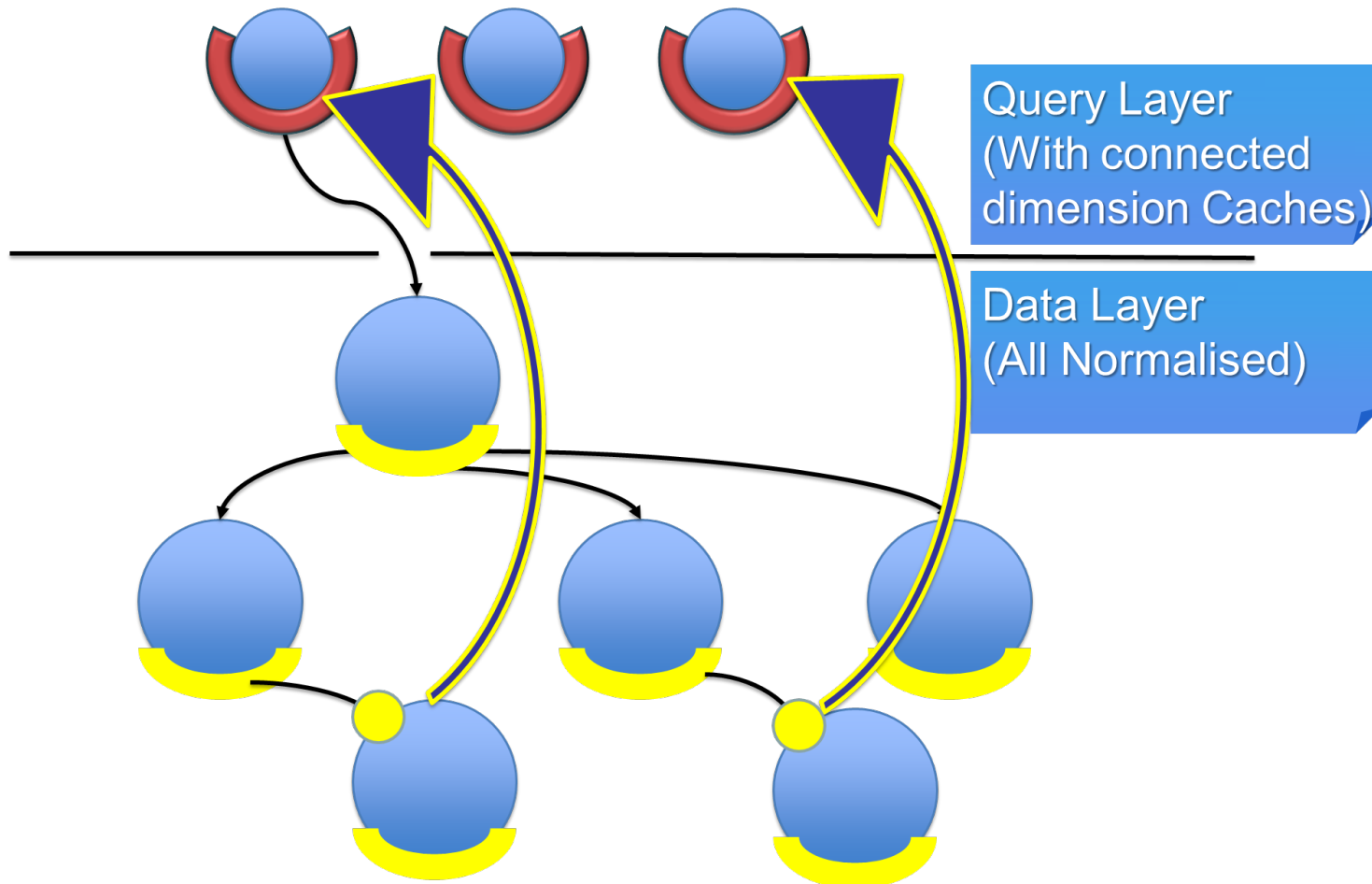
# Used Data:



# Connected Replication: Constantly track what data is used at a point in time, and only replicate this 'used' data

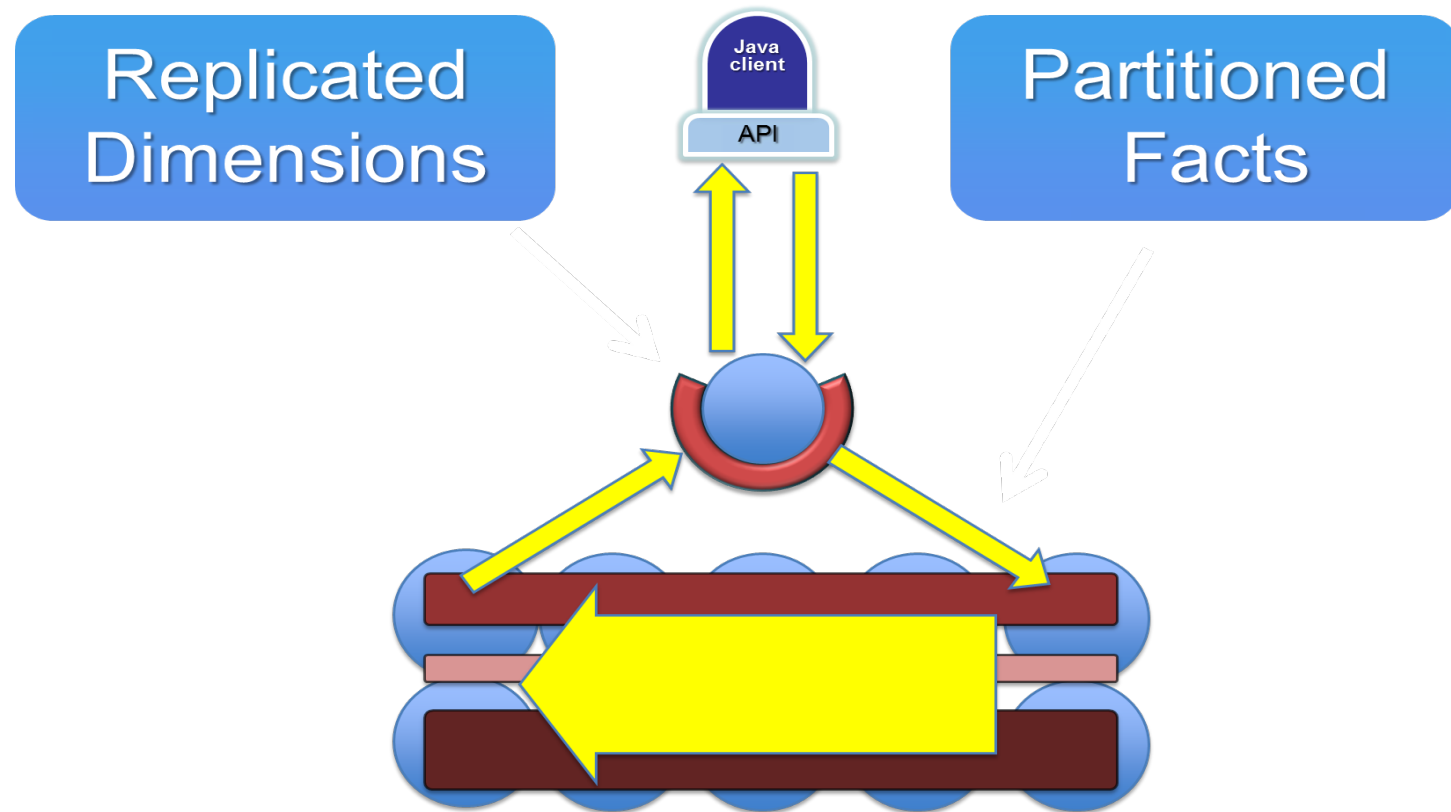


Writes cause items to become replicated if they were previously 'unused'



With 'Connected Replication' only  $1/10^{\text{th}}$  of the data needs to be replicated, which means we can do fast joins without eating into that much memory

It's a two tier approach: An In-Memory database on top, with a sharded In-memory database below



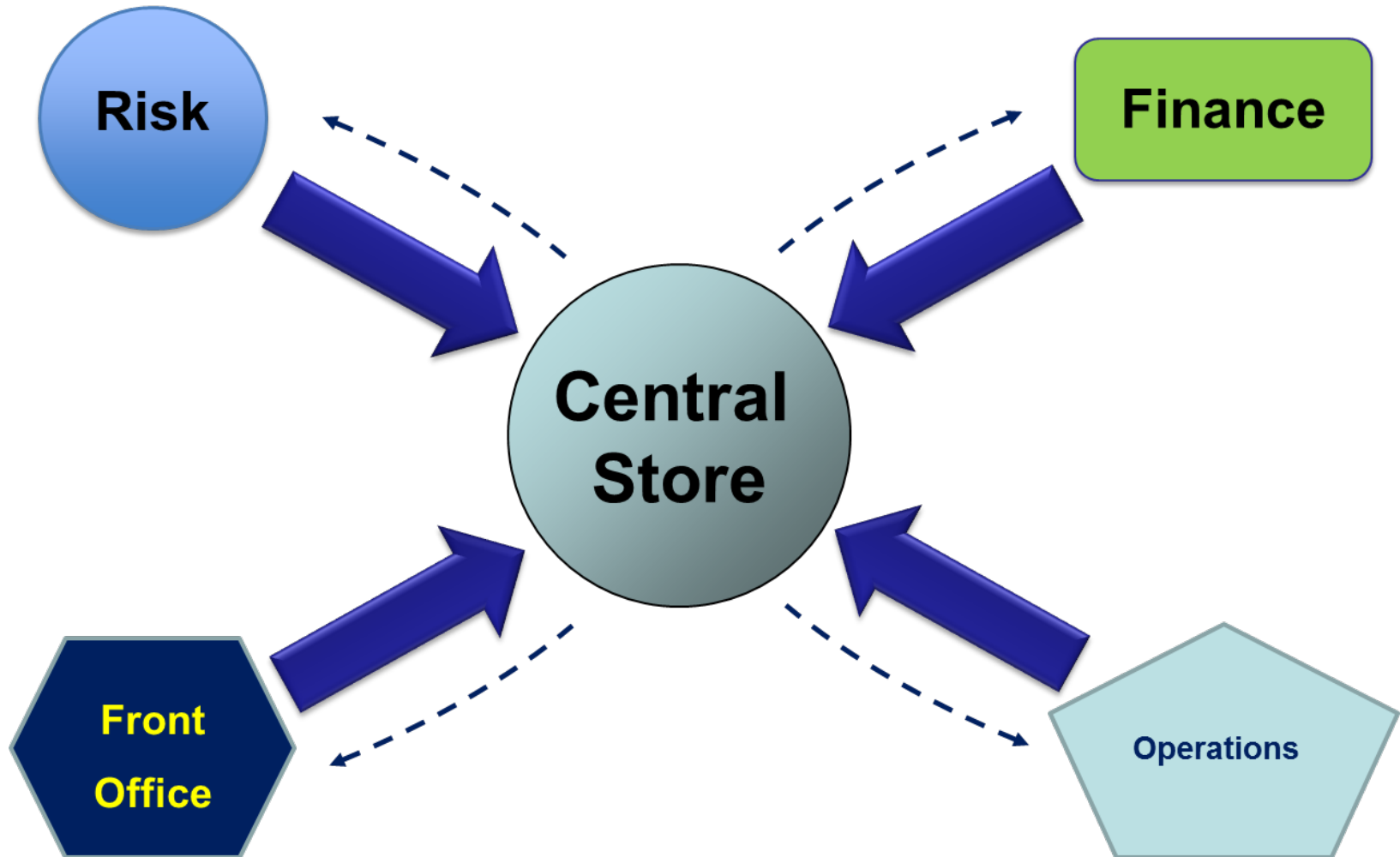
for example ODC can do  
a 34 join query data in a  
420 node cluster in 5ms  
(key based) and around  
20ms (broadcast)

So ODC provides

- Scalability
- Real-time interface
- Low latency access
- High throughput (throttled)



# One centralised version of the truth!



# Data Virtualisation



Why do Data Virtualisation?

Resisting the ‘YAS – Yet Another System (...and database & feeds & code &&&)’ Syndrome.

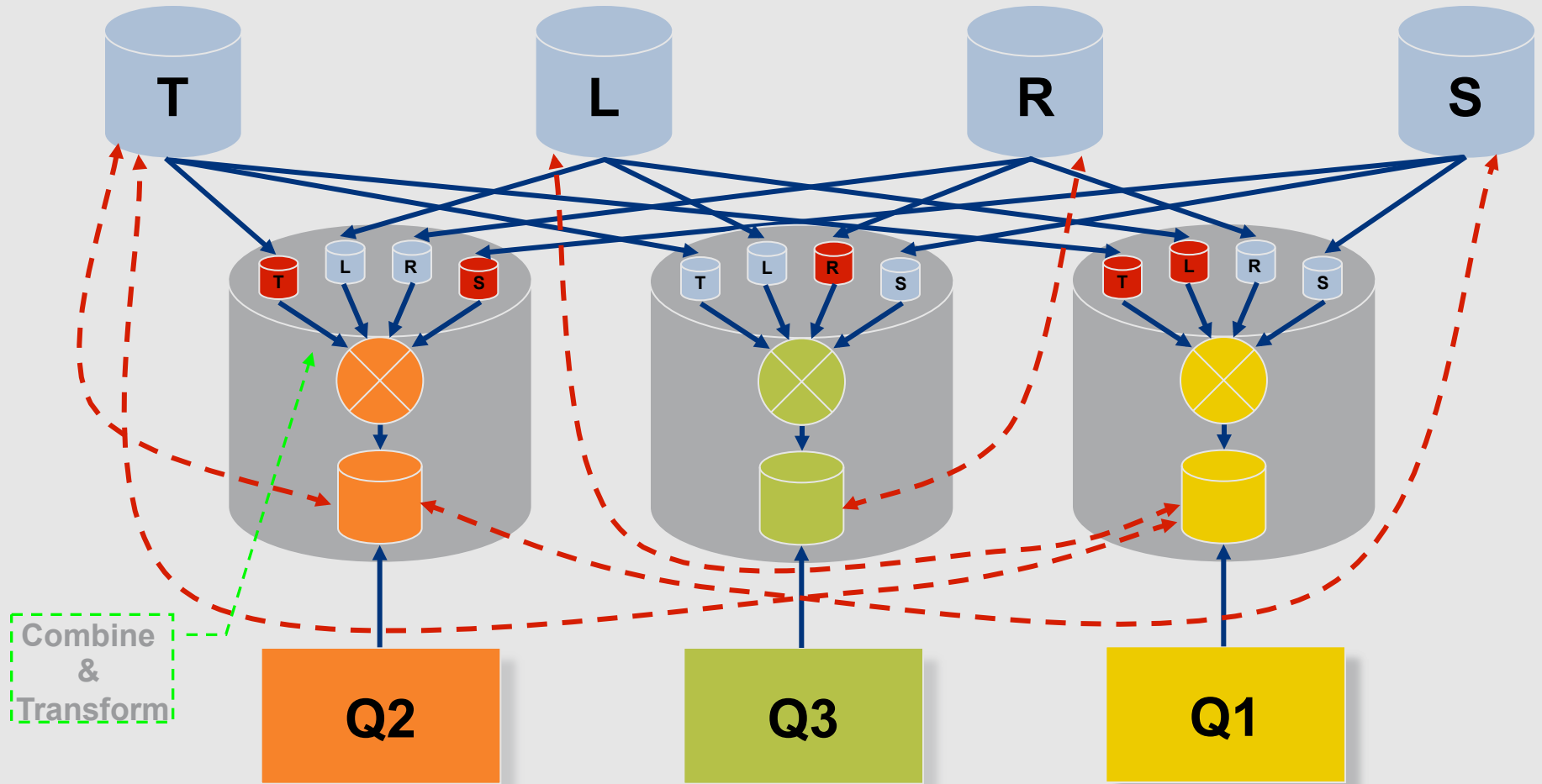
Aspirations:

- Leave data in situ
- Define an idealized data-schema

What’s different? The DAL – is not a point-solution, it’s a strategic platform and data-source

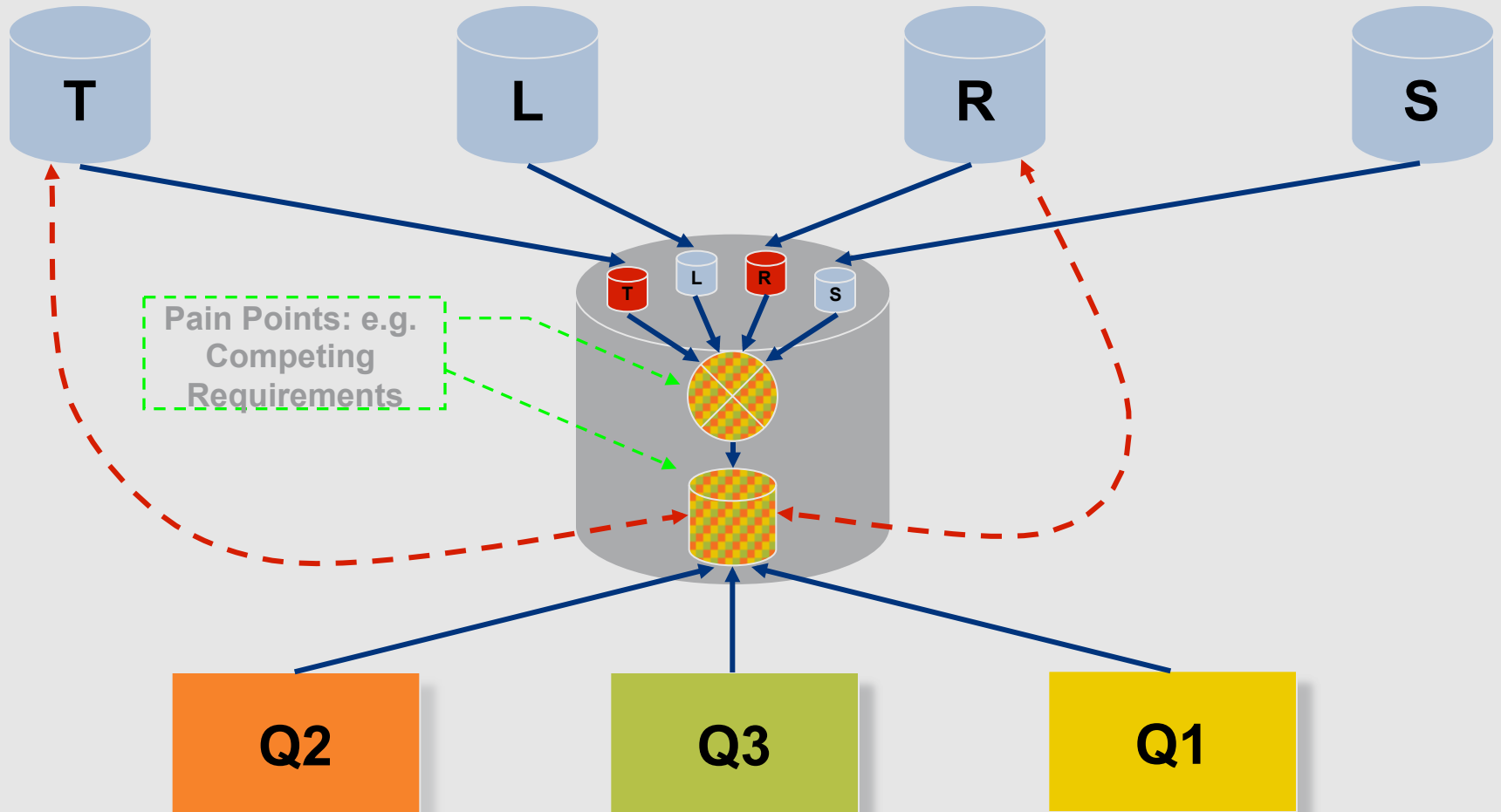


# YAS: Sources, Feeds, Replicas (Subsets) and Reconciliations



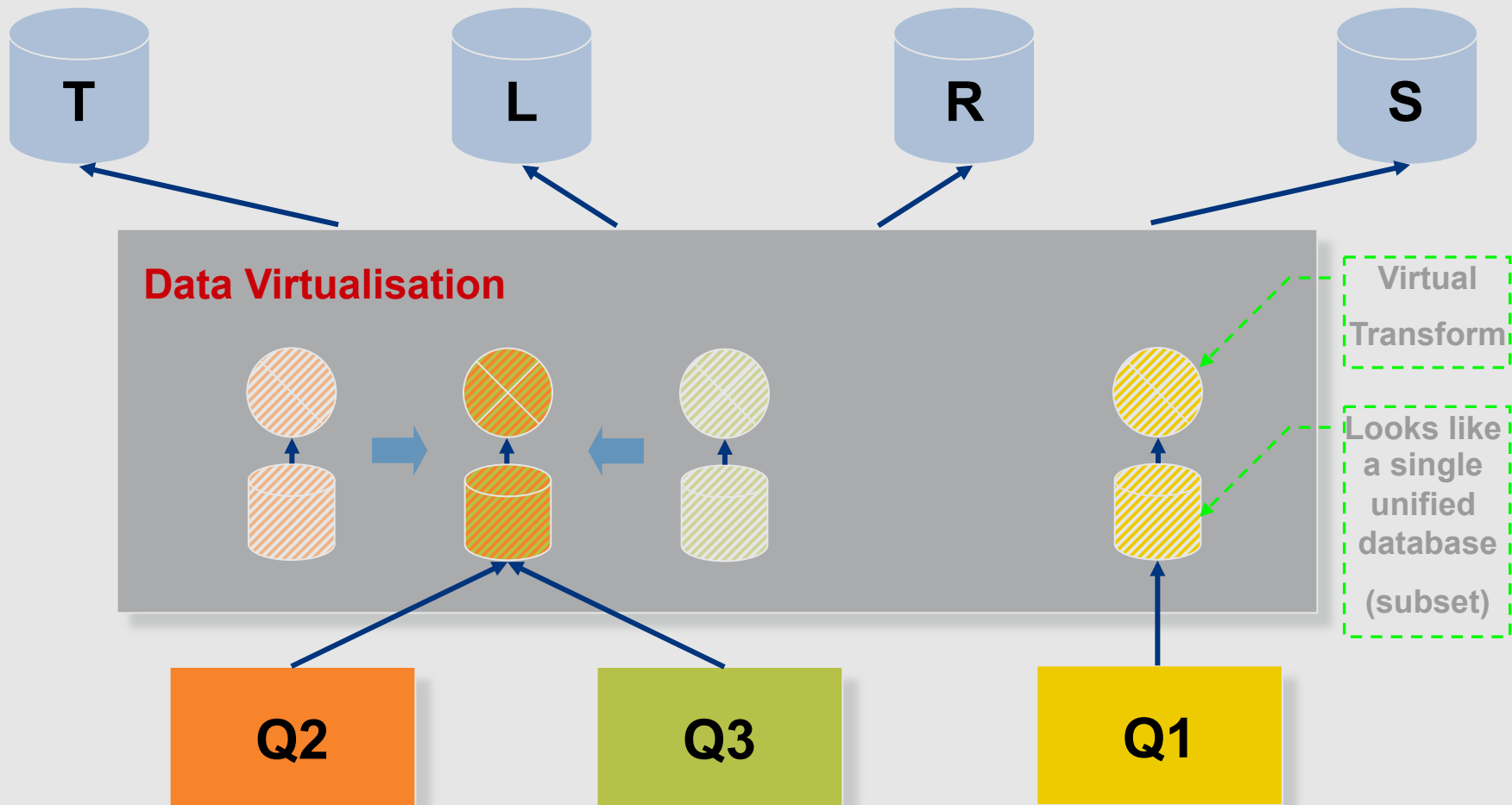
# Does the Enterprise Data Warehouse (EDW) Help?

Co-ordination of (potentially competing) requirements & releases, reconciliations still required



# Data Virtualisation

...and if synergies are identified, transitions can be straightforward:



# The Opportunity to Think Differently

However: It's not without its challenges...

## Technical:

Distributed joins & query optimisations are getting better

Pesky physics still get in the way!

Real World™ Data Consistency and Quality

'Unusual' sources require adapter development

Some sources can't (or don't want to) be hit directly

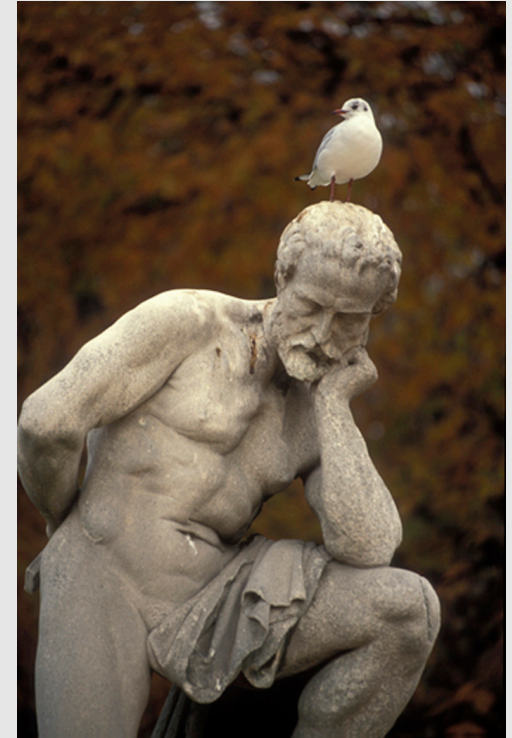
So in practice, we expect to create a hybrid **but all 'behind' the DAL.**

The developers in this space become Enterprise Data Engineers

So our aspirations are treated as our idealized design principles

## People:

New & Unfamiliar 'Stuff' often generates its own resistance



# Collaboration is Key

What we've learnt:

- This stuff is best shared!
- It lowers resistance to adoption.

Collaboration and Socialization is crucial:

- Guiding Principles / Thinking Frameworks
- 'Brown-Bag' Sessions
- Design & Architecture Discussions and Reviews
- Building the Network

**Effort invested in these types of activity  
always seems to pay off!**



Questions?

Thank you for your time!

