Pragmatism, Puritanism and Functional Programming

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Software in Industry ...



- Software in Industry ...
- * ... is broken.



- Software in Industry ...
- * ... is broken.
- * ... because it's too complex!



- Software in Industry ...
- ... is broken.
- * ... because it's too complex!



- My perspective on why FP techniques can help
 - * Based on 8 years of Haskell (5 full-time commercial)

- * Haskell makes *programs*:
 - Simpler to Write
 - * Simpler to Read (ie understand)

- * Haskell makes *programs*:
 - Simpler to Write
 - * Simpler to Read (ie understand)
- * Is Haskell simple?

- * Haskell makes *programs*:
 - Simpler to Write
 - * Simpler to Read (ie understand)
- * Is Haskell simple?
- * Is Haskell easy?

A Better Language?



What is important about FP?

- Purity
- Type Systems
- Higher Order Functions (HOFs)
- * Laziness
- Extensive abstract libraries

What makes FP hard to learn?

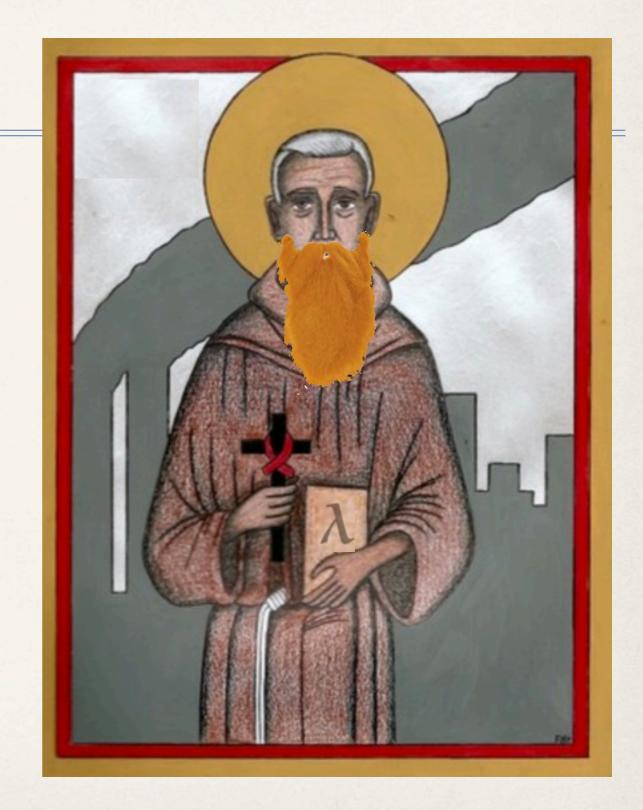
What makes FP hard to learn?

- Purity
- Type Systems
- * HOFs
- Laziness
- Extensive abstract libraries

Purity

"The functional programmer sounds rather like a medieval monk, denying himself the pleasures of life in the hope that it will make him virtuous."

-- John Hughes [~1984]



Modularity

- We build systems by naming and assembling chunks of code
- Interested in:
 - * Result of a chunk
 - * What a chunk Does

```
var y = chunk1(a,b);
var z = chunk2(c);
chunk3(x);
chunk4(y,z);
```

v0.9

```
var x = chunk1(a,b);
var y = chunk2(c);
chunk3(x,y);
```

v1.0

```
var x = chunk1(a,b);
var y = chunk2(c);
chunk3(x,y);
```

v1.0

```
var x = chunk1(a,b);
var y = chunk2(c);
chunk3(x,y);
```

v2.0

```
var y = chunk2(c);
var x = chunk1(a,b);
chunk3(x,y);
```

```
v1.0
```

```
var x = chunk1(a,b);
var y = chunk2(c);
chunk3(x,y);
```

v2.0

```
var y = chunk2(c);
var x = chunk1(a,b);
chunk3(x,y);
```

```
var x = chunk4(a,b);
var y = chunk4(a,b);
chunk5(x,y);
```

```
var x = chunk4(a,b);
var y = x;
chunk5(x,y);
```

What is Purity?

```
var x = chunk1(a,b);
var y = chunk2(c);
chunk3(x,y);
```

```
var x = chunk4(a,b);
var y = chunk4(a,b);
chunk5(x,y);
```

What is Purity?

```
var x = chunk1(a,b);
var y = chunk2(c);
chunk3(x,y);
```

```
var x = chunk4(a,b);
var y = chunk4(a,b);
chunk5(x,y);
```

```
f :: A -> B
```

* Pitfall 1 : Sometimes you want chunks which have state:

```
var y = chunk1(a,b);
                      var z = chunk1(a,b);
f :: A -> B
                      chunk3(x);
                      chunk4(y,z);
                      var (y, s2) = chunk1(a,b,s);
                      var (z, s3) = chunk1(a, b, s2);
f :: (A, s) -> (B, s)
                      var 54 = \text{chunk3}(x,53);
                      var s5 = chunk4(y,z,s4);
```

* Pitfall 1 : Sometimes you *need* chunks which can fail:

```
var y = chunk1(a,b);
                     var z = chunk2(c);
f :: A -> B
                     chunk3(x);
                     chunk4(y,z);
                     case chunk1(a,b) of
                       Just y ->
                         case chunk2(c) of
                           Just z ->
f :: A -> Maybe B
                              case chunk3(x) of
                                Just () ->
                                  case chunk4(y,z)
```

Pitfall 1b: The syntax is quite verbose

```
f :: (A,s) -> (B,s)
```

```
var (y,s2) = chunk1(a,b,s);
var (z,s3) = chunk1(a,b,s2);
var s4 = chunk3(x,s3);
var s5 = chunk4(y,z,s4);
```

Pitfall 1b: The syntax is quite verbose

```
var (y, s2) = chunk1(a,b,s);
                       var(z, s3) = chunk1(a, b, s2);
f :: (A,s) \rightarrow (B,s)
                       var s4 = chunk3(x,s3);
                       var s5 = chunk4(y,z,s4);
                         y < - chunk1(a,b)
f :: A -> State s B
                         z \leftarrow chunk1(a,b)
                         chunk3(x)
                         chunk4(y,z)
```

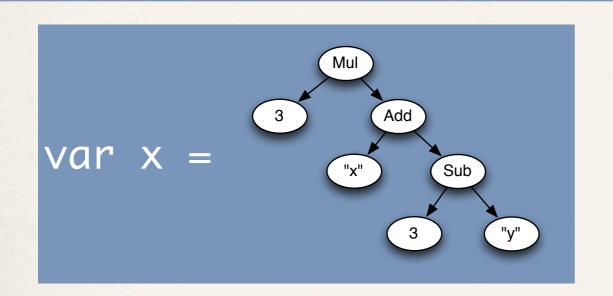
- * SIDE-EFFECTS pitfall is avoided by:
 - Take them away....
 - put them back

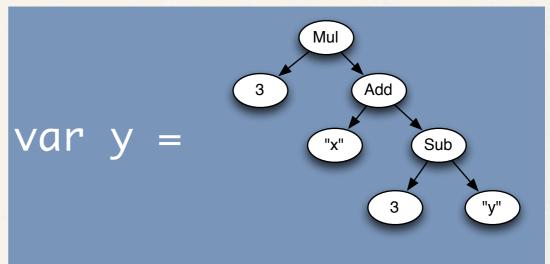
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 - * ... in a restricted way that makes the code simpler to understand!!

- SIDE-EFFECTS pitfall is avoided by:
 - Take them away....
 - put them back if necessary
 - * ... in a restricted way that makes the code simpler to understand!!
 - 1. Each chunk uses only the effects it needs
 - 2. The exact effects are visible in the type
 - 3. We can now control how different effects interact

Purity Pitfall 2 - Equality



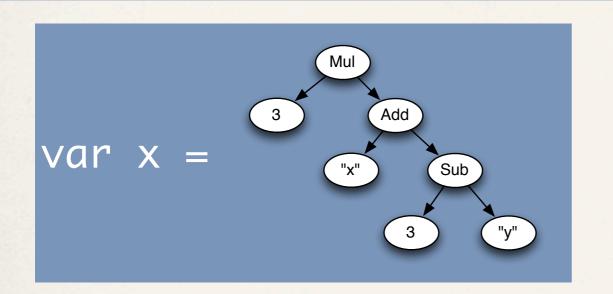


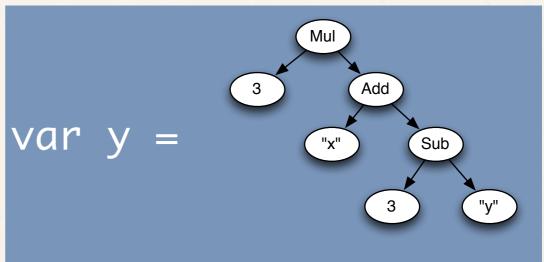
x == y

* Yes but

*

Purity Pitfall 2 - Equality



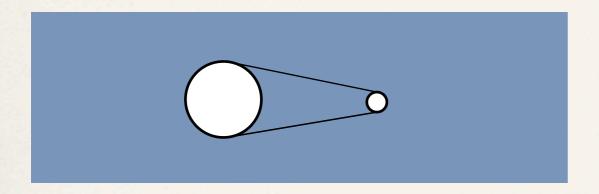


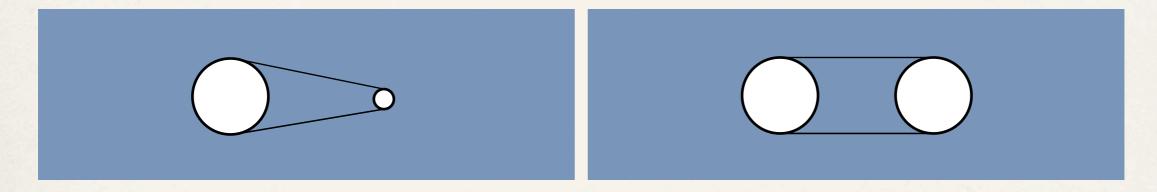
x == y

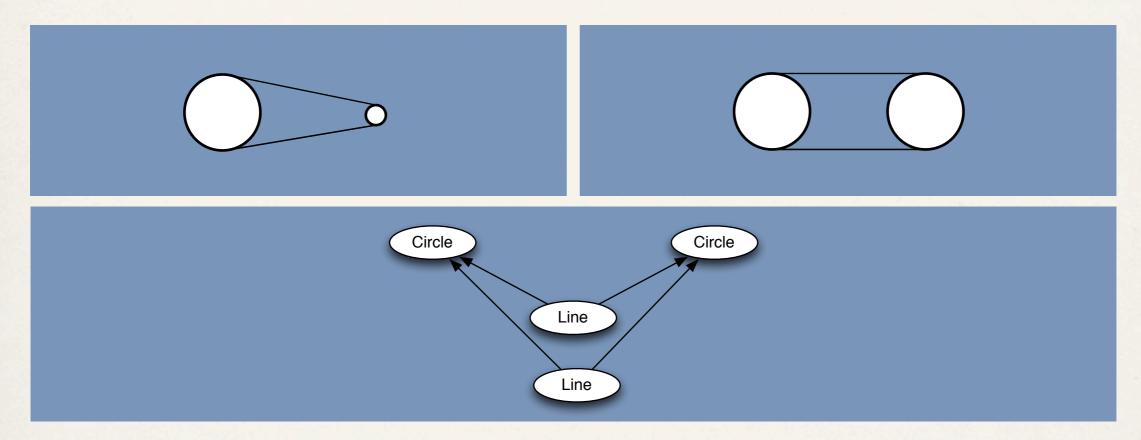
- * Yes but
- ... slow to find out

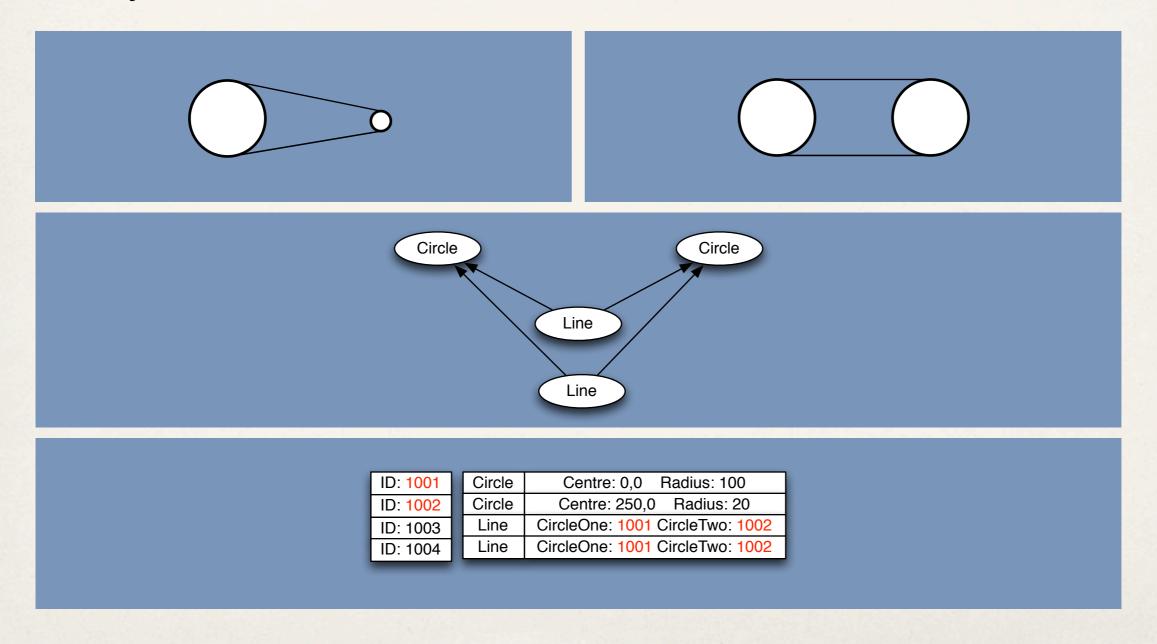
Purity Pitfall 2 - Equality

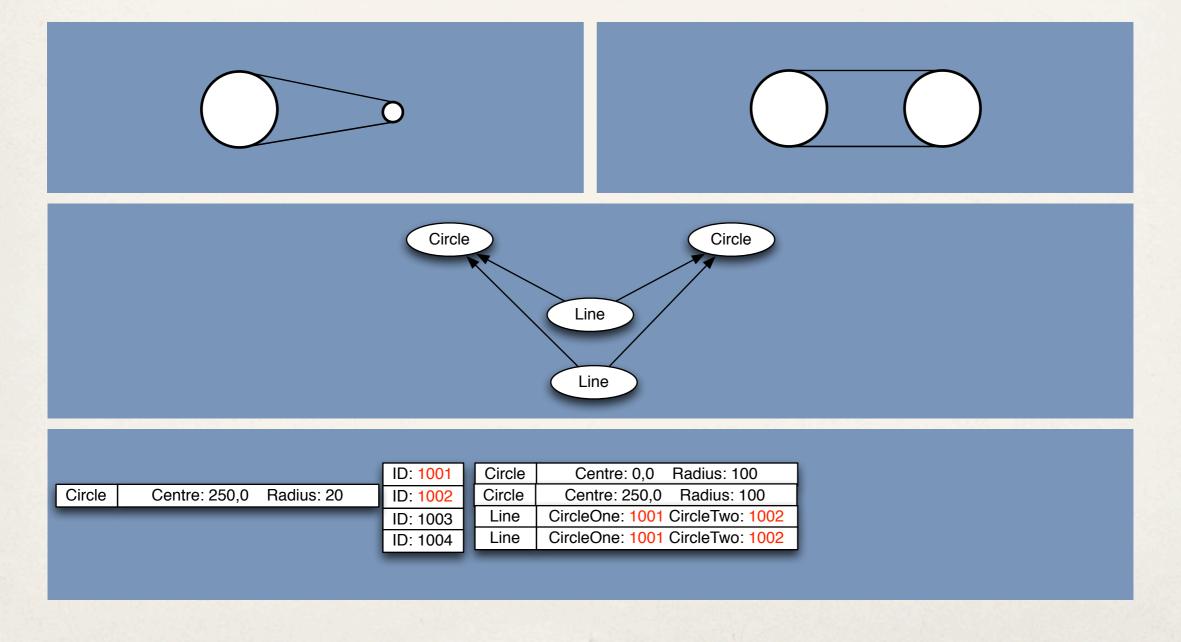
- POINTER EQUALITY pitfall is avoided by:
 - Take it away....
 - put it back if necessary
 - … in a restricted way that makes the code simpler to understand!!
 - Explicit
 - Localized











- IDENTITY pitfall is avoided by:
 - Take it away....
 - put it back if necessary
 - * ... in an explicit way that makes the code simpler to understand!!

Purity - Pitfalls

- 1. Sometimes you really want Side Effects (eg Failure / Mutability)
- 2. Sometimes you really want Pointer Equality
- 3. Sometimes you really want Identity

Purity - Pitfalls

- 1. Sometimes you really want Side Effects (eg Failure / Mutability)
- 2. Sometimes you really want Pointer Equality
- 3. Sometimes you really want Identity
- 4. Learning Curve Need to think in a different (unnatural) way

Type Systems

- Static vs Dynamic Types
- * Benefits...
- Pitfalls

Type Systems - Confusion

- * Statically Typed Languages <u>are</u> Dynamically Typed Languages
 - * ... with an extra feature:
 - * ... any "Classification" you make can be Static (or Dynamic)
 - * ... typically use this in their std libs

Types in Haskell

```
data Bool = True
| I False
```

Types in Haskell

Dynamic Types in Haskell

Strings vs Paths

Strings vs Paths

String

No Distinction

Strings vs Paths

String

No Distinction

appendPath :: (String, String) -> String

Strings vs Paths

String

No Distinction

```
appendPath :: (String, String) -> String
```

appendPath ("\usr\bm", "bar.csv")

Strings vs Paths

String

No Distinction

```
appendPath :: (String, String) -> String
```

appendPath ("line1\n line2\r\n...", "bar.csv")

...no error

Strings vs Paths

String

No Distinction

Runtime Distinction

appendPath :: (Value, Value) -> Value

Strings vs Paths

String

No Distinction

Runtime Distinction

...dynamic error

Strings vs Paths

String

No Distinction

```
data Value = ... | TagString String | I TagPath | String
```

Runtime

Strings vs Paths

String

No Distinction

Runtime

data Path = Path String

Compile time

Strings vs Paths

String

No Distinction

```
data Value = ... | TagString String
| I TagPath String
```

Runtime

data Path = Path String

Compile time

Path "/usr/bm" :: Path

Path "bar.csv" :: Path

Strings vs Paths

String

No Distinction

```
data Value = ... | TagString String
| I TagPath String
```

Runtime

data Path = Path String

Compile time

Path "/usr/bm" :: Path

Path "bar.csv" :: Path

appendPath :: (Path, Path) -> Path

```
Strings vs Paths

String

Adata Value = ... | TagString String | Runtime |

I TagPath String | Compile time
```

```
Path "/usr/bm" :: Path Path "bar.csv" :: Path
```

```
appendPath :: (Path, Path) -> Path
```

```
appendPath ("line1\n line2\r\n...", Path "bar.csv")
...static error
```

```
Strings vs Paths

String

String

No Distinction

data Value = . . . | TagString String | Runtime

| TagPath String | Compile time
```

```
Path "/usr/bm" :: Path Path "bar.csv" :: Path
```

```
appendPath :: (Path, Path) -> Path
```

```
appendPath ( Path "/usr/bm", Path "bar.csv")
```

Strings vs Paths

String

No Distinction

```
data Value = ... | TagString String
| I TagPath String
```

Runtime

data Path = Path String

Compile time

Path "/usr/bm" :: Path

Path "bar.csv" :: Path

```
appendPath :: (Path, Path) -> Path
```

```
appendPath (
```

Path "foo.txt",

Path "bar.csv")

...no error

Strings vs Paths

String

No Distinction

```
data Value = ... | TagString String
               I TagPath String
```

Runtime

data Path = Path String

Compile time

File Paths vs Dir Paths

```
data Path = FilePath String
          I DirPath String
```

Runtime

```
appendPath :: (Path, Path) -> Path
```

appendPath (Path "foo.txt",

Path "bar.csv")

Strings vs Paths

String

No Distinction

data Value = ... | TagString String | I TagPath String

Runtime

data Path = Path String

Compile time

File Paths vs Dir Paths

```
data Path = FilePath String
| DirPath String
```

Runtime

appendPath :: (Path, Path) -> Path

appendPath (FilePath "foo.txt", FilePath "bar.csv")

...dynamic error

```
String
Strings vs Paths
                                                No Distinction
           data Value = ... | TagString String
                                                   Runtime
                            I TagPath String
                 data Path = Path String
                                                 Compile time
File Paths vs
           data Path = FilePath String
                                                   Runtime
  Dir Paths
                      I DirPath String
           data Path
                         = Path String
                                                No Distinction
Path "/usr" :: Path
                            Path "fred.csv" :: Path
     appendPath :: (Path
                              , Path ) -> Path
```

```
String
Strings vs Paths
                                                No Distinction
           data Value = ... | TagString String
                                                  Runtime
                            I TagPath String
                 data Path = Path String
                                                 Compile time
File Paths vs
           data Path = FilePath String
                                                  Runtime
  Dir Paths
                      I DirPath String
           data Path fd = Path String
                                                 Compile time
Path "/usr" :: Path
                            Path "fred.csv" :: Path
     appendPath :: (Path
                              , Path ) -> Path
```

```
String
Strings vs Paths
                                                No Distinction
           data Value = ... | TagString String
                                                  Runtime
                           I TagPath String
                 data Path = Path String
                                                 Compile time
File Paths vs
           data Path = FilePath String
                                                  Runtime
  Dir Paths
                      I DirPath String
           data Path fd = Path String
                                                 Compile time
Path "/usr" :: Path Dir Path "fred.csv" :: Path File
     appendPath :: (Path
                              , Path ) -> Path
```

String Strings vs Paths No Distinction data Value = ... | TagString String Runtime I TagPath String data Path = Path String Compile time File Paths vs data Path = FilePath String Runtime Dir Paths I DirPath String data Path fd = Path String Compile time Path "/usr" :: Path Dir Path "fred.csv" :: Path File appendPath :: (Path Dir, Path fd) -> Path fd

Static Types - Sliding Scale

String Strings vs Paths No Distinction data Value = ... | TagString String Runtime I TagPath String data Path = Path String Compile time File Paths vs data Path = Path String No Distinction Dir Paths data Path = FilePath String Runtime I DirPath String

Compile time

data Path fd = Path String

Static Type Systems - Benefits

- Make code simpler to write
 - detect errors earlier
 - the first thing you write a design language
 - Refactoring, IDEs, Performance

Static Type Systems - Benefits

- Make code simpler to read
 - * Help locate bugs (in type-correct code!)
 - Help understand normal code
 - Help understand very abstract code
 - Document side-effects

```
mychunk :: [Int] -> [Int]
mychunk = ...
```

```
mychunk :: [Int] -> [Int]
mychunk = ...
```

```
mychunk :: [Int] -> [Int]
mychunk = ...
```

```
mychunk :: [a] -> [a] mychunk = ...
```

```
mychunk :: [a] -> [a]
mychunk = ...
```

```
mychunk :: [a] -> [a]
mychunk = ...
```

```
mychunk :: Eq a => [a] -> [a] mychunk = ...
```

```
mychunk :: Num a => [a] -> [a] mychunk = ...
```

```
mychunk :: Typeable a => [a] -> [a] mychunk = ...
```

```
mychunk :: [a] -> [a] mychunk = ...
```

Types: Locating Bugs

```
processWith f = g (f [100,101])

g \times s = \dots
```

Types: Locating Bugs

```
processWith :: (forall a. [a] -> [a]) -> [Int]
processWith f = g (f [100,101])

g :: [Int] -> [Int]
g xs = ...
```

Types: Locating Bugs

```
processWith :: (forall a. [a] -> [a]) -> [Int]
processWith f = g (f [100,101])

g :: [a] -> [a]
g xs = ...
```

```
trans xs ys =
    case xs of
      [] -> []
      ((a,b):xs) -> (map ((x,y) -> (a,y)) $ f ys)
                     ++ trans xs ys
          where
            f \square = \square
            f((c,d):ys) | c==b = (b,d) : f ys
                          l otherwise = f ys
```

```
trans :: [(Int,Int)] -> [(Int,Int)] -> [(Int,Int)]
trans xs ys =
    case xs of
       -> [ ]
      ((a,b):xs) \rightarrow (map ((x,y) \rightarrow (a,y)) $ f ys)
                      ++ trans xs ys
           where
             f \square = \square
             f((c,d):ys) | c==b = (b,d) : f ys
                            l otherwise = f ys
```

```
trans :: Eq b => [(b, b)] \rightarrow [(b, b)] \rightarrow [(b, b)]
trans xs ys =
    case xs of
       -> [ ]
       ((a,b):xs) \rightarrow (map ((x,y) \rightarrow (a,y)) $ f ys)
                        ++ trans xs ys
            where
              f \square = \square
              f((c,d):ys) \mid c==b = (b,d) : f ys
                              l otherwise = f ys
```

```
trans :: [a b \Rightarrow [(b, b)] \rightarrow [(b, b)] \rightarrow [(b, b)]
trans xs ys =
     case xs of
       -> [ ]
       ((a,b):xs) \rightarrow (map ((x,y) \rightarrow (a,y)) $ f ys)
                         ++ trans xs ys
            where
               f \square = \square
               f((c,d):ys) \mid c==b = (b,d) : f ys
                               l otherwise = f ys
```

```
trans :: Eq b => [(a, b)] -> [(b, c)] -> [(a, c)]
trans xs ys =
    case xs of
       -> [ ]
      ((a,b):xs) \rightarrow (map ((x,y) \rightarrow (a,y)) $ f ys)
                      ++ trans xs ys
           where
             f \square = \square
             f((c,d):ys) \mid c==b = (b,d) : f ys
                            l otherwise = f ys
```

```
trans :: Eq b => [(a, b)] \rightarrow [(b, c)] \rightarrow [(a, c)]
trans xs ys =
    case xs of
       -> [ ]
       ((a,b):xs) \rightarrow (map ((x,y) \rightarrow (a,y)) $ f ys)
                        ++ trans xs ys
            where
              f \square = \square
              f((c,d):ys) \mid c==b = (b,d) : f ys
                              l otherwise = f ys
```

```
trans :: Eq b => [(a, b)] -> [(b, c)] -> [(a, c)]
trans xs ys =
  [(a,c) | (a,b1) <- xs, (b2,c) <- ys, b1 == b2]
```

Types: Handling very abstract code

```
What is 'h'?
ala w h f = unwrap . h (w . f)
```

Types: Handling very abstract code

```
ala :: Newtype srb =>
        (b -> sb) ->
       ((a -> sb) -> ta -> srb) ->
        (a -> b) ->
       (ta -> Unwrap srb)
ala w h f = unwrap . h (w . f)
```

Types: Handling very abstract code

```
ala :: Newtype srb =>
        (b -> sb) ->
       ((a -> sb) -> ta -> srb) ->
        (a -> b) ->
       (ta -> Unwrap srb)
ala w h f = unwrap . h (w . f)
```

```
do
    y <- chunk1(a,b)
    z <- chunk1(a,b)
    chunk3(x)
    chunk4(y,z)</pre>
```

```
chunk0 a = do
    y <- chunk1(a,b)
    z <- chunk1(a,b)
    chunk3(x)
    chunk4(y,z)
    return z</pre>
```

```
chunk0 :: Int -> ReaderT Config
                      (StateT Connection
                       (EitherT DBError Int))
chunk0 a = do
            y \leftarrow chunk1(a,b)
            z \leftarrow chunk1(a,b)
            chunk3(x)
            chunk4(y,z)
            return z
```

Type Systems - Benefits

- * Make code simpler to read
 - Help locate bugs (in type-correct code!)
 - Help understand normal code
 - Help understand very abstract code
 - Document side-effects
- * Make code simpler to UNDERSTAND !!!!!!!

Type Systems - Pitfalls

Type Systems - Pitfalls

- Longer to get running code
- Changing your mind
- Learning Curve "A Second Language"

Summary

- * Functional Programming makes code **SIMPLER**
 - PURITY + TYPES are the two biggest ways it does this

End